

SKETCH OF SOLUTIONS (HOMEWORK II)

- 4.- a) yes, b)no, c) yes, d) yes, e) yes, f) no
16.- Suppose the division algorithm yields $a = qb + r$, $q = ct + s$ with $r < b$, $s < c$. Then substituting we get:

$$a = (cb)t + (sb + r)$$

If we show that $sb + r < cb$ we are done (*why?*). Notice $s < c \Rightarrow s + 1 \leq c$
Thus $(s + 1)b \leq cb$ but since $r < b$ we get $sb + r < sb + b \leq cb$

- 28.- We must require $m, n > 0$

$$\left[\frac{x+n}{m} \right] = \left[\frac{[x+n]}{m} \right] = \left[\frac{[x] + n}{m} \right]$$

The first equality comes from example 1.34 and the second from example 1.31

Section 3.1

- 2.- a) not prime, b) not prime, c) not prime, d) prime, e) not prime, f) not prime
6.- Notice $n^3 + 1 = (n+1)(n^2 - n + 1)$ Therefore $n^3 + 1$ is prime iff $n+1 = n^3 + 1$ (and $n+1 > 1$) iff $n = 1$
10.- Let p_i be a prime in the list. Without loss of generality we can say p_i divides Q and p_i does not divide R . But then p_i cannot divide $Q + R$. Since $Q + R > p_n$ there must be more than n primes.

Section 3.2

- 6.- If $b|a$ and $b|a+2$ then $b|a+2-a$ therefore $\gcd(a, a+2)$ is either 1 or 2. If a is even then $\gcd(a, a+2) = 2$
7.- By theorem 3.8 we know (ca, cb) is the least positive integer of the form $cma + cmb = |c||ma + nb|$ therefore $|ma + nb|$ is minimum i.e. (ma, nb)
31.- Suppose $\frac{ad+bc}{bd} = k$ then $kbd = ad + bc$ therefore $d|bc$. But $d \nmid c$, therefore we must have $d|b$. Analogously $b|d$