

# 1 What's on my whiteboard

## 1.1 left panel

$$(S, \sigma) \xrightarrow[f]{e} (T, \tau)$$

$$W(\sigma) \in \mathcal{WAD}(S)$$

$$W(\tau) \in \mathcal{WAD}(T)$$

$$W(\sigma) = f^*(W(\tau))$$

$$W(\sigma) + B \multimap W(\tau)$$

$$\|B\|_\infty \leq 2$$

## 1.2 middle

- $\mathcal{WAD}(S)$

$$f^*, \langle, \rangle, +, \| \cdot \|_\infty$$

- Path-family

- neat
- canonical
- valid

## 1.3 left panel

### 1 Lemma

If  $X \in \mathcal{WAD}(S)$  is valid for  $\sigma$ , then  $\exists B, \|B\|_\infty \leq 2$  such that

$$W_S(\sigma) + B \geq X.$$

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**2 Lemma**

If  $Y \in \mathcal{WAD}(T)$  is valid for  $\tau$ , and  $e : (S, \sigma) \rightarrow (T, \tau)$  is a conformal embedding, then  $\exists X \in \mathcal{WAD}(S)$  valid for  $\sigma$  such that

$$X \stackrel{e}{\dashv} Y.$$

## 2 What's on Misha's blackboard

### 2.1 The Degeneration Theorem

#### 3 Theorem

$$(S, \sigma_i) \xrightarrow[e]{f} (T, \tau_i)$$

If  $[\tau_i] \rightarrow \infty$  in  $\text{Teich}^*(T)$  then, letting  $W_i = W(T, \tau_i)$ , then  $\|W_i\|_\infty \rightarrow \infty$ , and, passing to a subsequence,

$$\frac{W_i}{\|W_i\|_\infty} \rightarrow \infty,$$

with

$$f^*(W_\infty) \xrightarrow{e} W_\infty.$$

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### 2.2 $B$ -invariant arc diagram, $B \in \mathbb{Z}^+$

$X \in \mathcal{AD}(T)$  is  $B$ -invariant if  $\forall [\alpha] \in X$ ,  $\exists [\alpha_1], \dots, [\alpha_n] \in f^*(X) (\in \mathcal{AD}(S))$  such that  $\alpha_1, \dots, \alpha_n \xrightarrow{e} \alpha$ .

$$f^*W \multimap \frac{1}{2}W \Rightarrow \text{supp}(W) \text{ is } B\text{-invariant}$$

### 2.3 The Arc Lemma

#### 4 Lemma

Given

$$S \xrightarrow[e]{[f]} T$$

and  $B$ , there exist finitely many  $B$ -invariant arc diagrams for  $(S, T, [e], [f])$ .

### 3 Lollypop

Given

$$S \xrightarrow{e} T$$

we say

$$\alpha_1, \dots, \alpha_n \rightarrow \alpha$$

if

$$e^{-1}(\alpha) = \bigcup_{i=1}^n \alpha_i \cup \{\text{homotopically trivial arcs}\}$$

(more precisely, if  $e^{-1}(\alpha)$  is  $\alpha_1, \dots, \alpha_n$  *in sequence*). Then

$$\sum_{i=1}^n w_i \cdot \alpha_i \multimap w \cdot \alpha$$

if

$$\alpha_1, \dots, \alpha_n \rightarrow \alpha$$

and

$$\sum_{i=1}^n \frac{1}{w_i} \leq \frac{1}{w}.$$

And

$$A + \sum A_j \multimap \sum B_j$$

if  $A \geq 0$  and  $A_j \multimap B_j$ .

## 4 Weighted arc diagrams

Let  $S$  be a compact surface with (non-trivial) boundary. An *arc* on  $S$  is an embedding  $(I, \partial I) \rightarrow (S, \partial S)$ , up to isotopy (that is not isotopic to a constant map). We denote the set of arcs on  $S$  by  $\mathcal{A}(S)$ . Any two arcs  $\alpha, \beta \in \mathcal{A}(S)$  have an intersection number  $\langle \alpha, \beta \rangle$  that is equal to the minimum number of times that representatives of  $\alpha$  and  $\beta$  intersect. We say that two arcs are *disjoint* if they have zero intersection number.

An *arc-diagram* is a (necessarily finite) set of arcs in  $S$  that are pairwise disjoint. We denote the set of arc-diagrams by  $\mathcal{AD}(S)$ .

A *weighted arc-diagram* is a arc-diagram with positive real weights assigned to each arc. We denote the set of weighted arc-diagrams by  $\mathcal{WAD}(S)$ . If  $X \in \mathcal{WAD}(S)$ , then the *support* of  $X$ , denoted  $\text{supp}(X)$ , is the underlying arc-diagram. If  $X, Y \in \mathcal{WAD}(S)$ , and  $\text{supp } X \cup \text{supp } Y \in \mathcal{AD}(S)$  (i.e. no arc for  $X$  intersects one for  $Y$ ) then we can form the weighted arc-diagram  $X + Y$  by adding the weights of arcs that appear in both. We say that  $X \geq Y$  if  $\exists Z \in \mathcal{WAD}(S)$  such that  $Y + Z = X$ .

Given  $X \in \mathcal{WAD}(S)$ , we can write

$$X = \sum_i w_i \alpha_i$$

where the  $\alpha_i$  are distinct. Then we write  $X|_\alpha = w_i$  if  $\alpha = \alpha_i$ , and  $X|_\alpha = 0$  if  $\alpha \notin \text{supp } X$ . (Note that  $X \geq Y$  if and only if  $X|_\alpha \geq Y|_\alpha$  for all  $\alpha \in \mathcal{A}(S)$ ). Also, we write

$$\|X\|_\infty = \sup_{\alpha \in \text{supp } X} X|_\alpha,$$

and

$$\|X\|_1 = \sum_{\alpha \in \text{supp } X} X|_\alpha.$$

If  $f : S \rightarrow T$  is a covering map, and  $Y \in \mathcal{WAD}(T)$ , we define  $f^*(Y) \in \mathcal{WAD}(S)$  by

$$f^*(Y)|_\alpha = Y|_{f_*\alpha}$$

(of course, if  $f_*\alpha$  is not an embedded arc, then  $Y|_{f_*\alpha} = 0$ ).

## 5 Neat path-families and weight

Let  $S$  now be a bordered Riemann surface. A *path* in  $S$  is a piecewise smooth map  $a : (I, \partial I) \rightarrow (S, \partial S)$  (*not* up to isotopy). A *path family* on  $S$  is a set of paths. A path family is *neat* if

- every path is an embedding,
- every two paths have disjoint images, and
- no path is homotopic (through paths) to a constant map.

Given any neat path family  $\mathcal{F}$  and arc  $\alpha \in \mathcal{A}(S)$  we define  $\mathcal{F}|_\alpha$  to be the set of paths in  $\mathcal{F}$  that are representatives of  $\alpha$ , and we let  $\text{supp } \mathcal{F} \in \mathcal{AD}(S)$  be those  $\alpha$  for which  $\mathcal{F}|_\alpha$  is non-empty. We then define  $W(\mathcal{F}) \in \mathcal{WAD}(S)$  by setting  $W(\mathcal{F})|_\alpha$  to be the reciprocal of the *extremal length* of  $\mathcal{F}|_\alpha$ ; we then have  $\text{supp } W(\mathcal{F}) = \text{supp } \mathcal{F}$ . We say that  $X \in \mathcal{WAD}(S)$  is *valid* for a given Riemann surface structure on  $S$  if there exists a neat path-family  $\mathcal{F}$  such that  $W(\mathcal{F}) = X$ .

## 6 The Canonical neat path-family

Now consider the set of all bordered Riemann surfaces. We can assign to each surface  $S$  a neat path-family, called the *canonical* neat path-family  $\mathcal{F}^N(S)$  in such a way that

- if  $\mathcal{F}'$  is any neat path-family on  $S$ , then

$$W(\mathcal{F}') \leq W(\mathcal{F}^N(S)) + B$$

for some  $B \in \mathcal{WAD}(S)$  with  $\|B\|_\infty \leq 2$ ;

- if  $f : S \rightarrow T$  is a conformal covering map, then  $\mathcal{F}^N(S) = f^*(\mathcal{F}^N(T))$ , in the sense that  $a \in \mathcal{F}^N(S)$  iff  $f \circ a \in \mathcal{F}^N(T)$ .

We define  $W(S) = W(\mathcal{F}^N(S)) \in \mathcal{WAD}(S)$ .

## 7 The inclusion lemma

### 5 Lemma

If  $Y \in \mathcal{WAD}(T)$  is valid for  $\tau$ , and  $e : (S, \sigma) \rightarrow (T, \tau)$  is a conformal embedding, then  $\exists X \in \mathcal{WAD}(S)$  valid for  $\sigma$  such that

$$X \overset{e}{\dashv} Y.$$



## 8 Almost maximality

### 6 Lemma

If  $X \in \mathcal{WAD}(S)$  is valid for  $\sigma$ , then  $\exists B, \|B\|_\infty \leq 2$  such that

$$W_S(\sigma) + B \geq X.$$