

Planar maps with at most six neighbors on average

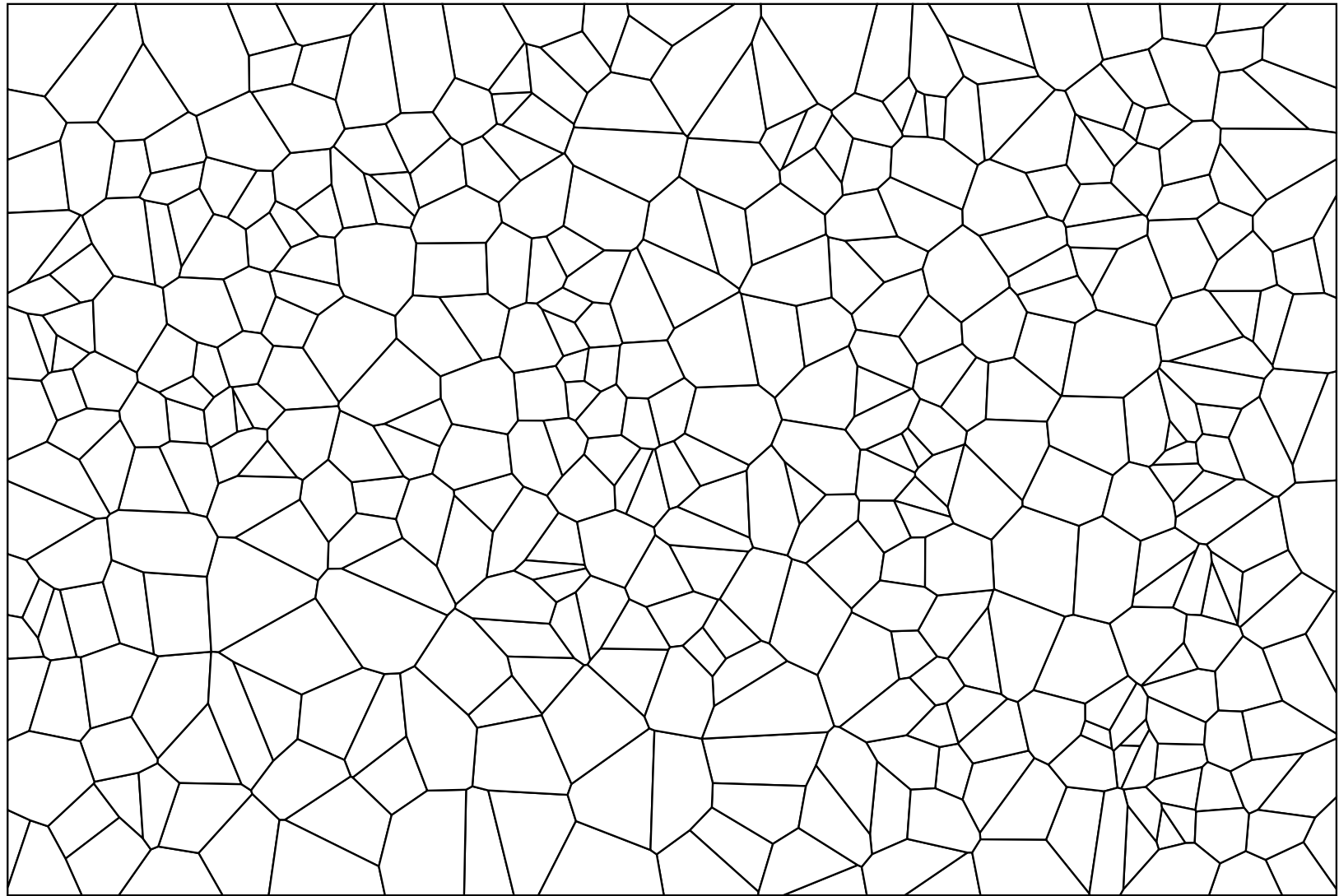
Christopher Bishop, Stony Brook

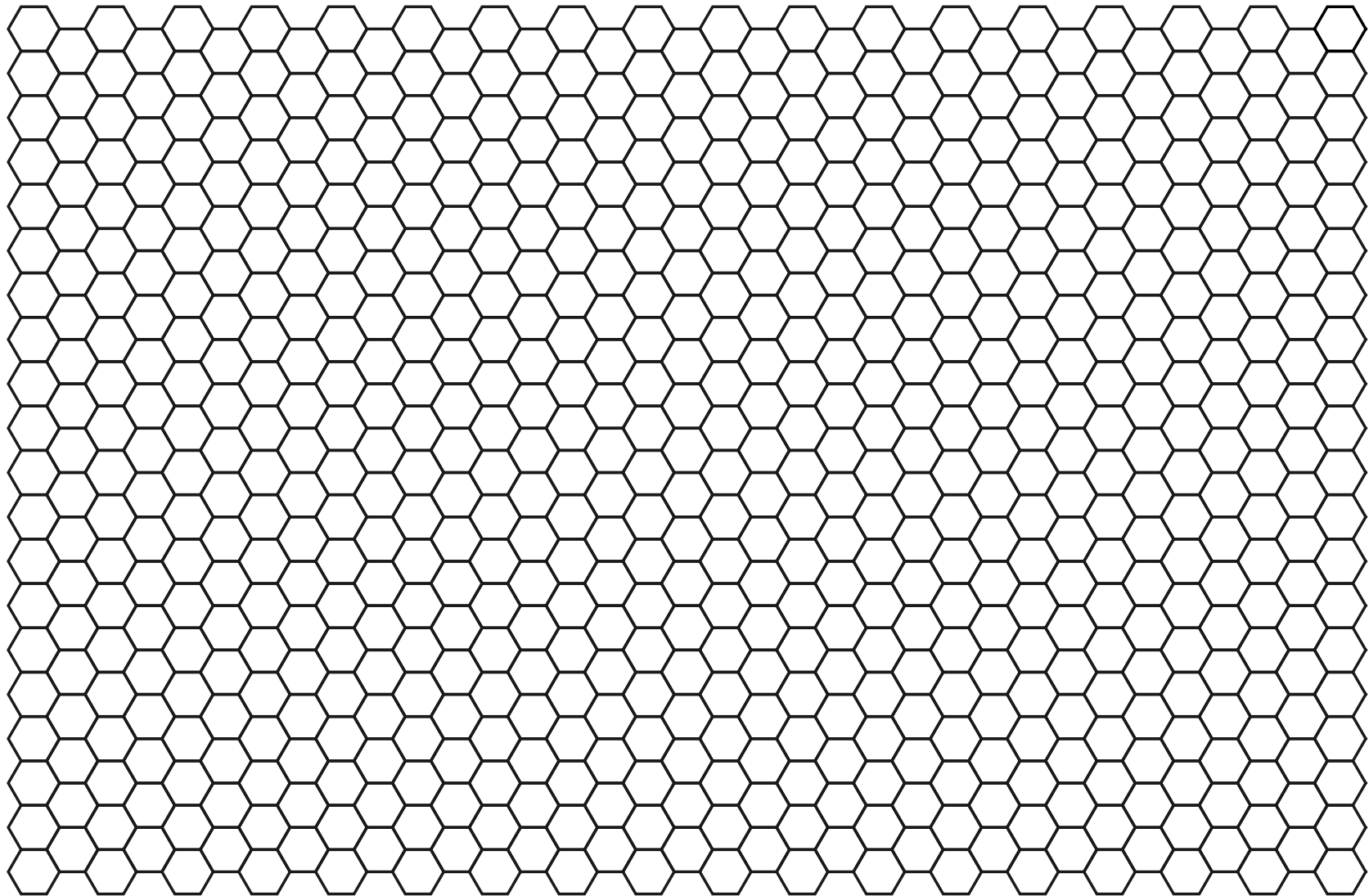
Dennis Sullivan, Stony Brook

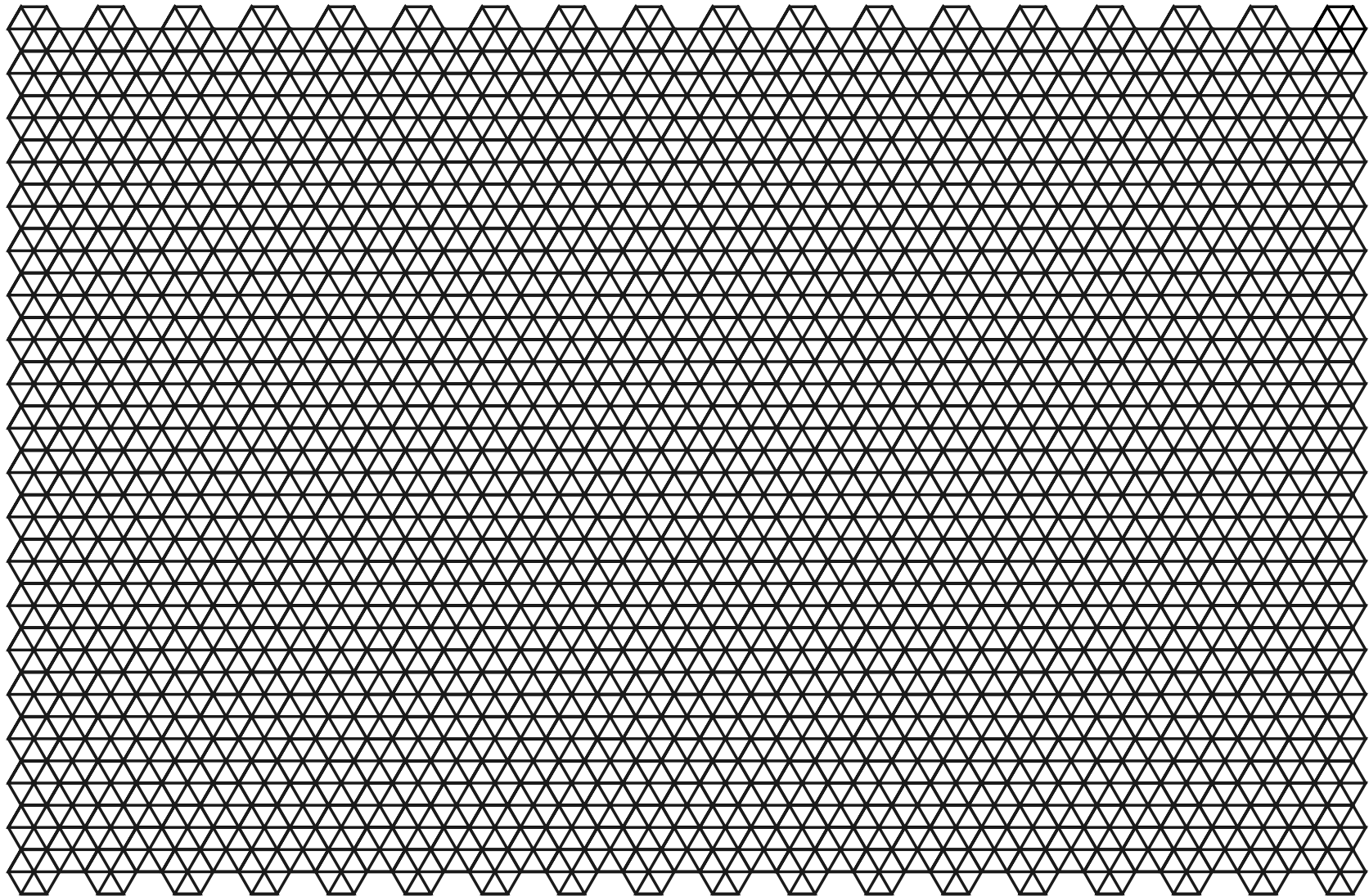
Michael Wigler, Cold Spring Harbor

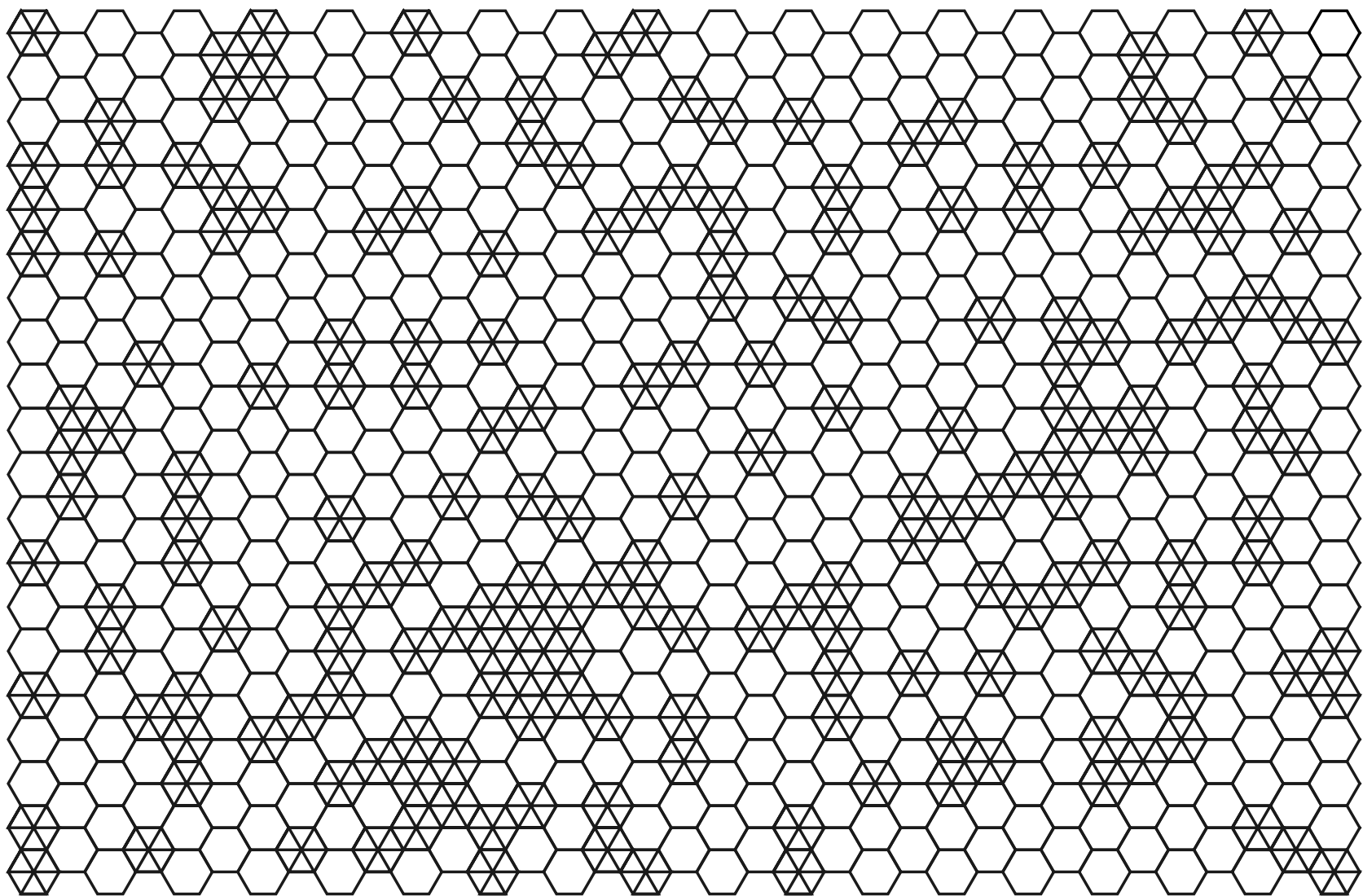
AMS Sectional Meeting – March 19, 2016 Stony Brook NY

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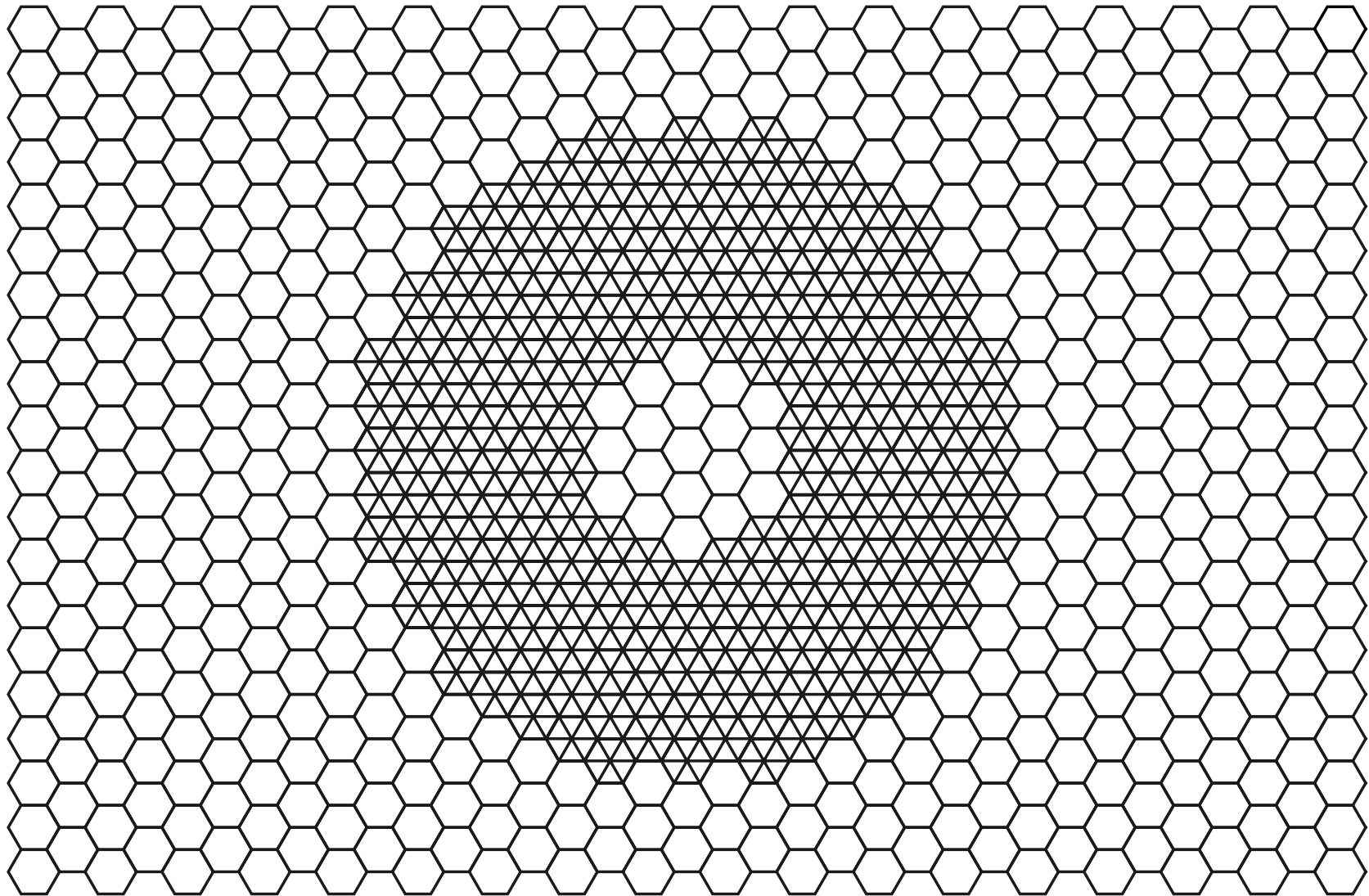




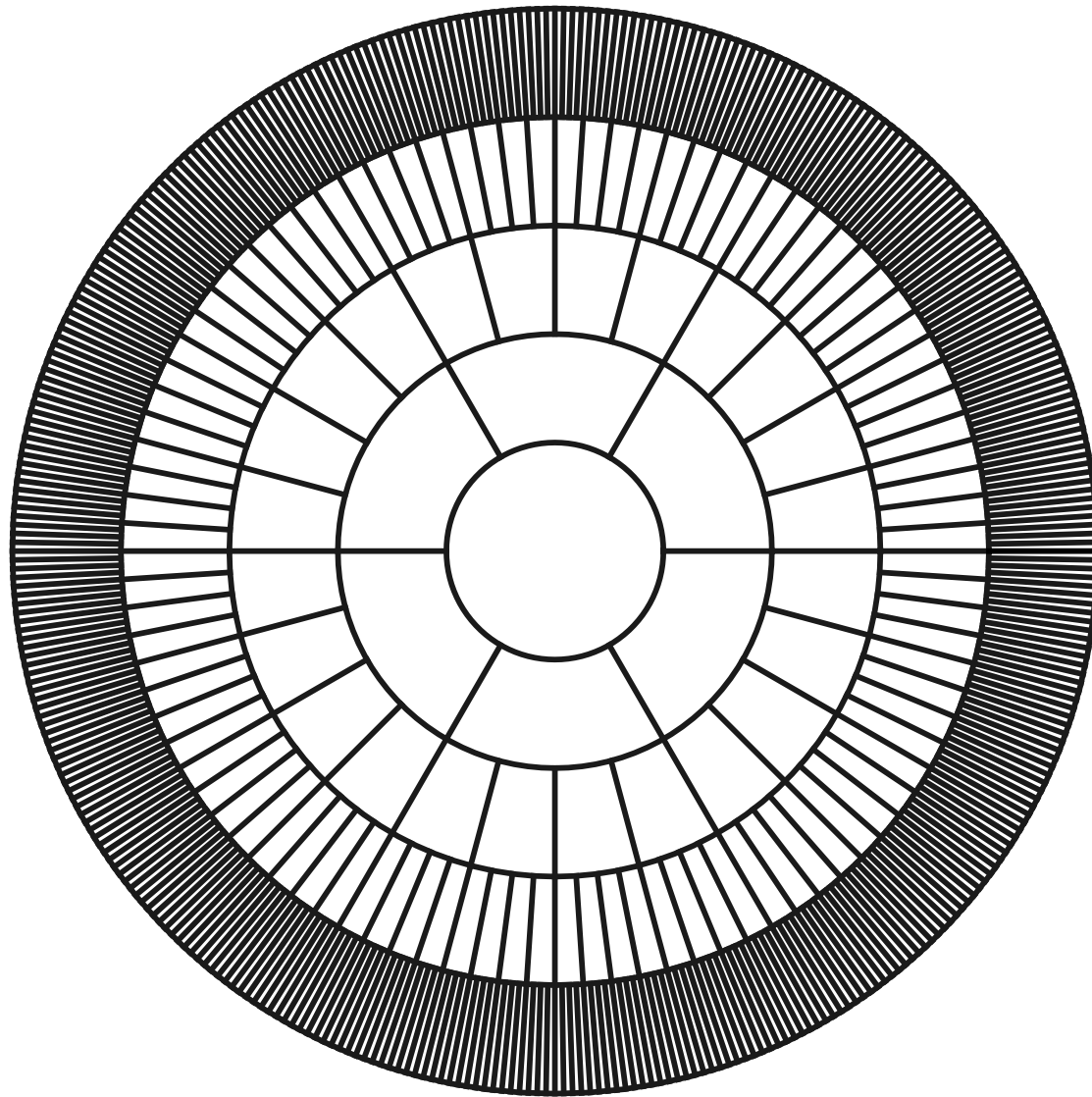




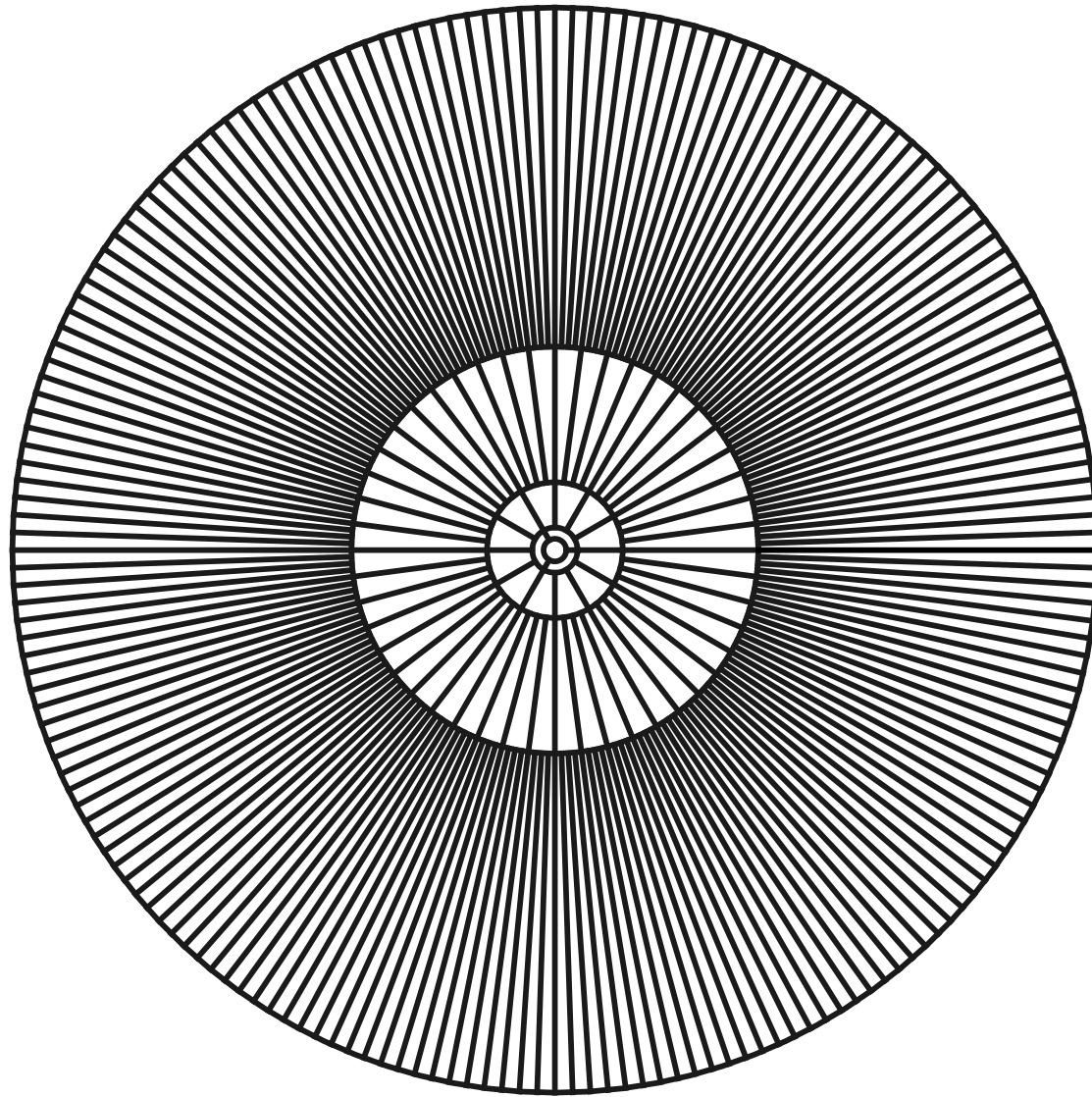
Hexagon \rightarrow triangles with probability p , ANS = $\frac{12p+6}{5p+1}$ ($p = 1/4$ shown).



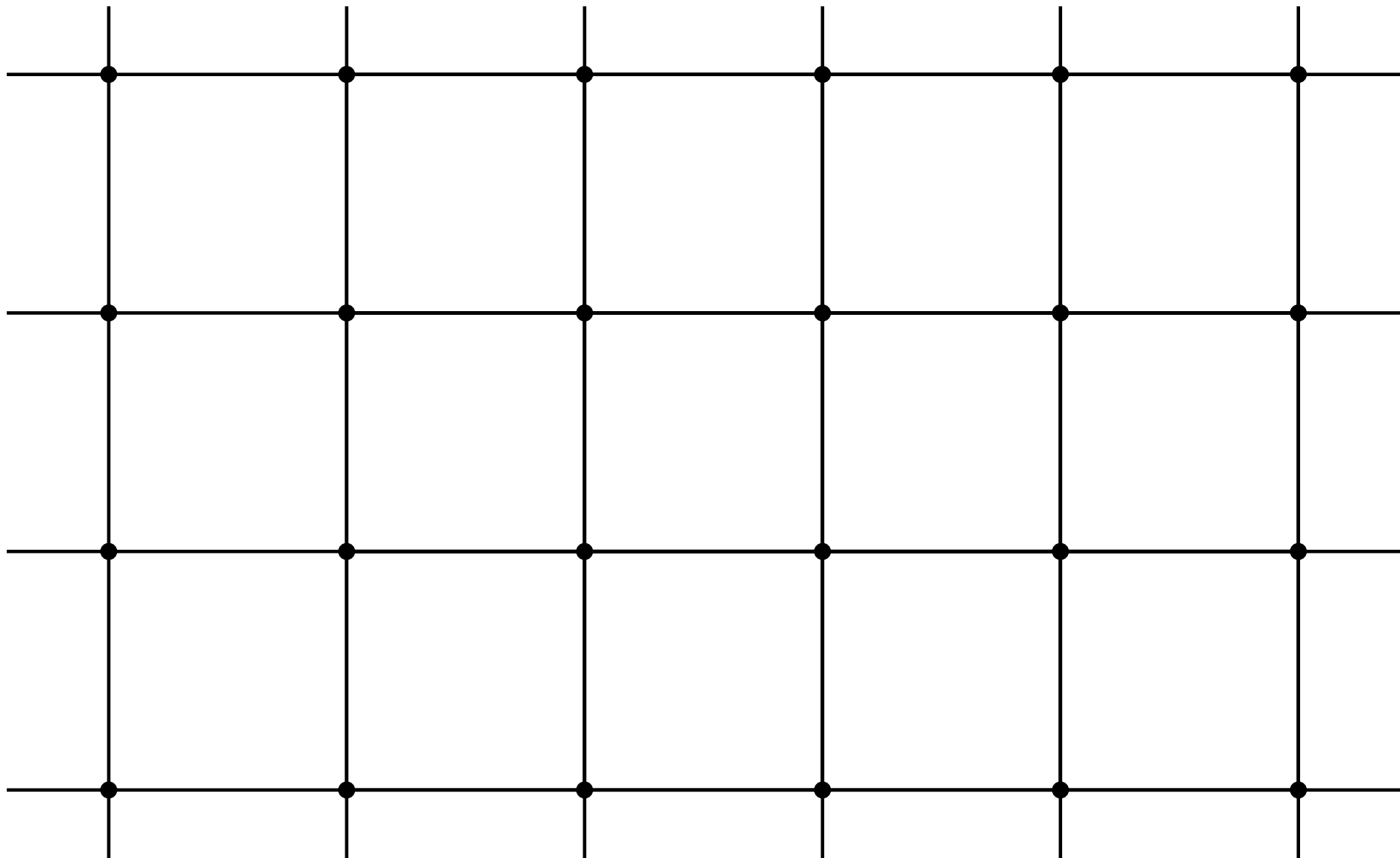
Limit need not exist

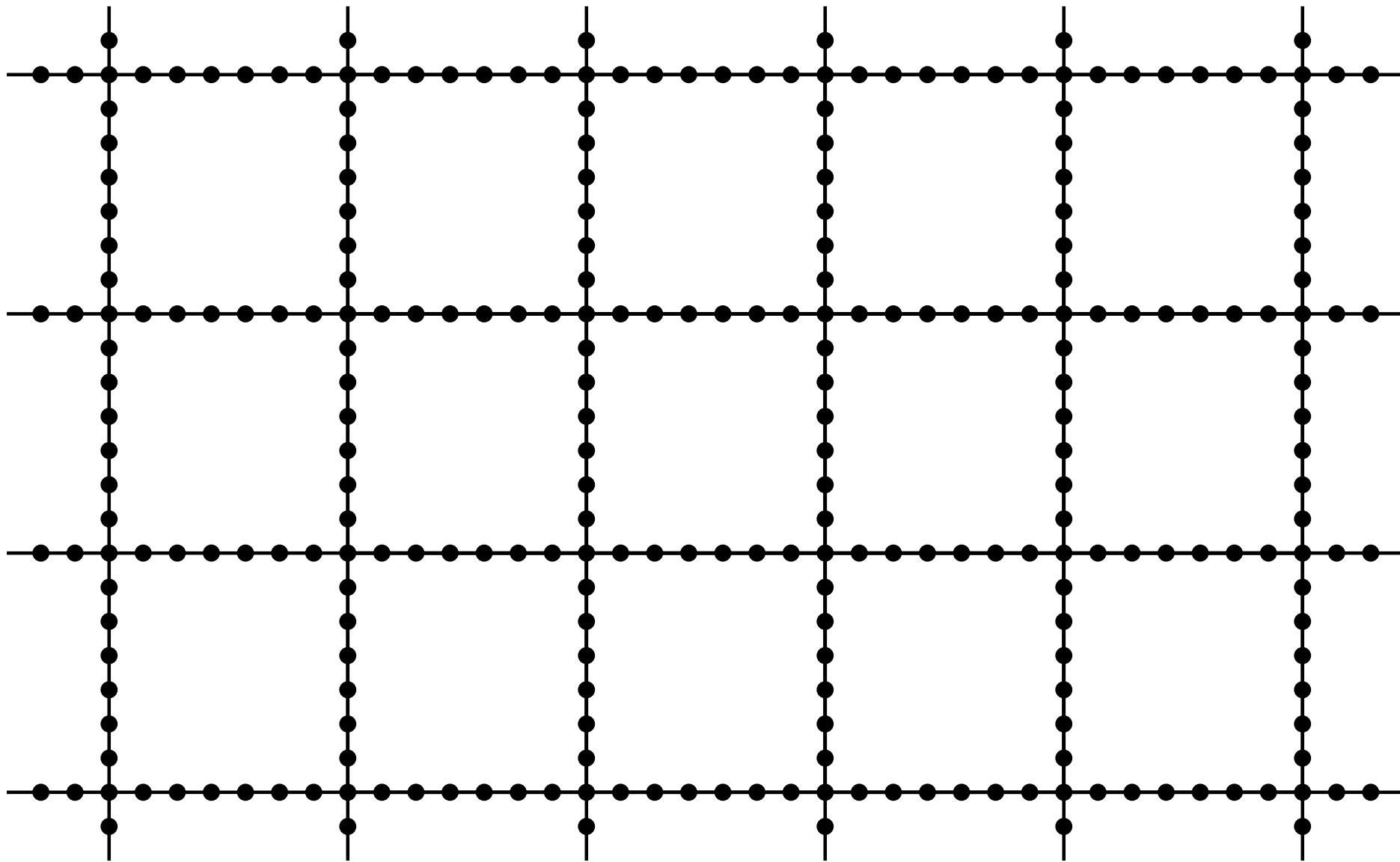


average number sides $\nearrow \infty$, areas $\rightarrow 0$, diameters bounded



average number sides $\nearrow \infty$, diameters $\rightarrow \infty$, areas bounded below



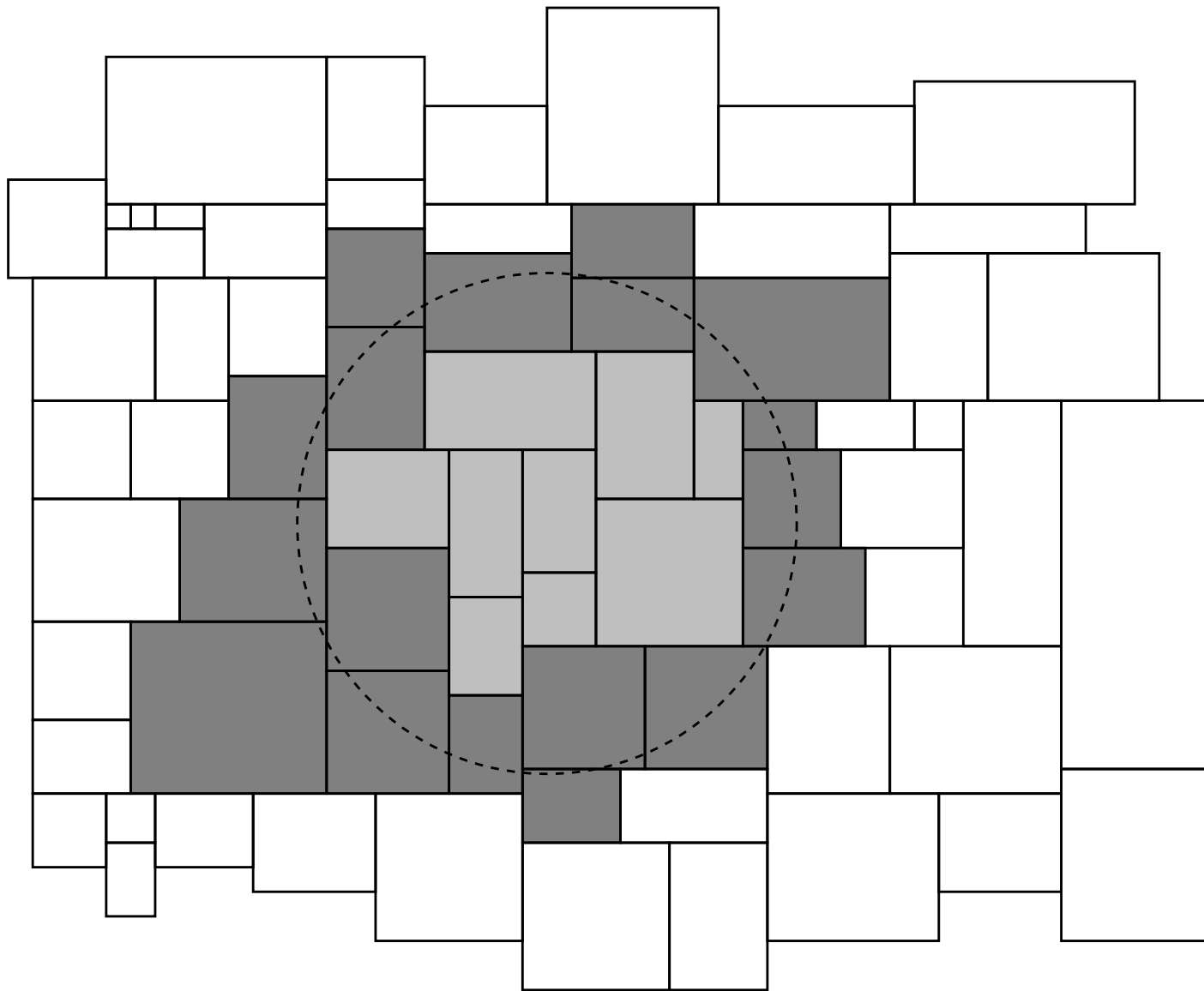


average number sides $\nearrow \infty$, forbid vertices of degree 1 and 2

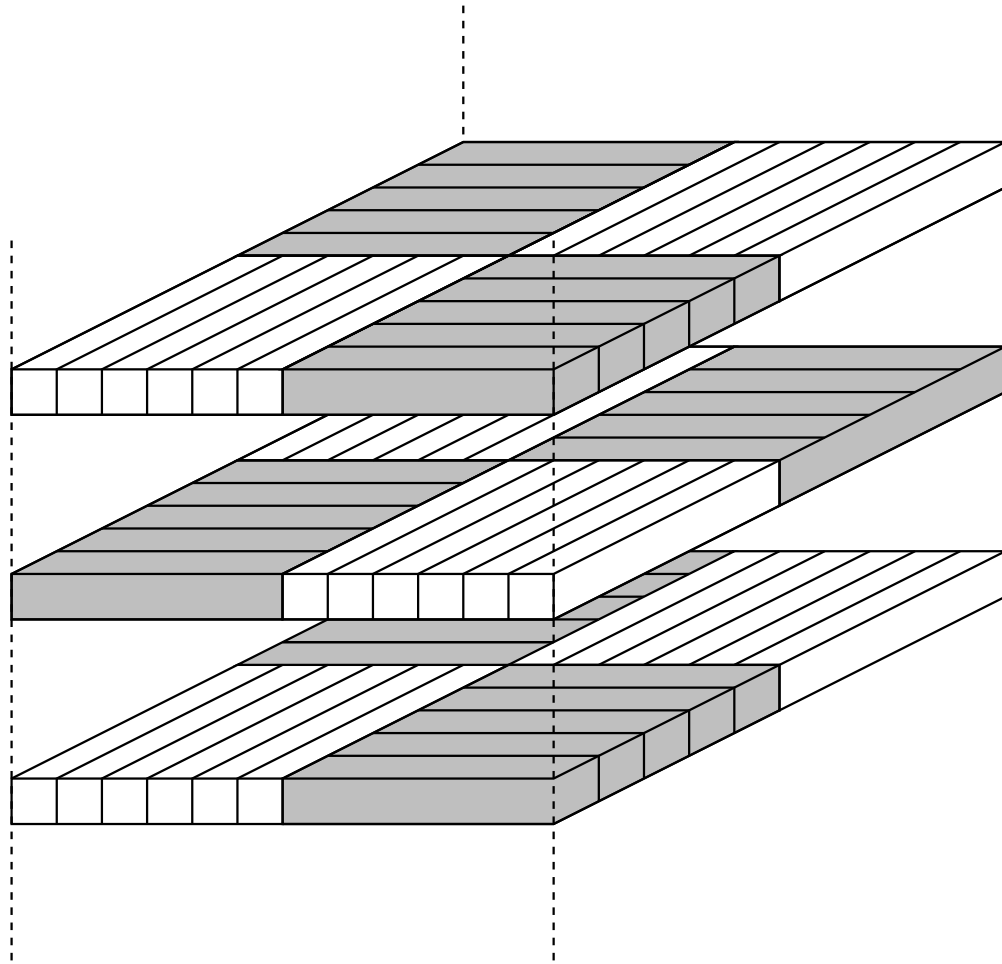
Theorem 1: Suppose faces have areas that are bounded below and diameters that are bounded above, and every vertex had degree ≥ 3 . Then

$$\limsup_{t \rightarrow \infty} \text{ANS}(t) \leq 6.$$

Average is over all faces contained in circle of radius t .



Average over faces inside/hitting an expanding circle



Theorem false in three dimensions

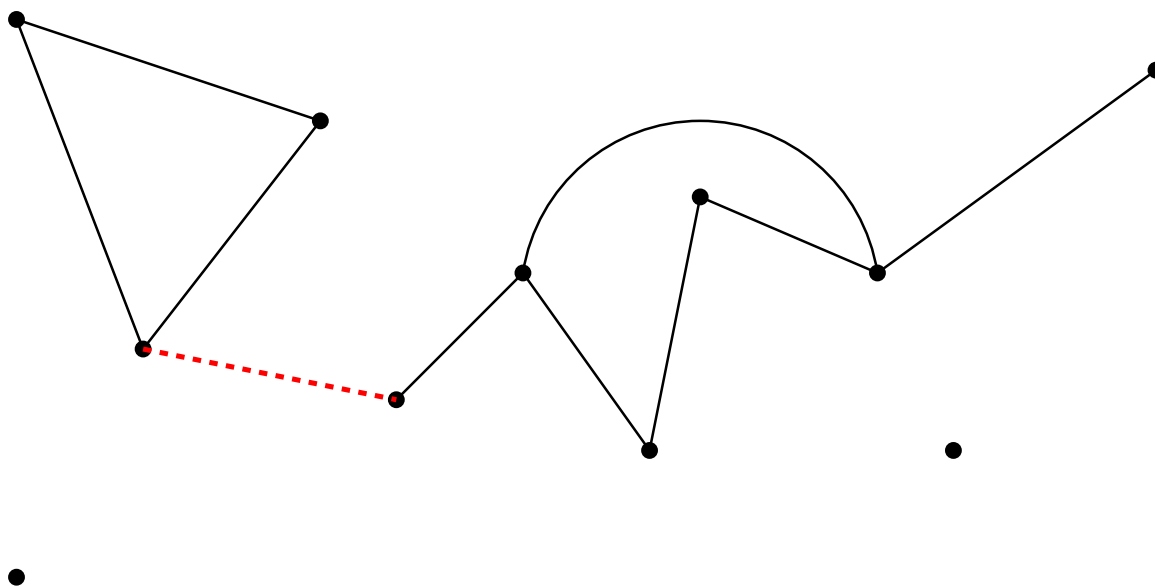
Euler's formula for finite planar graphs: $V - E + F = C + 1$

V = number vertices

E = number of edges

F = number of faces

C = number of components



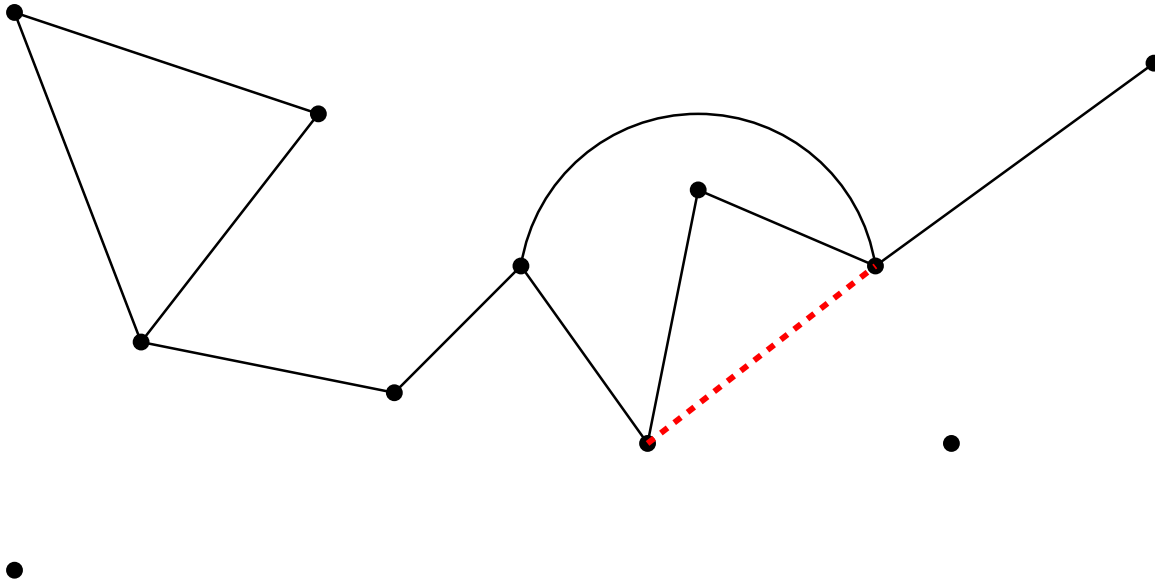
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Sphere Theorem: If H is a finite graph on the 2-sphere and every vertex has degree ≥ 3 , then $E < 3F$.

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Proof: Clearly

$$\sum_{v \in H} \deg(v) = 2E.$$

By assumption, $\deg(v) \geq 3$ for all vertices so $3V \leq 2E$. Hence $V \leq \frac{2}{3}E$.

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Plugging this into Euler's formula $V - E + F = C + 1$ gives

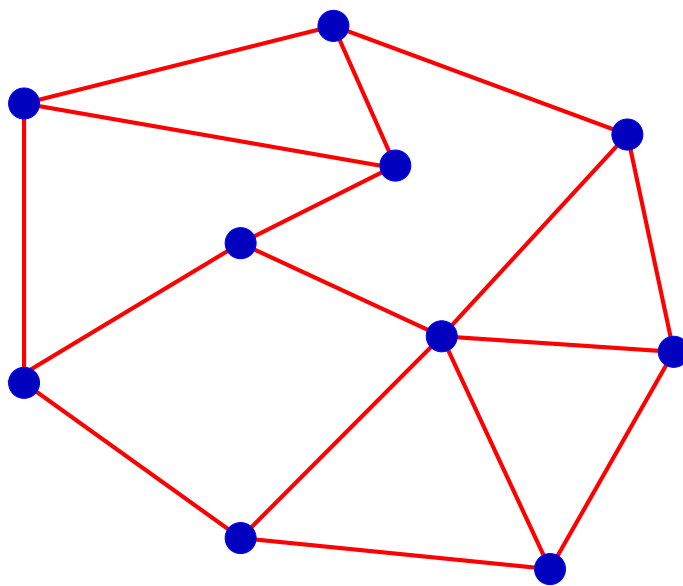
$$-\frac{1}{3}E + F \geq C + 1$$

or

$$E \leq 3F - 3(C + 1) < 3F.$$

Sphere Theorem: If H is a finite graph on the 2-sphere and every vertex has degree ≥ 3 , then $E < 3F$.

Corollary: Number of sides = $2E < 6F$. So for a finite planar graph (all degrees ≥ 3), the average number of sides per face is ≤ 6 .



$$E = 16, \quad F = 8, \quad \text{ANS} = 4$$

This is for finite H . What about infinite maps?

Choose a piecewise smooth region R in the plane and let $tR + x$ denote the region dilated by a factor of $t > 0$ and translated by x .

Let $H = H(R, t, x)$ be the sub-map of G consisting of the 2-cells of G that lie inside $tR + x$.

Let $\text{ANS}(t)$ be average number of sides over faces in $H(R, t, x)$.

Usually take circle around origin.

Theorem 1: Suppose all faces of G have diameter $\leq D < \infty$, and have area $\geq A > 0$, and that every vertex has degree ≥ 3 . Then

$$\limsup_{t \rightarrow \infty} \text{ANS}(H(R, t, x)) \leq 6.$$

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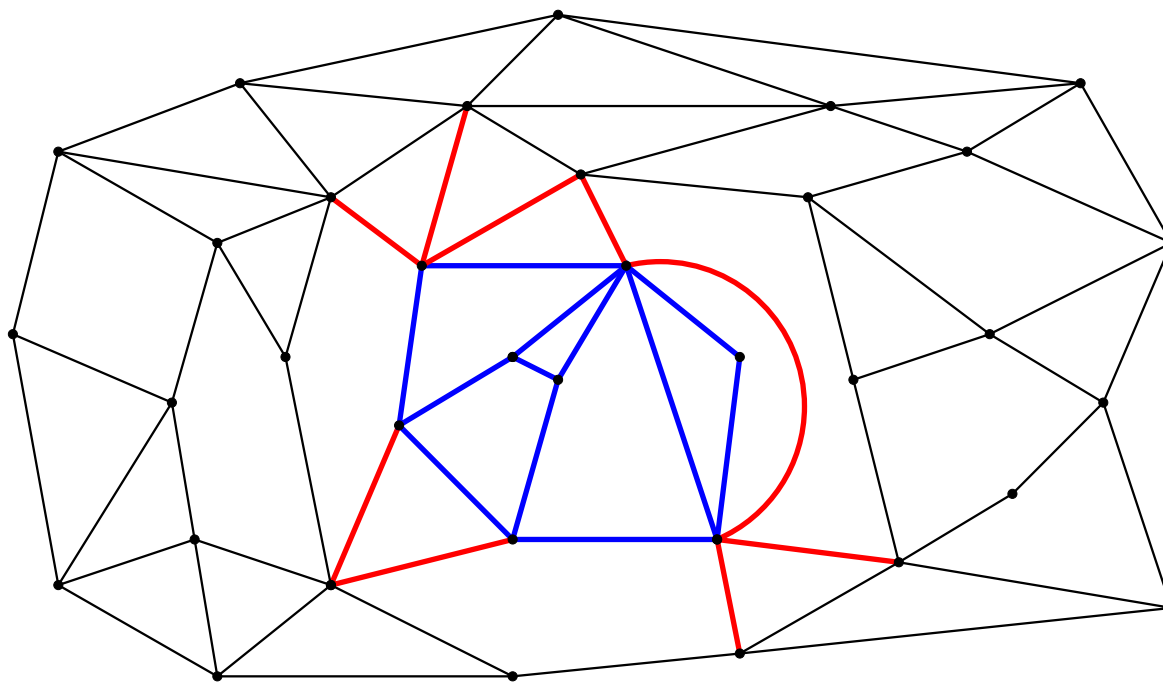
Average is taken over faces contained in $tR + x$.

Same conclusion holds for faces hitting $tR + x$

Really only need R to have non-empty interior and zero area boundary.

Definition: edge boundary: If $H \subset G$, then $\partial_E H$ is set of edges in $G \setminus H$ with at least one endpoint in H .

$$e(H) = \frac{|\partial_E H|}{F}.$$



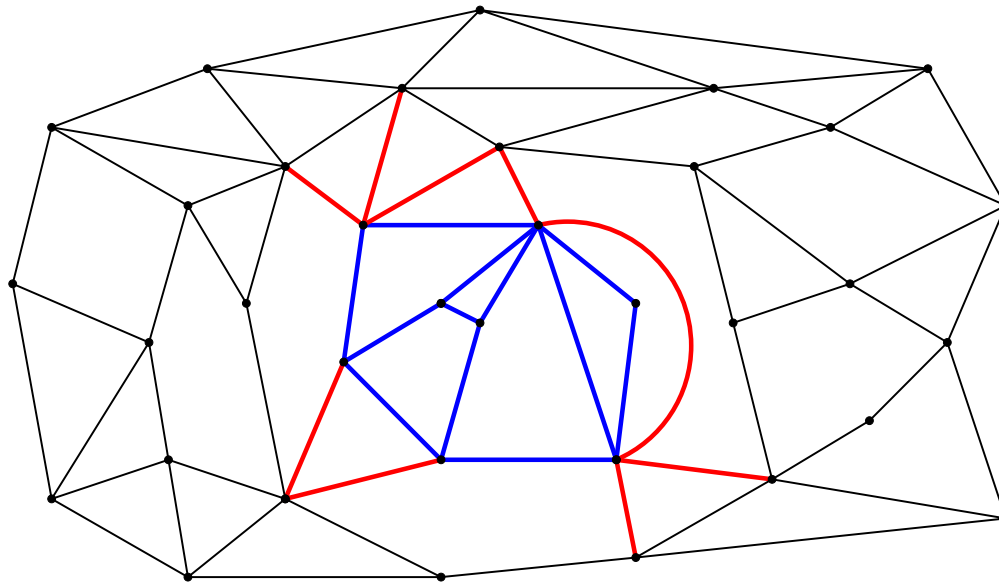
Here $e(H) = 9/5$. Related to Cheeger constant of G .

Lemma 1: If H is a sub-map of G then $\text{ANS}(H) \leq 6 + 4e(H)$.

Proof: Note that

$$2E + |\partial_E H| \leq \sum_{v \in H} \deg_G(v) \leq 2E + 2|\partial_E H|,$$

The left hand inequality would be an equality, except that some edges in $\partial_E H$ might have both endpoints on ∂H .



Lemma 1: If H is a sub-map of G then $\text{ANS}(H) \leq 6 + 4e(H)$.

Proof continued: Every vertex has degree ≥ 3 , so

$$3V \leq \sum_{v \in H} \deg_G(v) \leq 2E + 2|\partial_E H|.$$

Divide by 3 and put estimate for V into Euler's formula:

$$\left(\frac{2}{3}E + \frac{2}{3}|\partial_E H|\right) - E + F \geq C$$

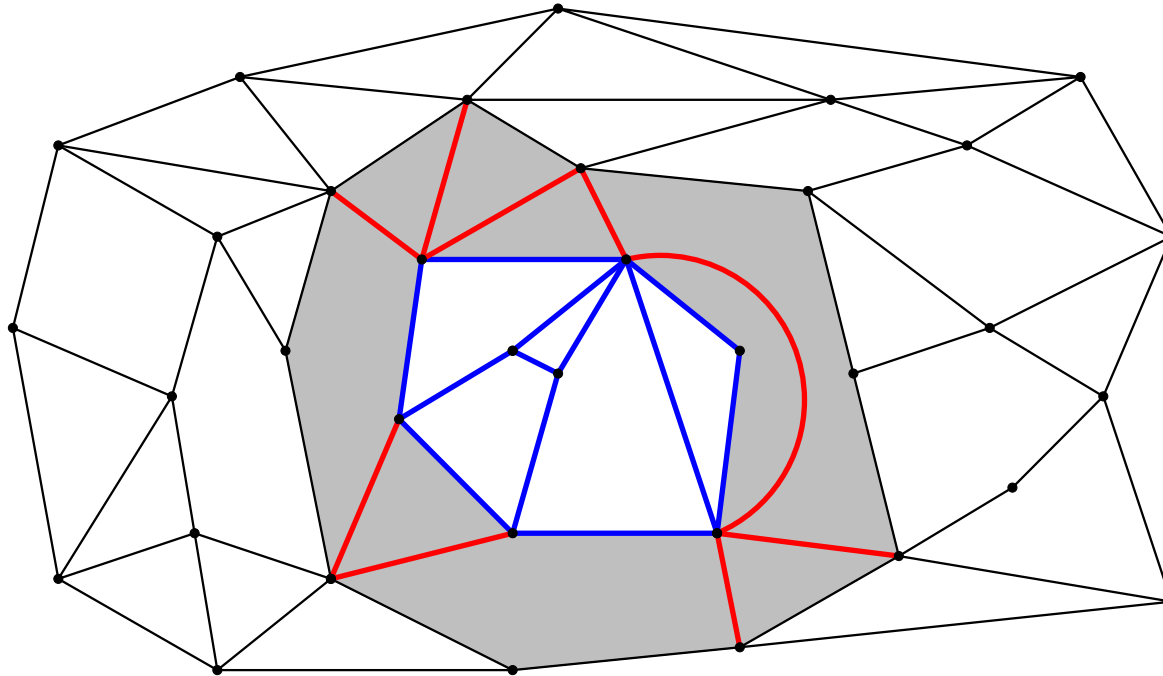
Simplifying gives

$$\frac{E}{F} \leq 3 + 2\frac{|\partial_E H|}{F} - 3\frac{C}{F} \leq 3 + 2e(H).$$

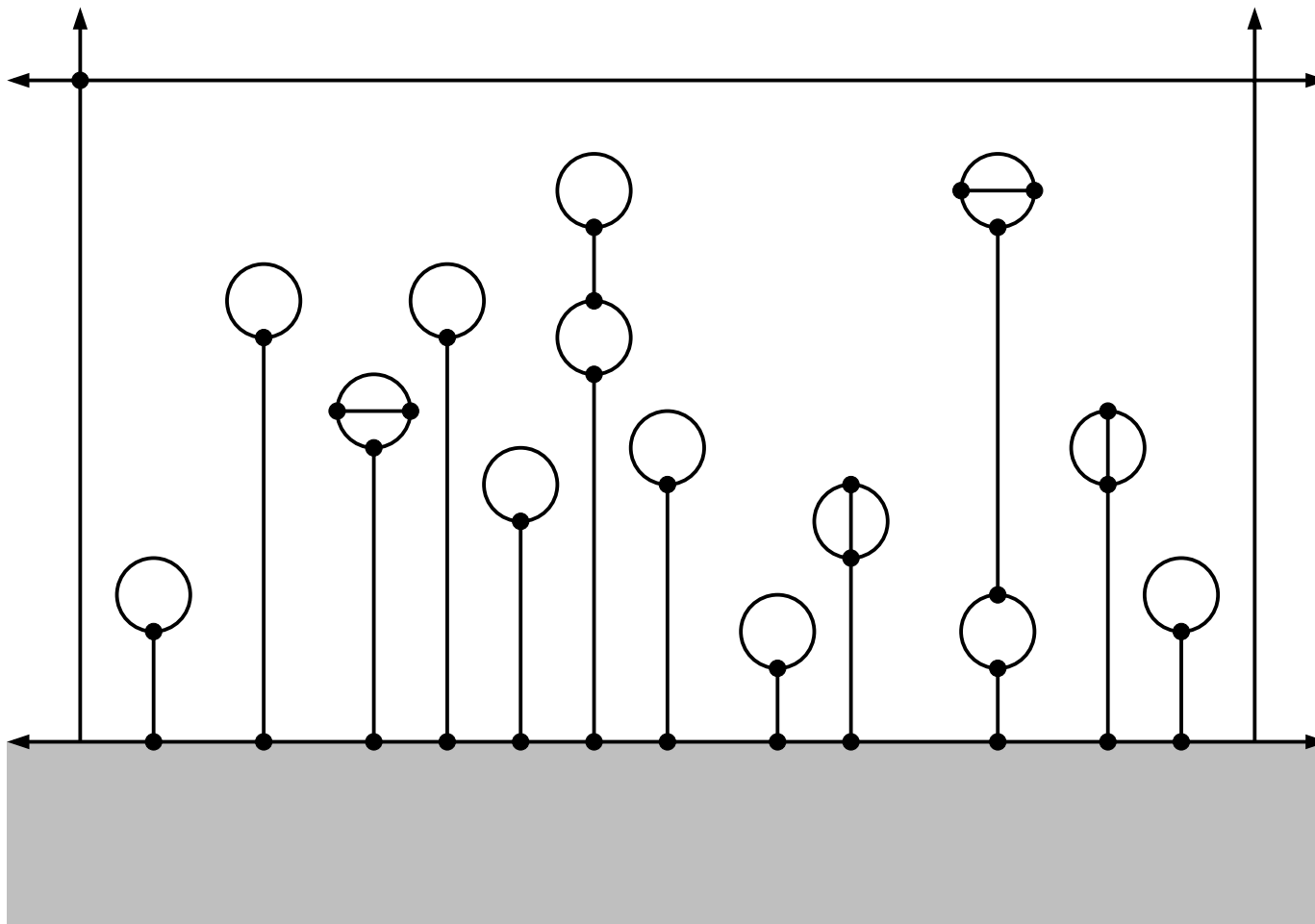
Hence

$$\text{ANS}(H) \leq 2E/F \leq 6 + 4e(H).$$

Definition: the face boundary: $\partial_F H$ is the set of faces in $G \setminus H$ that touch H .



Easy to see $|\partial_F H| \leq |\partial_E H|$. We need converse direction.



One adjacent face, many adjacent edges

Lemma 2: Suppose the faces of G have diameters $\leq D < \infty$ and H is a sub-map. Let N be the number of faces that lie inside a $3D$ -neighborhood of ∂H . Then $|\partial_E H| \leq 3N$.

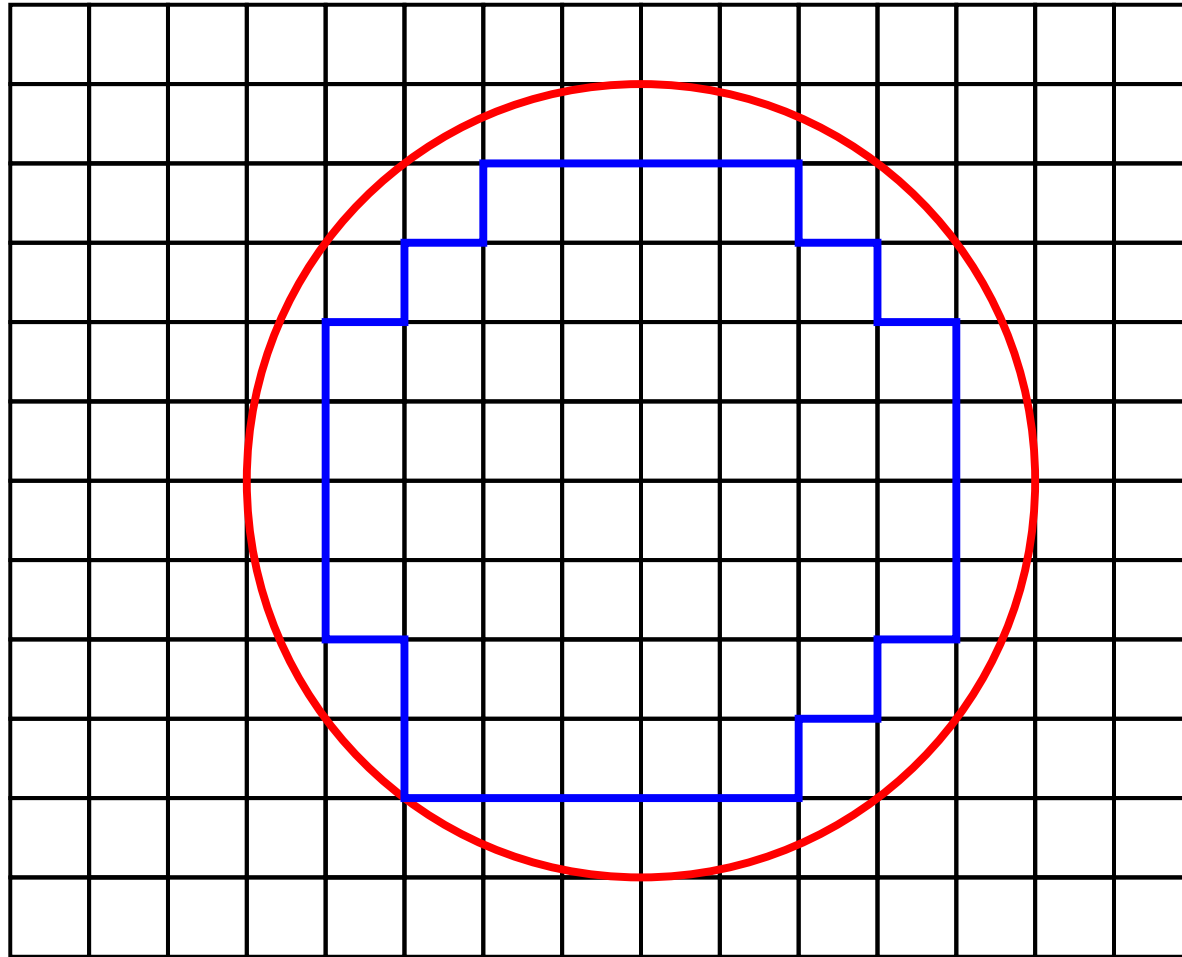
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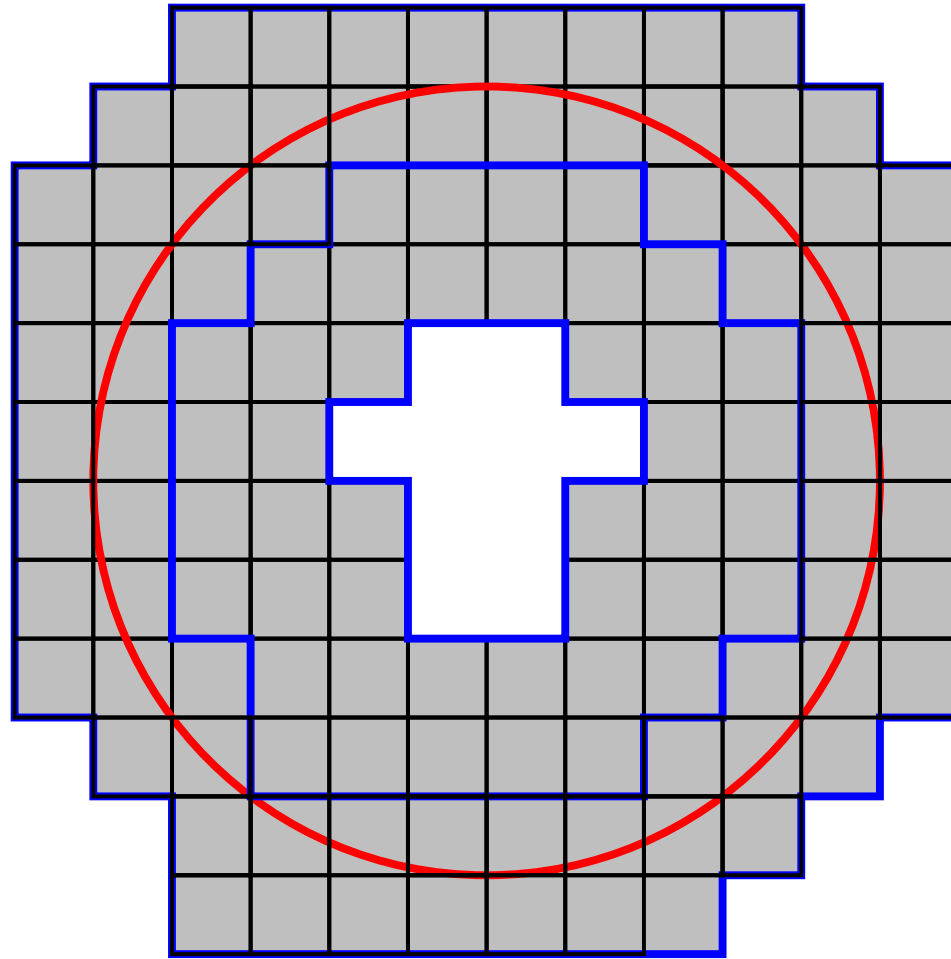
Proof:

Let G' be the finite graph on the sphere consisting of the faces of G that lie within the $2D$ -neighborhood of ∂H , together with their edges and vertices.

Check that edges of G' include all edges in $\partial_E H$.

Each face of G' is either a face of G or contains a face of G that is within $3D$ of ∂H . Thus number of faces of G' is at most N .





Lemma 2: Suppose the faces of G have diameters $\leq D < \infty$ and H is a sub-map. Let N be the number of faces that lie inside a $3D$ -neighborhood of ∂H . Then $|\partial_E H| \leq 3N$.

Proof continued:

Define G'' by removing any vertices of degree 2 from G' ; combine edges.

G' and G'' have same number of faces, $\leq N$.

If $e \in \partial_E H$, endpoints have degree ≥ 3 in G' (any edge touching e in G is in G'). So e is an edge of G'' .

By Sphere Theorem, edges in G'' bounded by three times faces in G'' .

Hence $|\partial_E H| \leq 3N$.

Theorem 1: Suppose all faces of G have diameter $\leq D < \infty$, and have area $\geq A > 0$, and that every vertex has degree ≥ 3 . Then

$$\limsup_{t \rightarrow \infty} \text{ANS}(H(R, t, x)) \leq 6.$$

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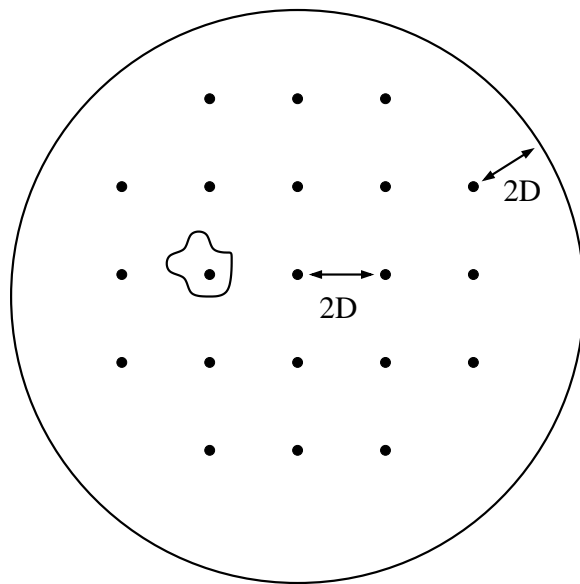
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Proof: Only need show $e(H(R, t, x)) \rightarrow 0$.

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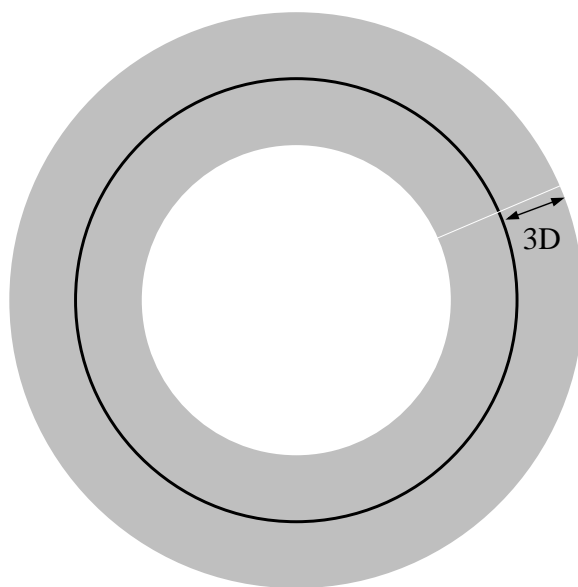
Proof continued: The number of faces in $H(R, t, x)$ is $\geq ct^2$.



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Proof continued: Area within $3D$ of ∂H is $O(t \cdot D)$,
so $|\partial_E H| = O(tD/A)$.



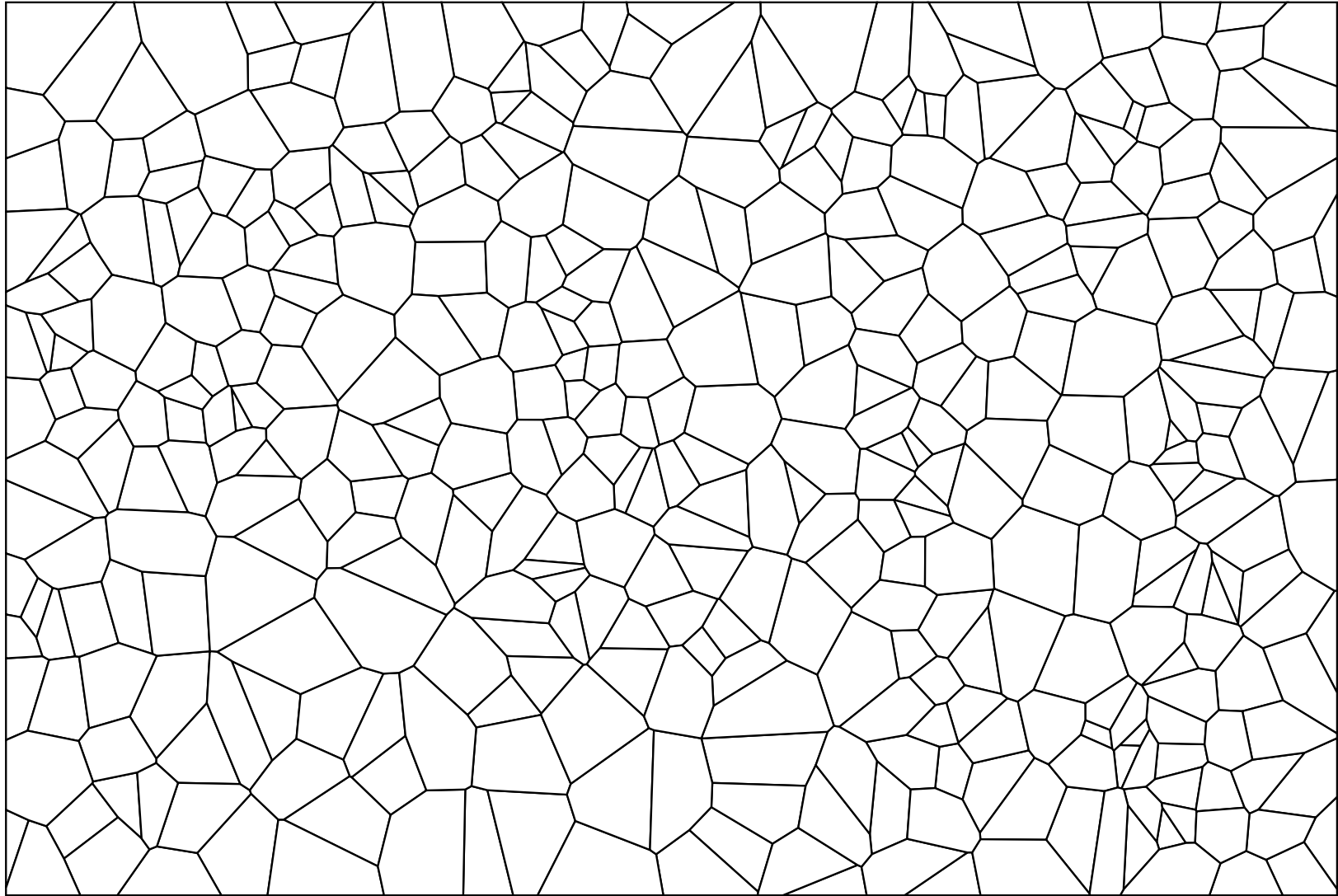
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Proof continued: Hence $e(H) = O(t/t^2) = O(1/t)$.

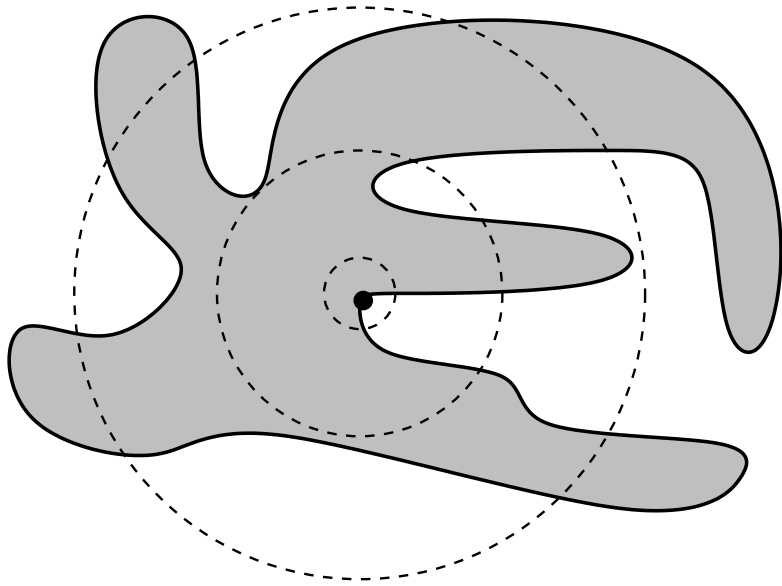
Theorem 2: If G satisfies the area lower bound and diameter upper bounded and every vertex has degree 3, then

$$\lim_{t \rightarrow 0} \text{ANS}(t) = 6.$$

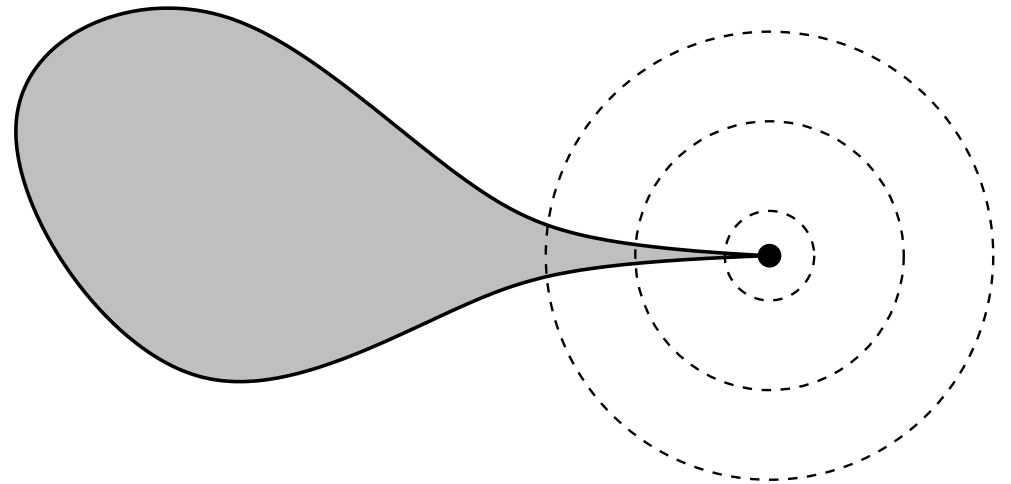


Voronoi diagram of Poisson point process

Ω is δ -thick if for all $0 < r < \text{diam}(\Omega)$ and all $x \in \partial\Omega$ we have

$$\text{area}(\Omega \cap D(x, r)) \geq \delta r^2.$$


Thick



Not Thick

True if Ω is convex, with a δ that depends on the aspect ratio.

Holds for quasidisks.

Theorem 3: Suppose G is a planar map and every face is δ -thick, for some $\delta > 0$. Then there is a nested, increasing sequence of sub-maps $\{H_n\}$ so that $\lim_{n \rightarrow \infty} \text{ANS}(H_n) \leq 6$.

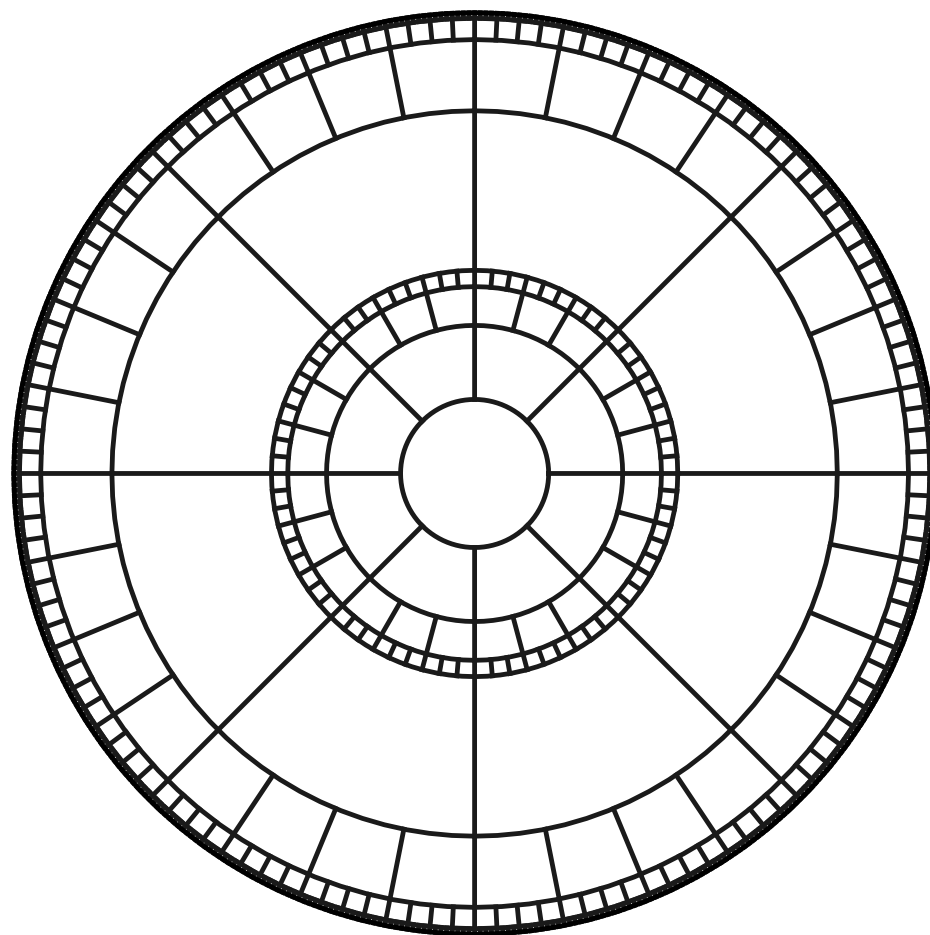
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Idea of proof:

If $\liminf > 6$, then $e(H(t)) > \epsilon > 0$ for all $t > t_0$.

We can show this forces $|\partial_F H(t)| \nearrow \infty$ in a finite time.

This contradicts local finiteness of G .



$$\limsup = \infty, \liminf \leq 6.$$

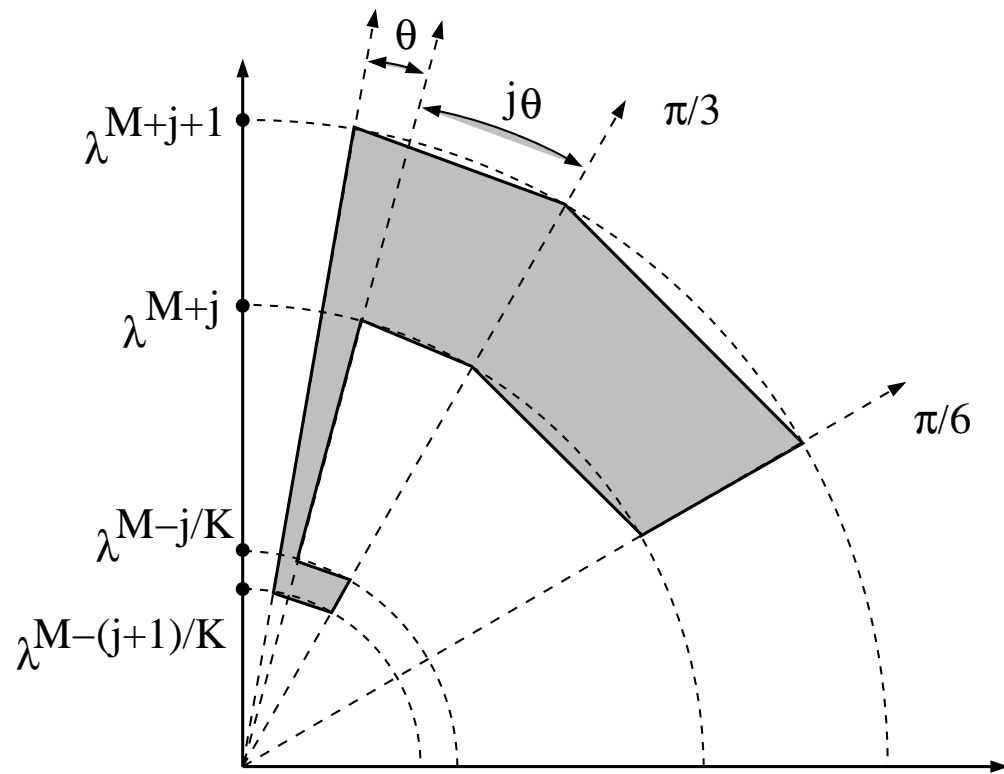
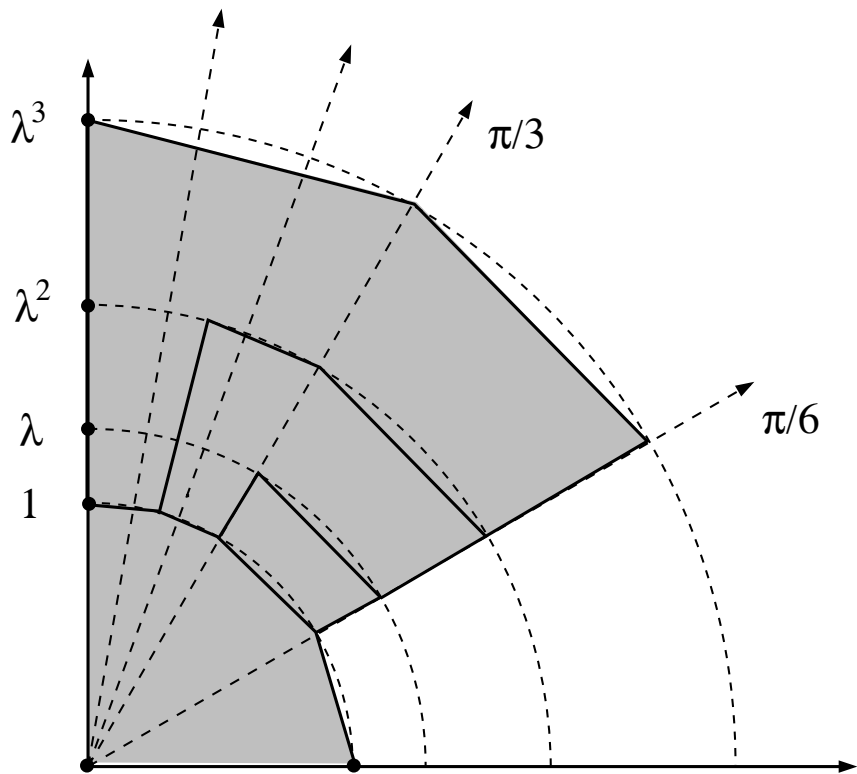
Aspect Ratio(K) = $\inf R/r$ where $D(x, r) \subset K \subset D(y, R)$.

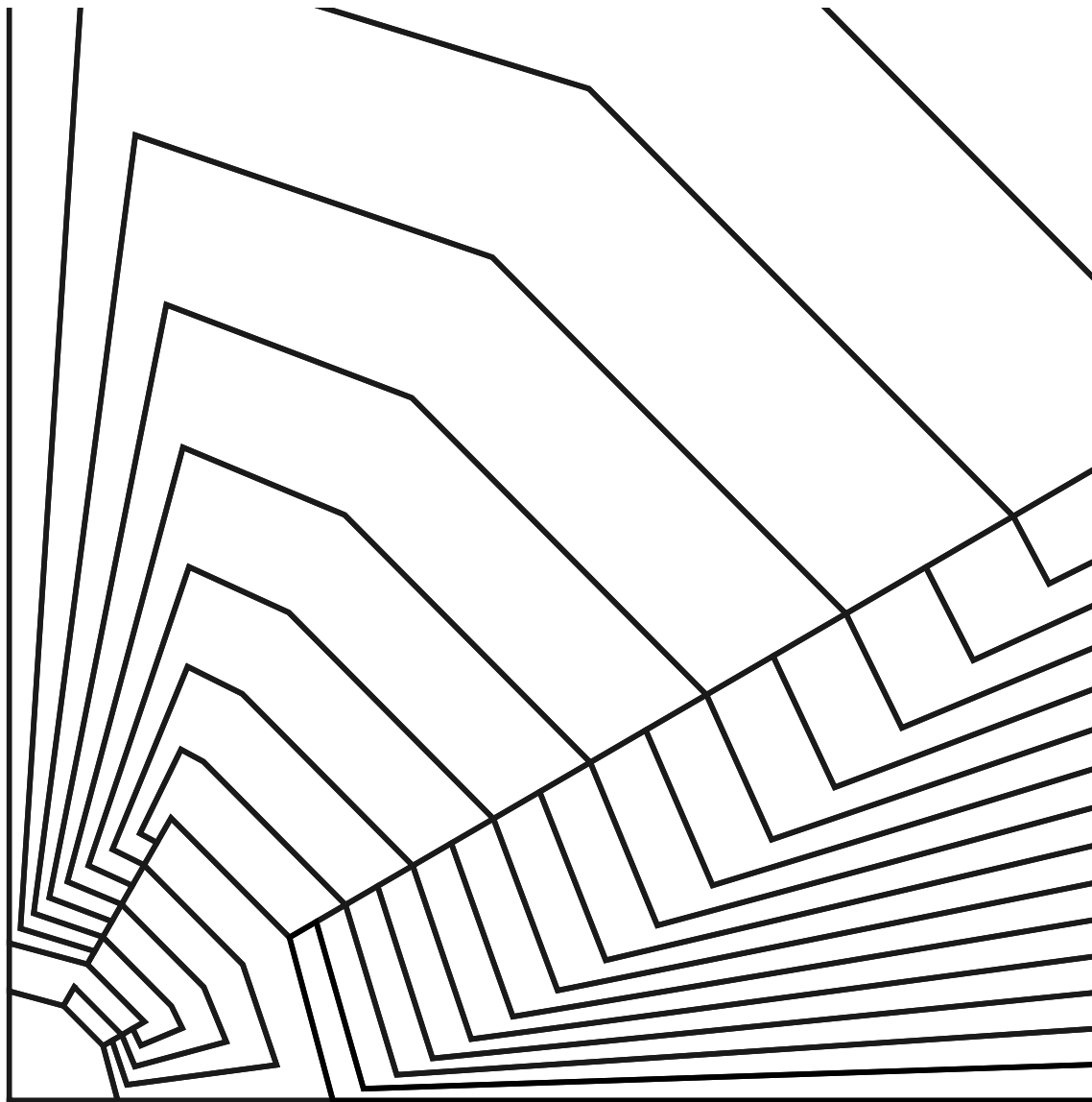
Is there a “6” theorem if faces have uniformly bounded aspect ratios?

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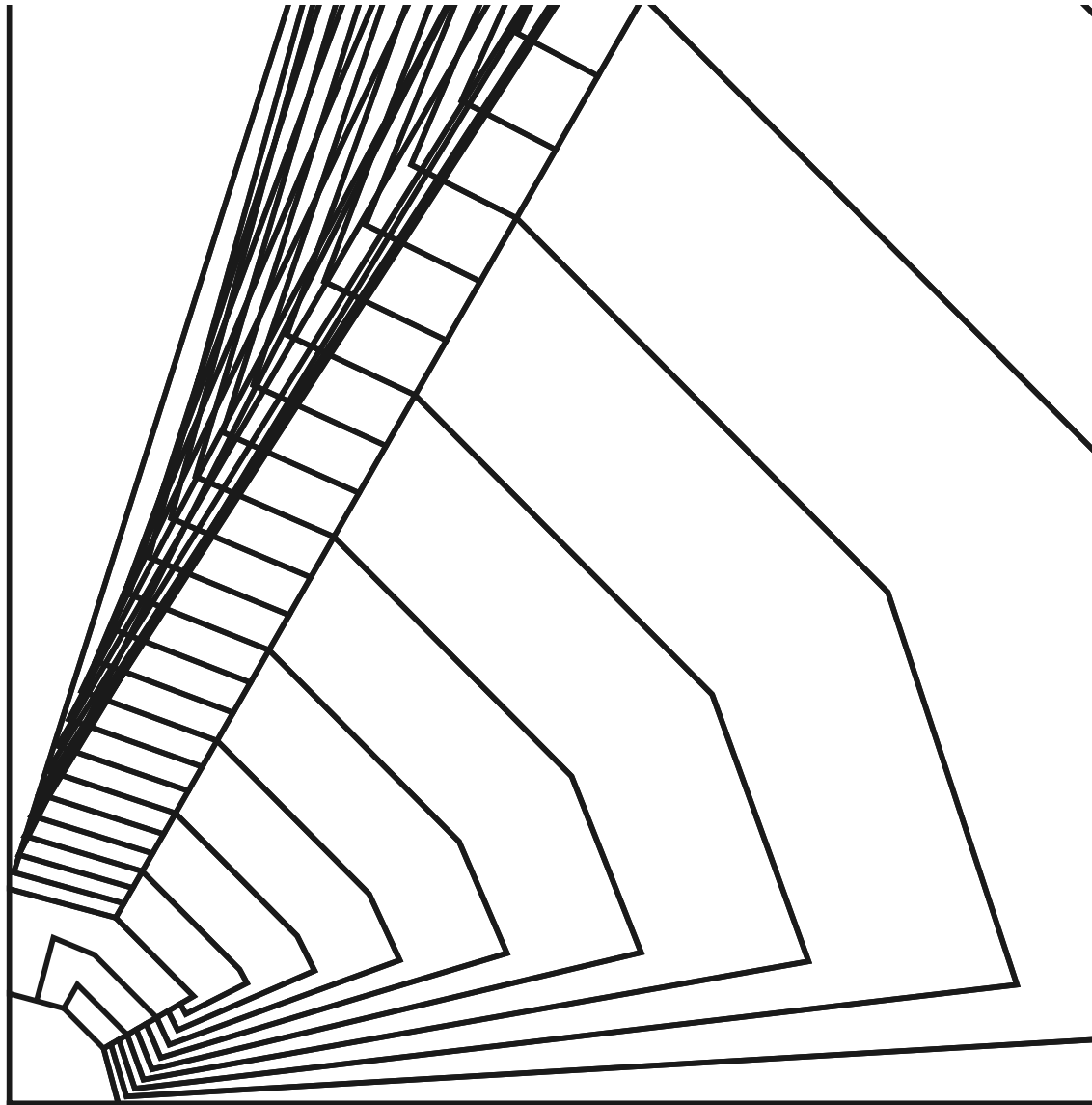
Is there a “6” theorem if faces have uniformly bounded aspect ratios?

No.





Out-radius/in-radius is bounded, but number of neighbors $\nearrow \infty$



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