

Enumerative Geometry: from Classical to Modern

Aleksey Zinger

Stony Brook University

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Summary

- Classical enumerative geometry: examples
- Modern tools: Gromov-Witten invariants
counts of holomorphic maps
- Insights from string theory:
 - quantum cohomology: refinement of usual cohomology
 - mirror symmetry formulas
duality between symplectic/holomorphic structures
 - integrality predictions for GW-invariants
geometric explanation yet to be discovered

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What is Classical EG about?

How many geometric objects satisfy given geometric conditions?

objects = curves, surfaces, ...

conditions = passing through given points, curves, ...

tangent to given curves, surfaces, ...

having given shape: genus, singularities, degree

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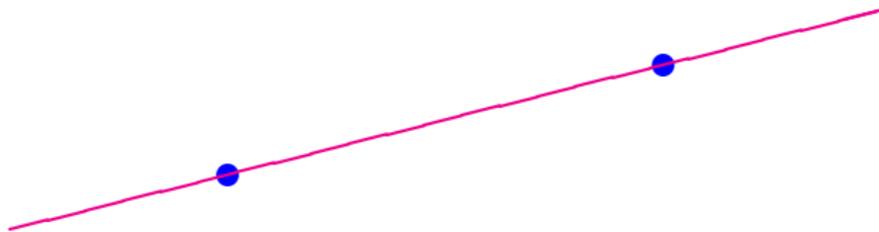
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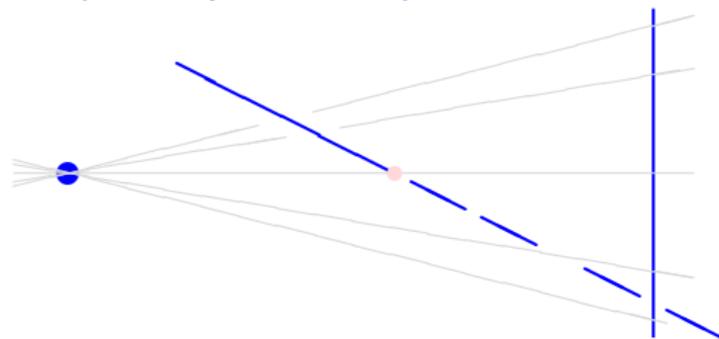
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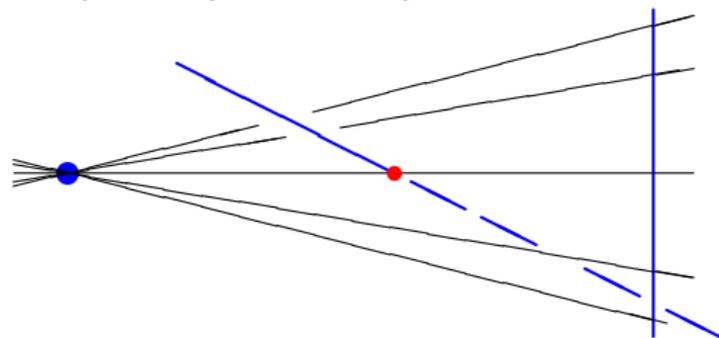
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lines thr. the **point** and **1st line** form a plane
 2nd line intersects the plane in **1 point**

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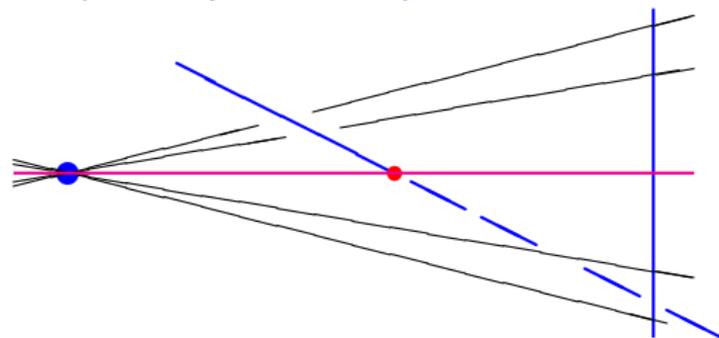
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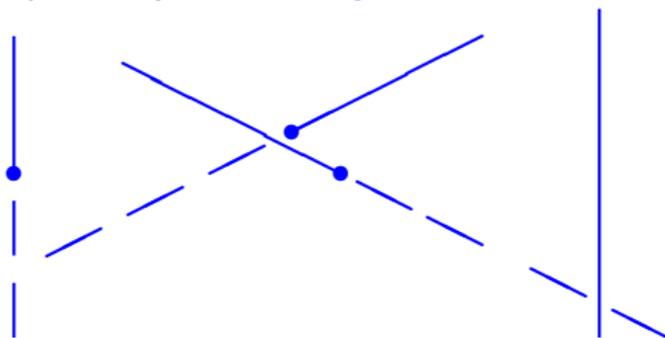
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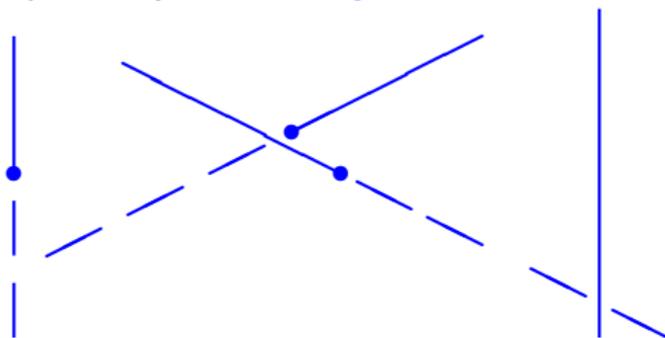
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bring two of the lines together so that they intersect in a point and form a plane

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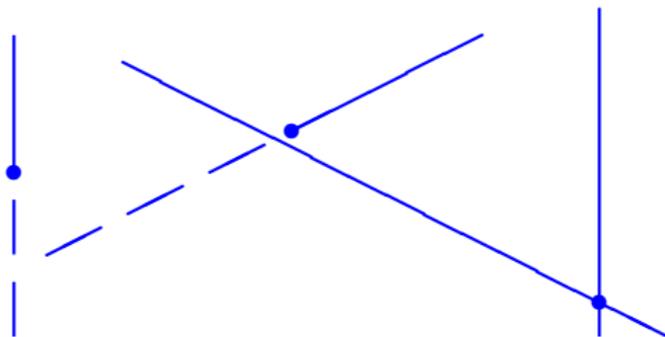
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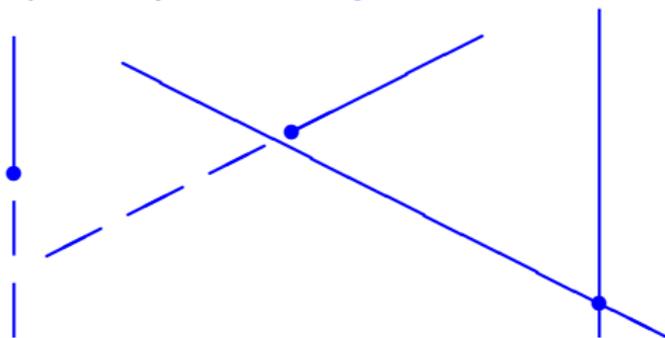
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1 **line** passes thr the **intersection pt** and lines #3,4
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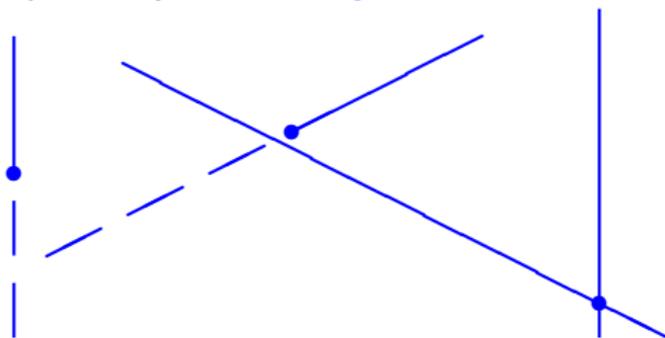


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General Line Counting Problems

All/most line counting problems in vector space V reduce to computing intersections of cycles on

$$\begin{aligned} G(2, V) &\equiv \{2 \text{ dim linear subspaces of } V\} \\ &\cong \{(\text{affine}) \text{ lines in } V\} \end{aligned}$$

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Example 0 (a semi-modern view)

Q: How many **lines** pass through **2 distinct points**?

A **line** in the plane is described by $(A, B, C) \neq 0$:

$$Ax + By + C = 0.$$

(A, B, C) and (A', B', C') describe the same line iff

$$(A', B', C') = \lambda(A, B, C)$$

$$\therefore \{\text{lines in } (x, y)\text{-plane}\} = \{1\text{-dim lin subs of } (A, B, C)\text{-space}\} \\ \equiv \mathbb{P}^2.$$

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degree d curve in (x, y) -plane

\equiv 0-set of nonzero degree d polynomial in (x, y)

polynomials Q and Q' determine same curve iff $Q' = \lambda Q$

coefficients of Q is $\binom{d+2}{2} \implies$

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"Passing thr a point" = 1 linear eqn on coefficients of Q

\implies get hyperplane in $\binom{d+2}{2}$ -dim v.s. of coefficients

$$\{\text{deg } d \text{ curves in } (x, y)\text{-plane thr. } (x_i, y_i)\} \approx \mathbb{P}^{N(d)-1} \subset \mathbb{P}^{N(d)}$$

intersection of $\binom{d+2}{2} - 1$ HPs in $\binom{d+2}{2}$ -dim v.s. is 1 1dim lin subs

intersection of $N(d)$ HPs in $\mathbb{P}^{N(d)}$ is 1 point

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Count **complex** curves = (singular) Riemann surfaces Σ

of fixed genus g , fixed degree d

in $\mathbb{C}^n, \mathbb{C}P^n = \mathbb{C}^n \sqcup \mathbb{C}^{n-1} \sqcup \dots \sqcup \mathbb{C}^0$

in a hypersurface $Y \subset \mathbb{C}^n, \mathbb{C}P^n$ (0-set of a polynomial)

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$n_d \equiv$ # genus 0 degree d **plane curves** thr. $(3d-1)$ **general pts**

$n_1 = 1$: # **lines** thr **2 pts**

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$n_3 = 12$: # **nodal cubics** thr **8 pts** $\implies \int_{\overline{\mathcal{M}}_{1,3}} \psi_1 = \frac{1}{24}$

$n_3 =$ # zeros of transverse bundle section over $\mathbb{C}P^1 \times \mathbb{C}P^2$
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Genus 0 Plane Quartics thr 11 pts

$n_4 = \#$ plane quartics thr 11 pts with 3 non-separating nodes

Zeuthen'1870s: $n_4 = 620 = 675 - 55$

$3! \cdot 675 =$ euler class of rank 9 vector bundle over $\mathbb{C}P^3 \times (\mathbb{C}P^2)^3$
 minus excess contributions of a certain section

$\mathbb{C}P^3 =$ quartics thr 11 pts; $\mathbb{C}P^2 =$ possibilities for i -th node

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Kontsevich's Formula (Ruan-Tian'1993)

$n_d \equiv$ # genus 0 degree d **plane curves** thr. $(3d-1)$ general pts

$$n_1 = 1$$

$$n_d = \frac{1}{6(d-1)} \sum_{d_1+d_2=d} \left(d_1 d_2 - 2 \frac{(d_1 - d_2)^2}{3d-2} \right) \binom{3d-2}{3d_1-1} d_1 d_2 n_{d_1} n_{d_2}$$

$$n_2 = 1, n_3 = 12, n_4 = 620, n_5 = 87,304, n_6 = 26,312,976, \dots$$

Kontsevich's Formula (Ruan-Tian'1993)

$n_d \equiv$ # genus 0 degree d plane curves thr. $(3d-1)$ general pts

$$n_1 = 1$$

$$n_d = \frac{1}{6(d-1)} \sum_{d_1+d_2=d} \left(d_1 d_2 - 2 \frac{(d_1 - d_2)^2}{3d-2} \right) \binom{3d-2}{3d_1-1} d_1 d_2 n_{d_1} n_{d_2}$$

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Gromov's 1985 paper

Consider equivalence classes of **maps** $f: (\Sigma, j) \longrightarrow \mathbb{C}P^n$
 (Σ, j) =connected Riemann surface, possibly with nodes

$f: (\Sigma, j) \longrightarrow \mathbb{C}P^n$ and $f': (\Sigma', j') \longrightarrow \mathbb{C}P^n$ are **equivalent** if
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Gromov's Compactness Theorem

genus of $f : (\Sigma, j) \longrightarrow \mathbb{C}P^n$ is # of holes in Σ ($\geq g(\Sigma)$)

degree d of $f \equiv |f^{-1}(H)|$ for a generic hyperplane:

$$f_*[\Sigma] = d[\mathbb{C}P^1] \in H_2(\mathbb{C}P^n; \mathbb{Z}) = \mathbb{Z}[\mathbb{C}P^1]$$

Theorem: With respect to a natural topology,

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Maps vs. Curves

Image of holomorphic $f: (\Sigma, j) \rightarrow \mathbb{C}P^n$ is a curve
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$$\begin{aligned} \implies n_d &\equiv \# \text{ genus 0 degree } d \text{ curves thr. } (3d-1) \text{ pts in } \mathbb{C}P^2 \\ &= \# \text{ degree } d \text{ } f: (S^2, j) \rightarrow \mathbb{C}P^2 \text{ s.t. } p_i \in f(\mathbb{C}P^1) \\ &\quad i = 1, \dots, 3d-1 \\ &= \# \{ [f: (\Sigma, j) \rightarrow \mathbb{C}P^2] \in \overline{\mathfrak{M}}_{10}(\mathbb{C}P^2, d) : p_i \in f(\Sigma) \} \end{aligned}$$

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Use counts of genus 0 maps to $\mathbb{C}P^n$ to deform \cup -product on H^* ,

$$H^*(\mathbb{C}P^n) = \mathbb{Z}[x]/x^{n+1}, \quad x^a \cup x^b = x^{a+b},$$

to $*$ -product on $H^*(\mathbb{C}P^n)[q_0, \dots, q_n]$

$x^a * x^b = x^{a+b} + q$ -corrections counting genus 0 maps
thr. $\mathbb{C}P^{n-a}, \mathbb{C}P^{n-b}$

Theorem (McDuff-Salamon'93, Ruan-Tian'93, ...)

The product $*$ is associative

$*$ generalizes to all cmpt algebraic/symplectic manifolds

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Remark: Classical proof of Kontsevich's formula for $\mathbb{C}P^2$ only:
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Other Enumerative Applications of Stable Maps

- Genus 0 with singularities: Pandharipande, Vakil, Z.-
- Genus 1: R. Pandharipande, Ionel, Z.-
- Genus 2,3: Z.-

Gromov-Witten Invariants

$Y = \mathbb{C}P^n$, = hypersurface in $\mathbb{C}P^n$ (0-set of a polynomial),...
 $\mu_1, \dots, \mu_k \subset Y$ cycles

$$\text{GW}_{g,d}^Y(\mu) \equiv \# \{ [f: (\Sigma, j) \rightarrow Y] \in \overline{\mathfrak{M}}_g(Y, d) : f(\Sigma) \cap \mu_i \neq \emptyset \}$$

$g = 0$, $Y = \mathbb{C}P^n$: $\overline{\mathfrak{M}}_g(Y, d)$ is smooth, of expected dim, "#"=#

Typically, $\overline{\mathfrak{M}}_g(Y, d)$ is highly singular, of wrong dim

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Example: Quintic Threefold

$Y_5 \subset \mathbb{C}P^4$ 0-set of a degree 5 polynomial Q

Schubert Calculus: Y_5 contains 2,875 (isolated) lines

S. Katz'86 (via Schubert): Y_5 contains 609,250 conics

For each line $L \subset Y_5$ and conic $C \subset Y_5$,

$$\{[f: (\Sigma, j) \longrightarrow Y_5] \in \overline{\mathfrak{M}}_0(Y_5, 2) : f(\Sigma) \subset L\} \approx \overline{\mathfrak{M}}_0(\mathbb{C}P^1, 2)$$

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are connected components of $\overline{\mathfrak{M}}_0(Y_5, 2)$ of dimensions 2 and 0

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holomorphic degree d $f: \mathbb{C}P^1 \longrightarrow \mathbb{C}P^4$ has the form

$$f([u, v]) = [R_1(u, v), \dots, R_5(u, v)]$$

$R_1, \dots, R_5 =$ homogeneous polynomials of degree d

$$\implies \dim \overline{\mathfrak{M}}_0(\mathbb{C}P^4, d) = 5 \cdot (d+1) - 1 - 3$$

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$$\implies \dim^{vir} \overline{\mathfrak{M}}_0(Y_5, d) = \dim \overline{\mathfrak{M}}_0(\mathbb{C}P^4, d) - (5d + 1) = 0$$

A more elaborate computation gives

$$\dim^{vir} \overline{\mathfrak{M}}_g(Y_5, d) = 0 \quad \forall g$$

\implies want to define

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ν = small generic deformation of $\bar{\partial}$ -equation

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$$\bar{\partial}_j f \equiv df + J_{Y_5} \circ df \circ j$$

$$N_{g,d} \equiv |\overline{\mathfrak{M}}_g(Y_5, d)|^{vir}$$

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Y_5 is Calabi-Yau 3-fold:

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A-Model partition function
for Calabi-Yau 3-fold Y

MIRROR
principle

B-Model partition function
for mirror (family) of Y

generating function for GWs of Y :

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Mirror Symmetry Predictions and Verifications

Predictions

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Theorem (Givental'96, Lian-Liu-Yau'97, ... 2000)

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Need to compute each $N_{g,d}$ and all of them (for fixed g):

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From $Y_5 \subset \mathbb{C}P^4$ to $\mathbb{C}P^4$

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 $\overline{\mathfrak{M}}_g(\mathcal{L}, d)$

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Genus 0 vs. Positive Genus

$g = 0$ everything is as expected:

- $\overline{\mathfrak{M}}_g(\mathbb{C}P^4, d)$ is smooth
- $[\overline{\mathfrak{M}}_g(\mathbb{C}P^4, d)]^{vir} = [\overline{\mathfrak{M}}_g(\mathbb{C}P^4, d)]$
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Thm. A (J. Li–Z.'04): HP holds for **reduced** genus 1 GWs

$$|\overline{\mathfrak{M}}_1^0(Y_5, d)|^{vir} = e(\mathcal{V}_{1,d}) \cap \overline{\mathfrak{M}}_1^0(\mathbb{C}P^4, d).$$

This generalizes to complete intersections $Y \subset \mathbb{C}P^n$.

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- $(\mathbb{C}^*)^5$ acts on $\mathbb{C}P^4$ (with 5 fixed pts)
- \implies on $\overline{\mathfrak{M}}_g(\mathbb{C}P^4, d)$ (with simple fixed loci)
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- $\int_{\overline{\mathfrak{M}}_g^0(\mathbb{C}P^4, d)} e(\mathcal{V}_{g,d})$ localizes to fixed loci
 - $g = 0$: Atiyah-Bott Localization Thm reduces \int to \sum_{graphs}
 - $g = 1$: $\overline{\mathfrak{M}}_g^0(\mathbb{C}P^4, d), \mathcal{V}_{g,d}$ singular \implies AB does not apply

Genus 1 Bypass

Thm. C (Vakil–Z.'05): $\mathcal{V}_{1,d} \rightarrow \overline{\mathfrak{M}}_1^0(\mathbb{C}P^4, d)$ admit
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 \end{array}$$

$$\Rightarrow \int_{\overline{\mathfrak{M}}_1^0(\mathbb{C}P^4, d)} e(\mathcal{V}_{1,d}) = \int_{\widetilde{\overline{\mathfrak{M}}}_1^0(\mathbb{C}P^4, d)} e(\tilde{\mathcal{V}}_{1,d})$$

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Computation of Genus 1 GWs of CIs

Thm. C generalizes to all $\mathcal{V}_{1,d} \longrightarrow \overline{\mathfrak{M}}_{1,k}^0(\mathbb{C}P^n, d)$:

$$\begin{array}{ccc}
 \mathcal{L} \equiv \mathcal{O}(a) & & \overline{\mathfrak{M}}_{1,k}(\mathcal{L}, d) \\
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