

Math Camp

Public Key Cryptography

Merkle-Hellman Knapsack Algorithm

Based on the subset-sum problem.

Example:

$A=(2, 3, 5, 10), \quad c = 18$

$A= (-7, -3, -2, 9000, 5, 8)$ and $c = 0$

$A = (5, 43, 17, 9, 21, 12)$ and $c = 47$

$A = (295, 592, 301, 14, 28, 353, 120, 236)$ and $c = 1129$

Euclidean Algorithm

1. Find the greatest common divisor of 72, 23

2. Find the greatest common divisor of 156, and 61

Definitions

Superincreasing sequence = each element is larger than the sum of all previous elements.

Relatively prime numbers = two numbers with no common factors

X modulo C = divide X by C and take the remainder

Merkle-Hellman Knapsack Algorithm.

Set up:

1. Choose your private key: pick a **superincreasing** set A of **7** numbers.
2. Choose a random number C that is: **larger than the sum of elements of A.**
3. Choose a random number R that is **relatively prime** to C.
4. Construct the public key B: take **$A \times R \text{ modulo } C$.**
5. Publish the public key B.

How to Encrypt each letter:

1. Convert letter to binary.
2. Suppose $B = \{b_1, b_2, b_3, b_4, b_5, b_6, b_7\}$, and your letter in binary is $n = z_1z_2z_3z_4z_5z_6z_7$.
3. Take the sum **$(z_1 \times b_1) + (z_2 \times b_2) + \dots + (z_7 \times b_7) = m$**
4. Send m.

Letter	Binary
A	01000001
B	01000010
C	01000011
D	01000100
E	01000101
F	01000110
G	01000111
H	01001000
I	01001001
J	01001010
K	01001011
L	01001100
M	01001101
N	01001110
O	01001111
P	01010000
Q	01010001
R	01010010
S	01010011
T	01010100
U	01010101
V	01010110
W	01010111
X	01011000
Y	01011001
Z	01011010

An Example:

Setting Up Your Private Key

Message to send: MATH

In Binary: 1001101 / 1000001 / 1010100 / 1001000

1. Choose a super increasing sequence of length 7

A =

2. Pick a number bigger than the sum of elements in A

C =

3. Pick a number relatively prime to C

R =

4. Find the remainder when you divide C by R.

5. Write $xR+yC=1$ for some numbers x and y

Record:

A =

C =

R =

X =

6. Calculate your public key $B = A * R$, then remainder dividing by C .

Now we can encrypt!

Decryption

We've been sent the cipher

$m =$ _____

We have our private key $A =$

the number $C =$ _____

and the number $R =$ _____

used to construct the public key B .

1. We have our secret number $x =$

The number x modulo C is called the inverse of r modulo C .

2. Calculate the number $m' = mx$ modulo C .

3. Solve subset sum problem for A, m' .

4. Decrypt the message!!!