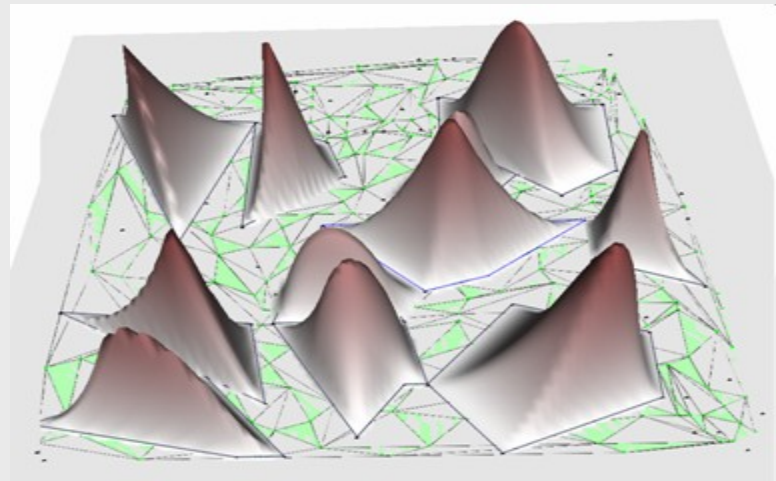
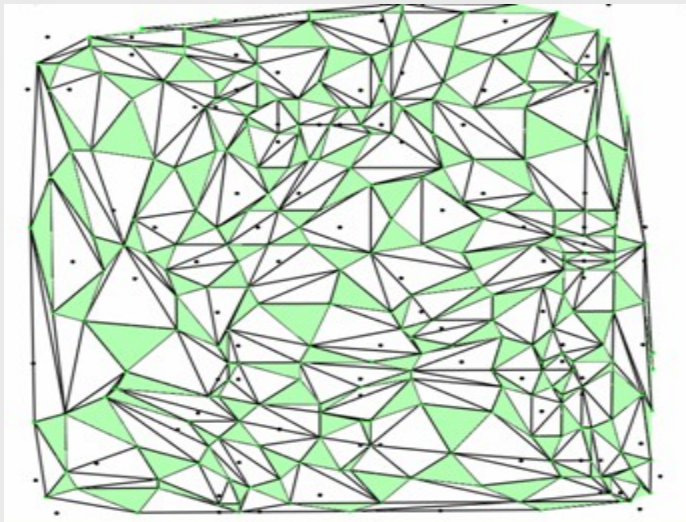


# Bivariate B-Splines From Centroid Triangulations



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UNC Chapel Hill

# Q Motivating Questions

Comput'l Geometry:

"PL surface meshes can be constructed from (irregular) points by triangulating.

*What about smooth surfaces?"*

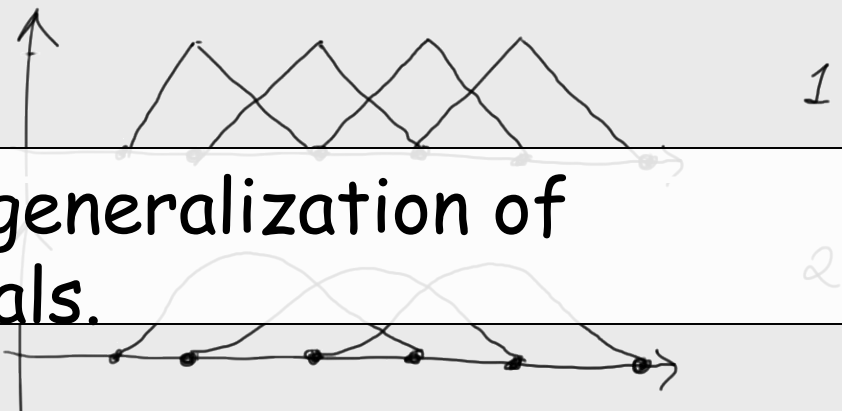
CAGD:

"Smooth B-splines can be constructed over (irregular) points along the real line.

*How do we make bivariate B-splines?"*

**A?**

centroid triangulations, a generalization of higher order Voronoi duals.



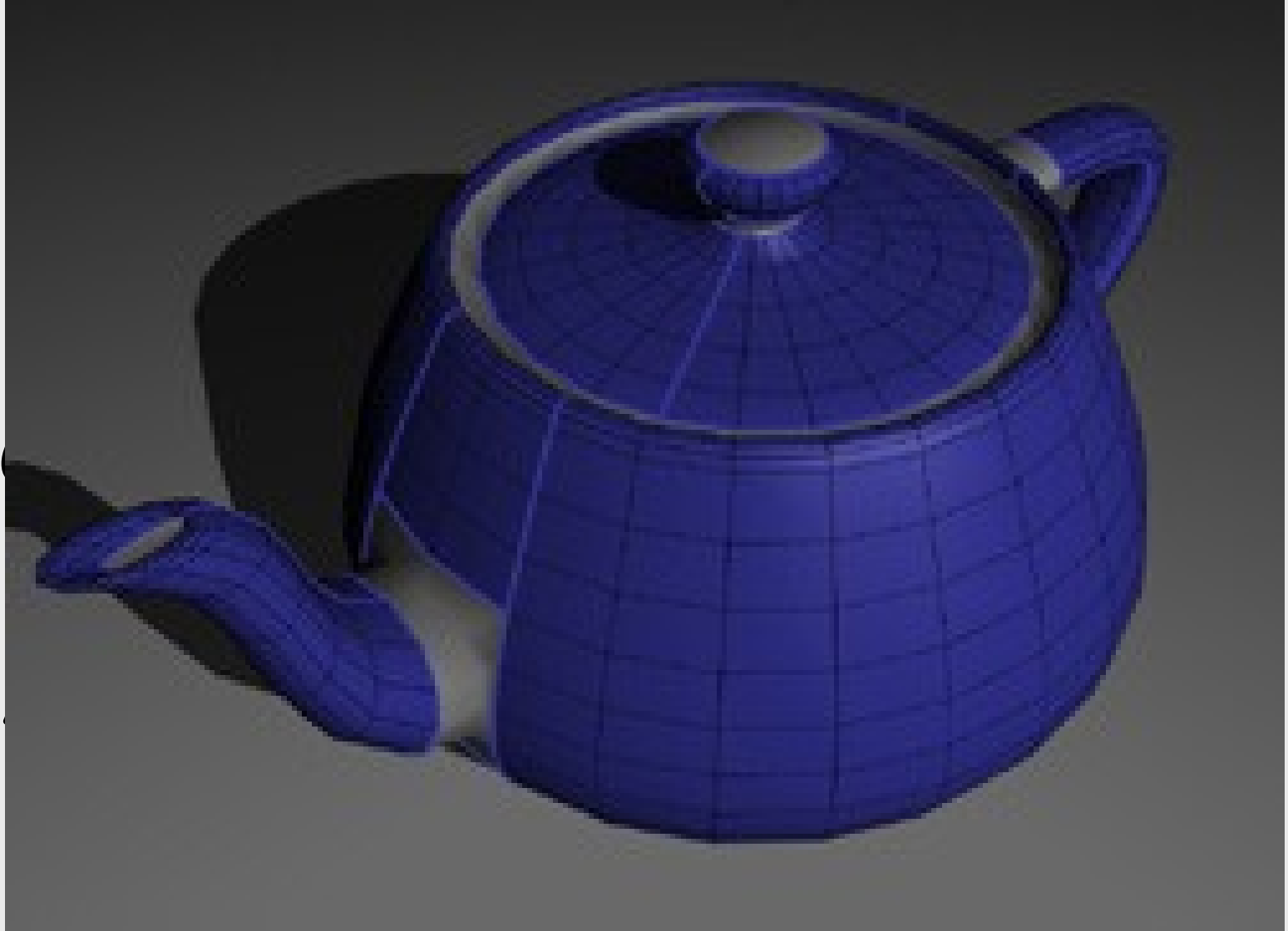
# Outline

- Context & Motivation

- 

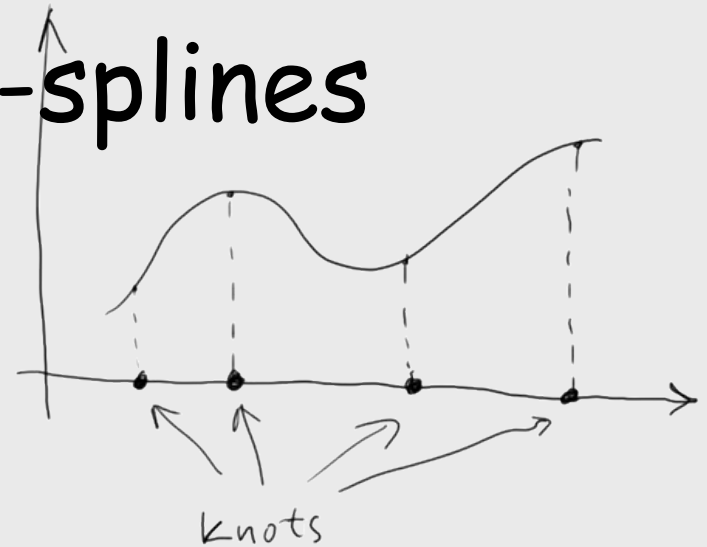
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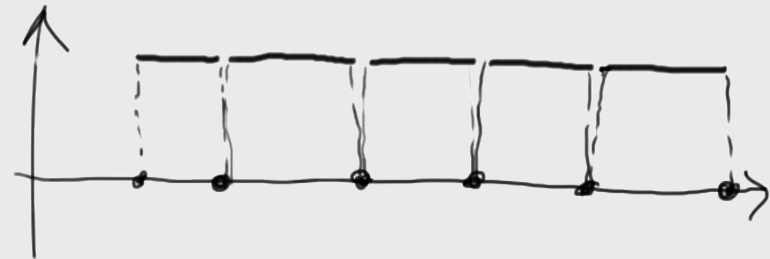


# Univariate B-splines

- *splines* :  
piecewise polynomials



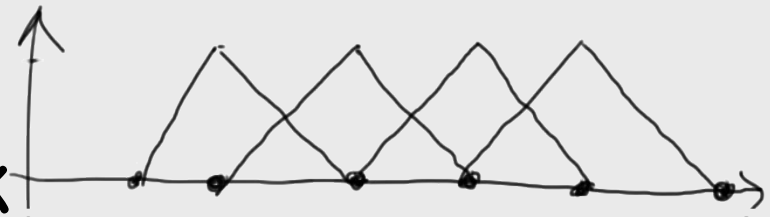
- *B-spline space*:  
linear combination  
of basis functions



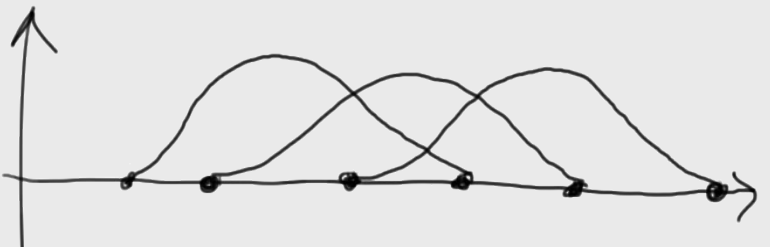
$k$

$\emptyset$

- A *B-spline* of deg.  $k$   
is defined for  
any  $k+2$  knots.



1



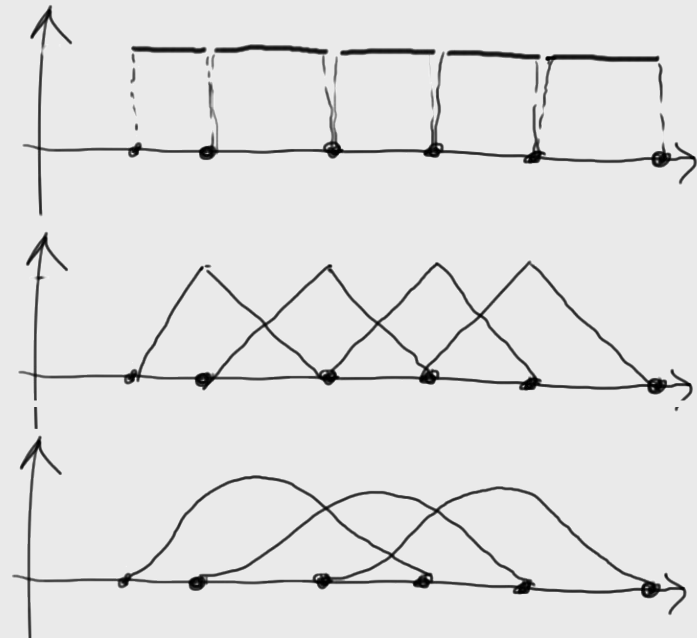
2

# Univariate B-splines

## Properties

- *local support*
- *optimal smoothness*
- *partition of unity*  
 $\sum B_i = 1$
- *polynomial reproduction*,  
 for any deg.  $k$  polynomial  $p$ ,  
 with polar form  $P$ ,  

$$p = \sum P(S_{i+1} \dots S_{i+k}) B_i(\cdot | S_i \dots S_{i+k+1})$$



$k$

$\emptyset$

1

2

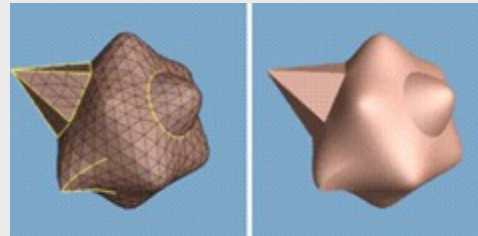
# What are multivariate splines?

Are they B-splines?

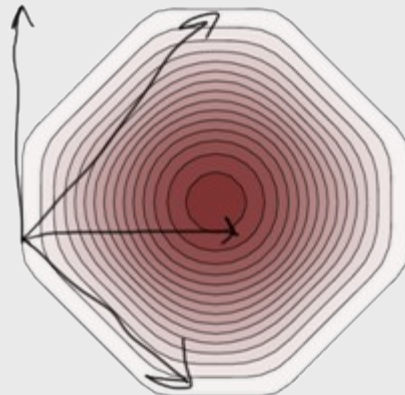
- tensor product



- subdivision



- box splines



# What are multivariate B-splines?

(multivariate) B-splines should define basis functions with no restriction on knot positions and have these properties of the classic B-splines:

- local support
- optimal smoothness
- partition of unity  $\sum B_i = 1$
- polynomial reproduction:  
for any degree  $k$  polynomial  $p$ ,  
with polar form  $P$ ,  
$$p = \sum P(S_{i+1} \dots S_{i+k}) B_i(\cdot | S_i \dots S_{i+k+1})$$

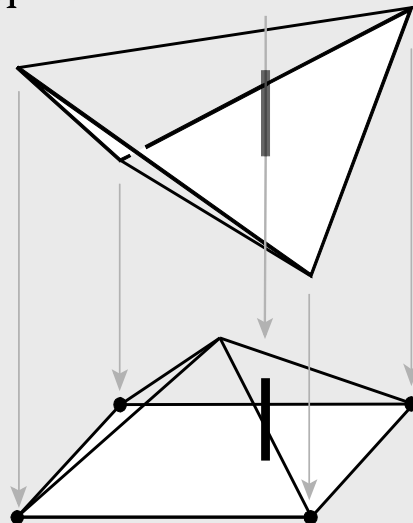
# Simplex spline [dB76]

A degree  $k$  polynomial defined on a set  $X$  of  $k+s+1$  points in  $\mathbb{R}^s$ .

Lift  $X$  to  $Y \in \mathbb{R}^{k+s}$  and take relative measure of the projection of this simplex:  
 $M(x | X) := \underline{\text{vol} \{ y \mid y \in [Y] \text{ and projects to } x \}}$

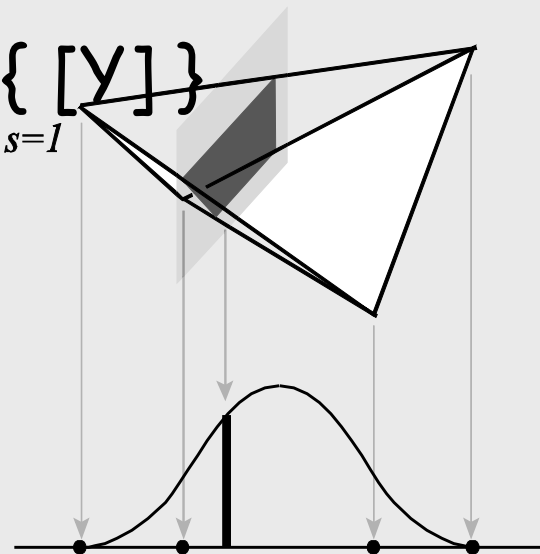
simplex spline

$k=1, s=2$



$\text{vol} \{ [Y] \}$

$k=2, s=1$

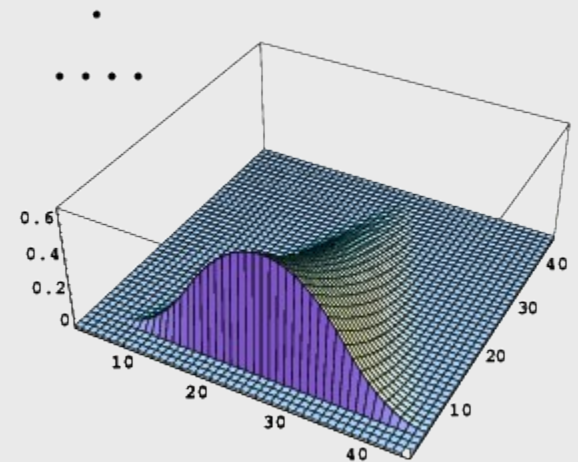
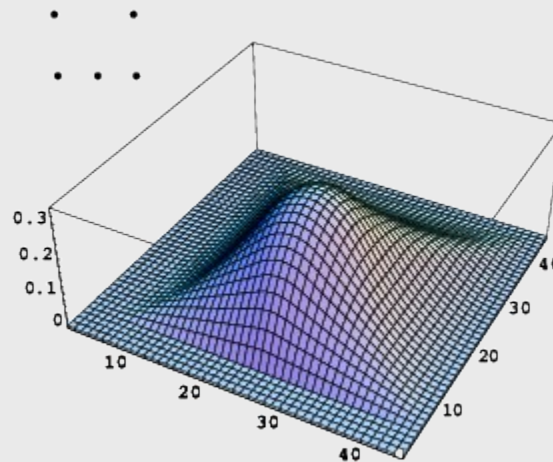
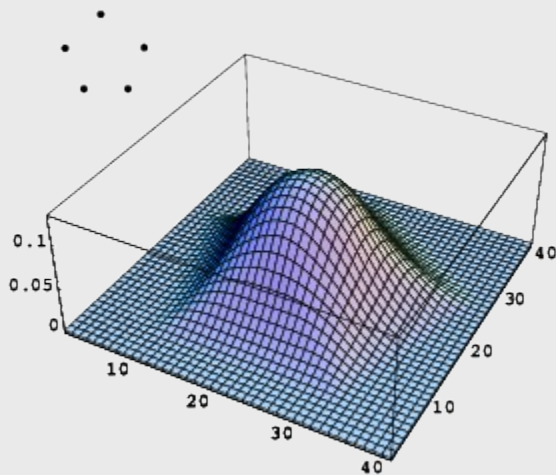




# Simplex spline

## Properties

- ✓ local support:  $M(\cdot|X)$  is non-zero only over the convex hull of  $X$ .
- ✓ optimally smooth, assuming  $X$  is in general position.



# What are multivariate B-splines?

Using simplex splines as basis:

The task of building multivariate B-splines becomes choosing the “right” configurations.

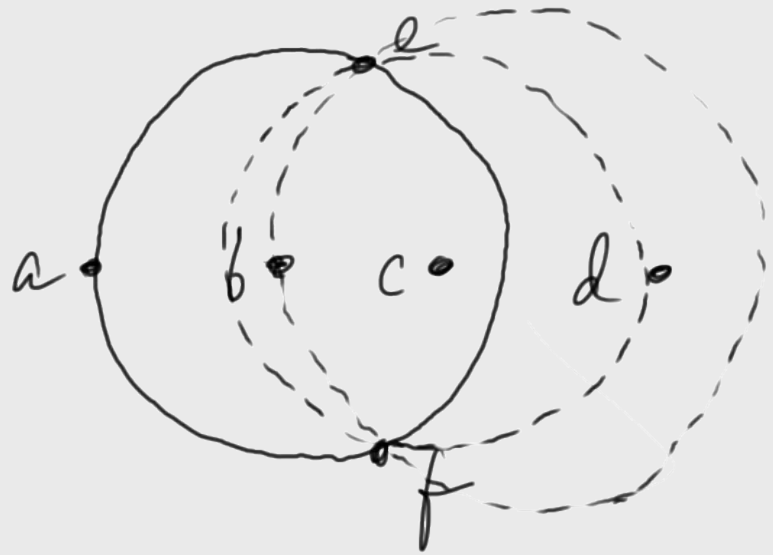
- All  $k$ -tuple configs [Dahmen & Micchelli 83]
- DMS-splines [Dahmen, Micchelli & Seidel 92]
- Delaunay configurations [Neamtu 01]

# Neamtu's Delaunay configurations

A degree  $k$  Delaunay configuration  $(t, I)$  is defined by a circle through  $t$  containing  $I$  inside.

$\Gamma_{\text{Del}}^2$

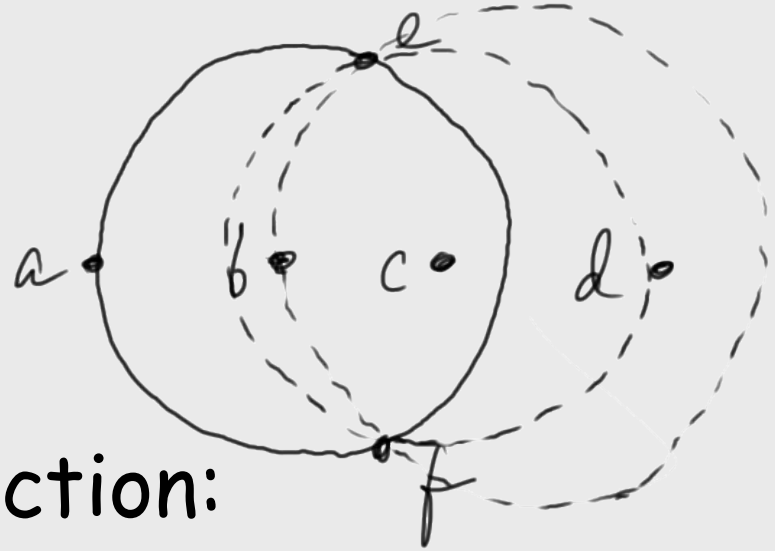
$= \{ (aef, bc),$   
 $(def, bc),$   
 $(bef, cd) \dots \}$



# Delaunay configurations

Spline space for  $\Gamma_{\text{Del}}^k$ :

$$\text{span}\{ d(t) M(\cdot | t \cup I) \}_{(t,I) \in \Gamma_{\text{Del}}^k}$$



## Properties

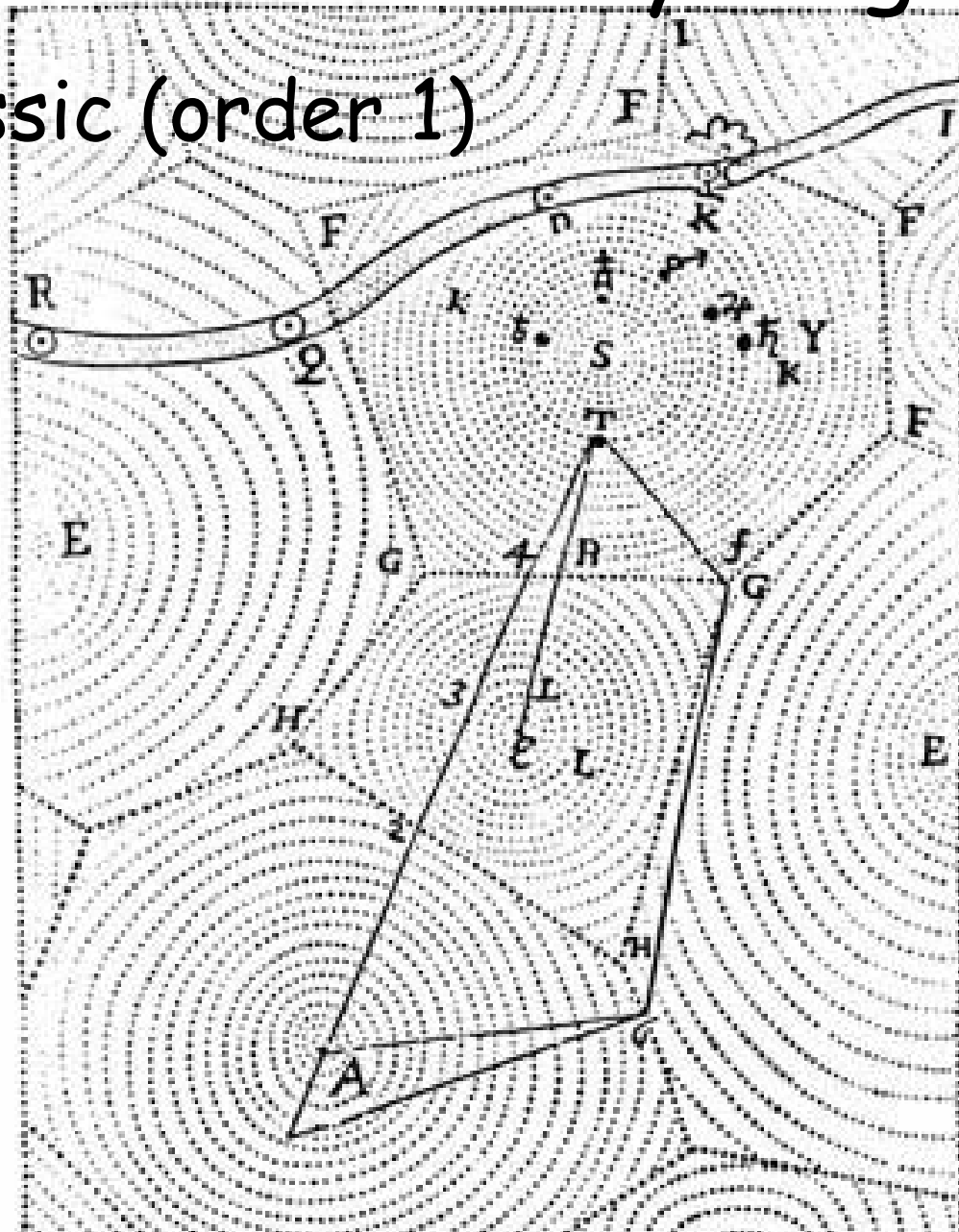
✓ polynomial reproduction:

for any deg.  $k$  polynomial  $p$ ,

$$p = \sum_{(t,I) \in \Gamma_{\text{Del}}^k} P(I) d(t) M(\cdot | t \cup I)$$

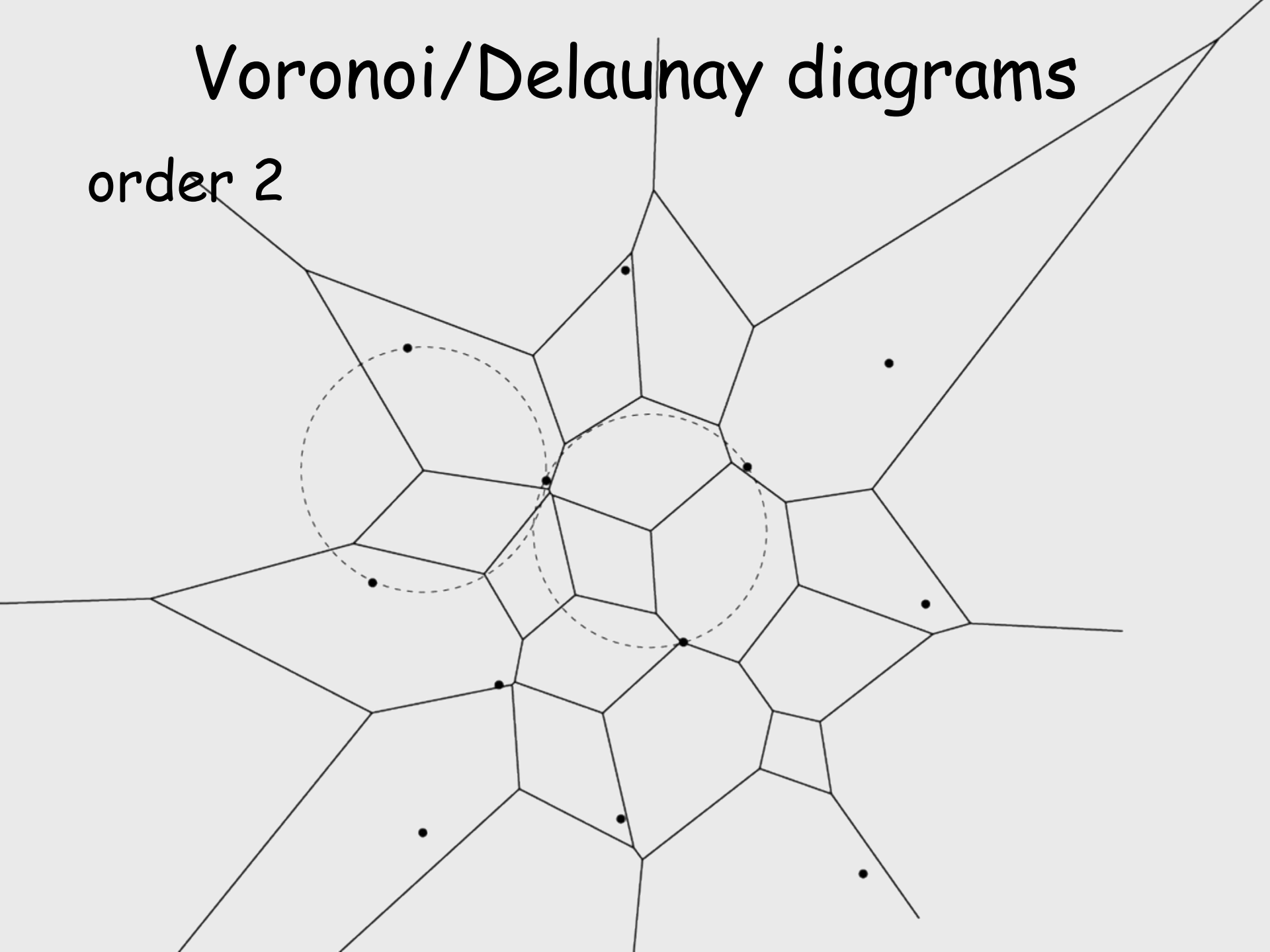
# Voronoi/Delaunay diagrams

the classic (order 1)



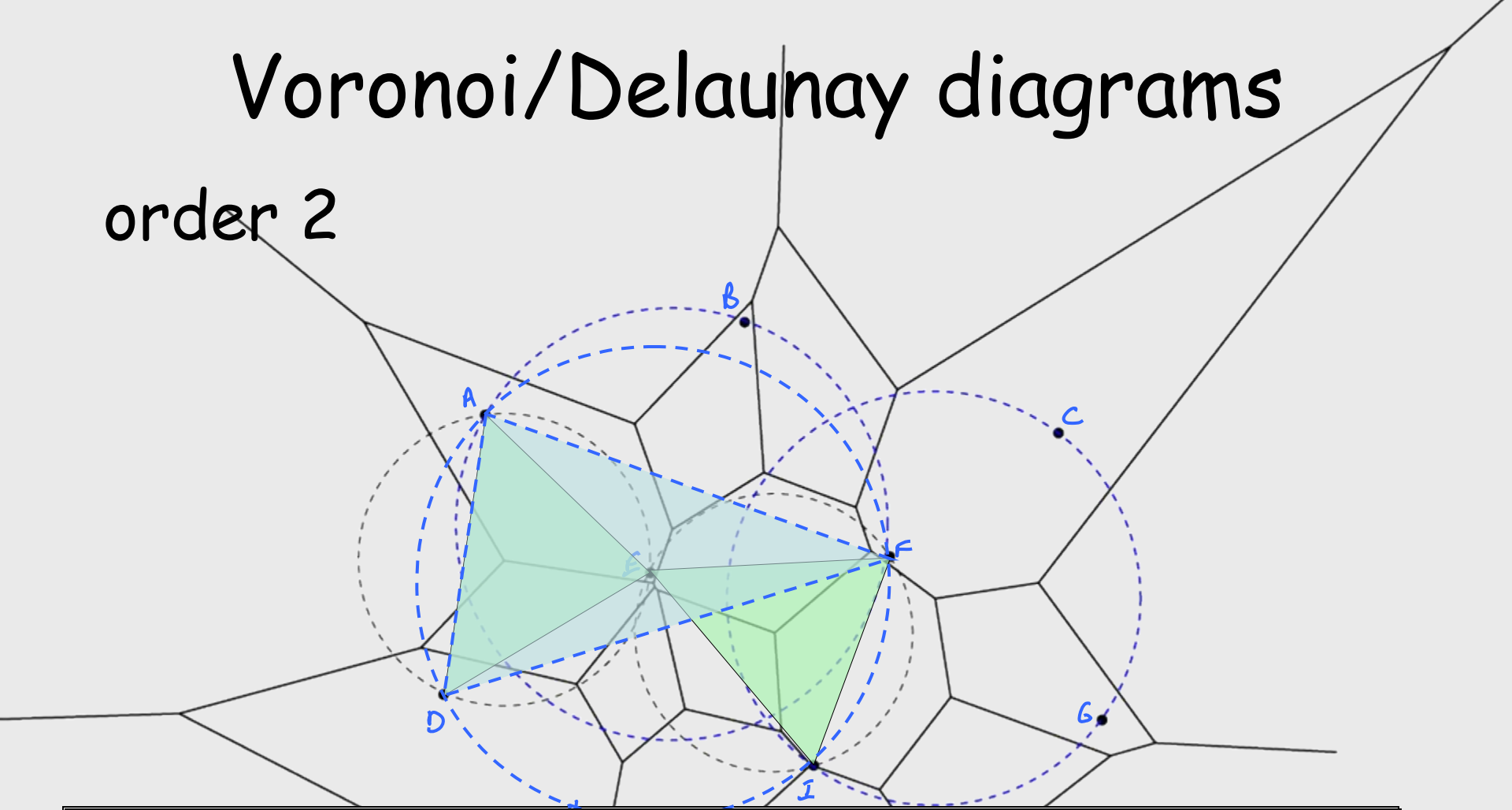
# Voronoi/Delaunay diagrams

order 2



# Voronoi/Delaunay diagrams

order 2

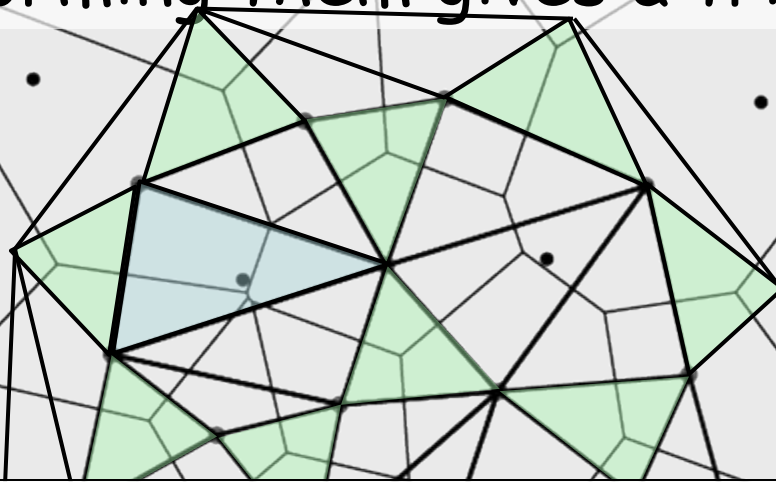


Delaunay triangles for these vertices may overlap, ...

# Voronoi/Delaunay diagrams

order 2

but transforming them gives a triangulation

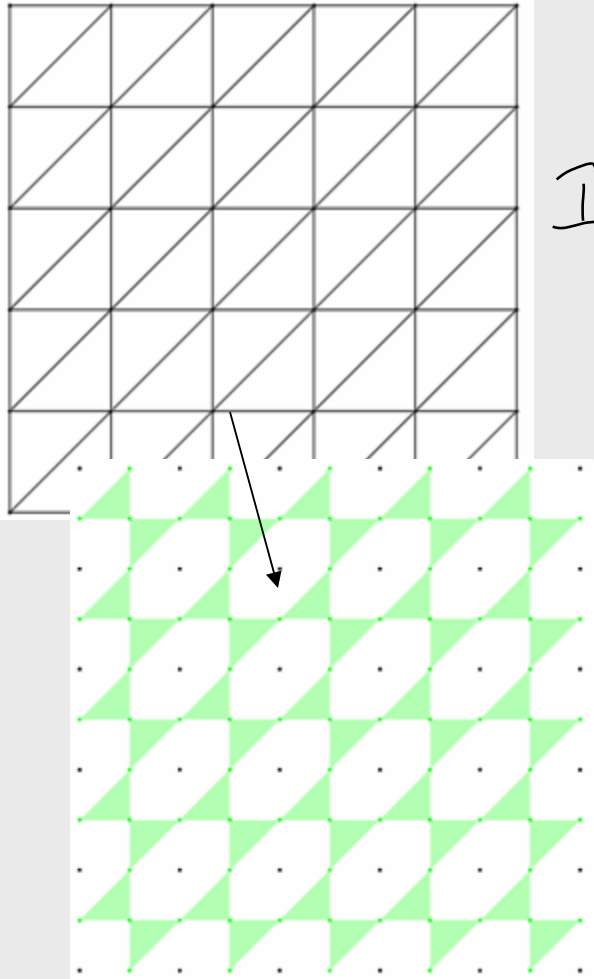


- Two ways to get the centroid triangulation:
- Project the lower hull of the centroids of all  $k$ -subsets of the lifted sites. [Aurenhammer 91]
  - Map *Delaunay configurations* to centroid triangles. [Schmitt 95, Andrzejak 97]



# Centroid triangulations

Dualizing Lee's algorithm to compute order  $k$  Voronoi diagrams

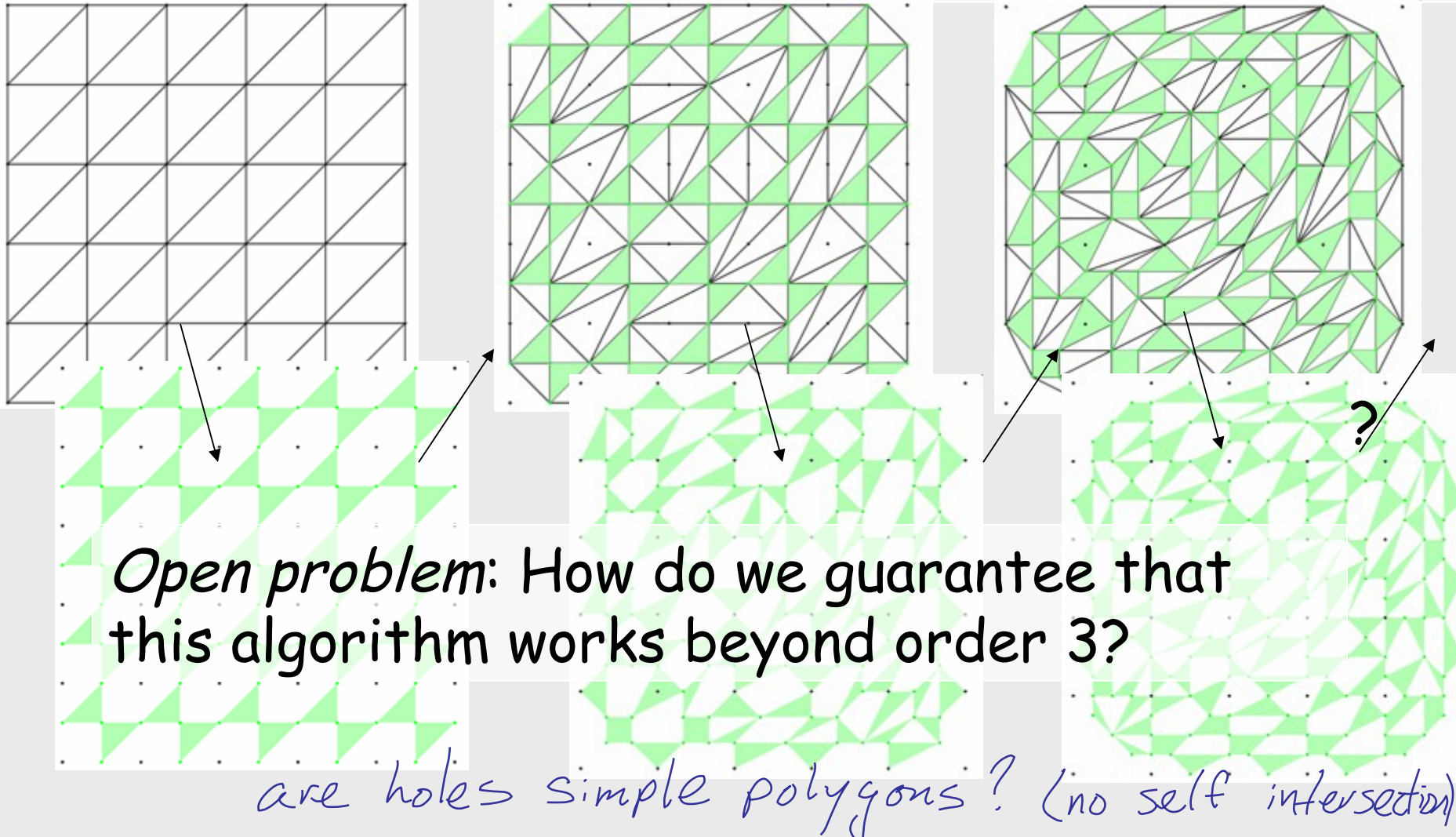


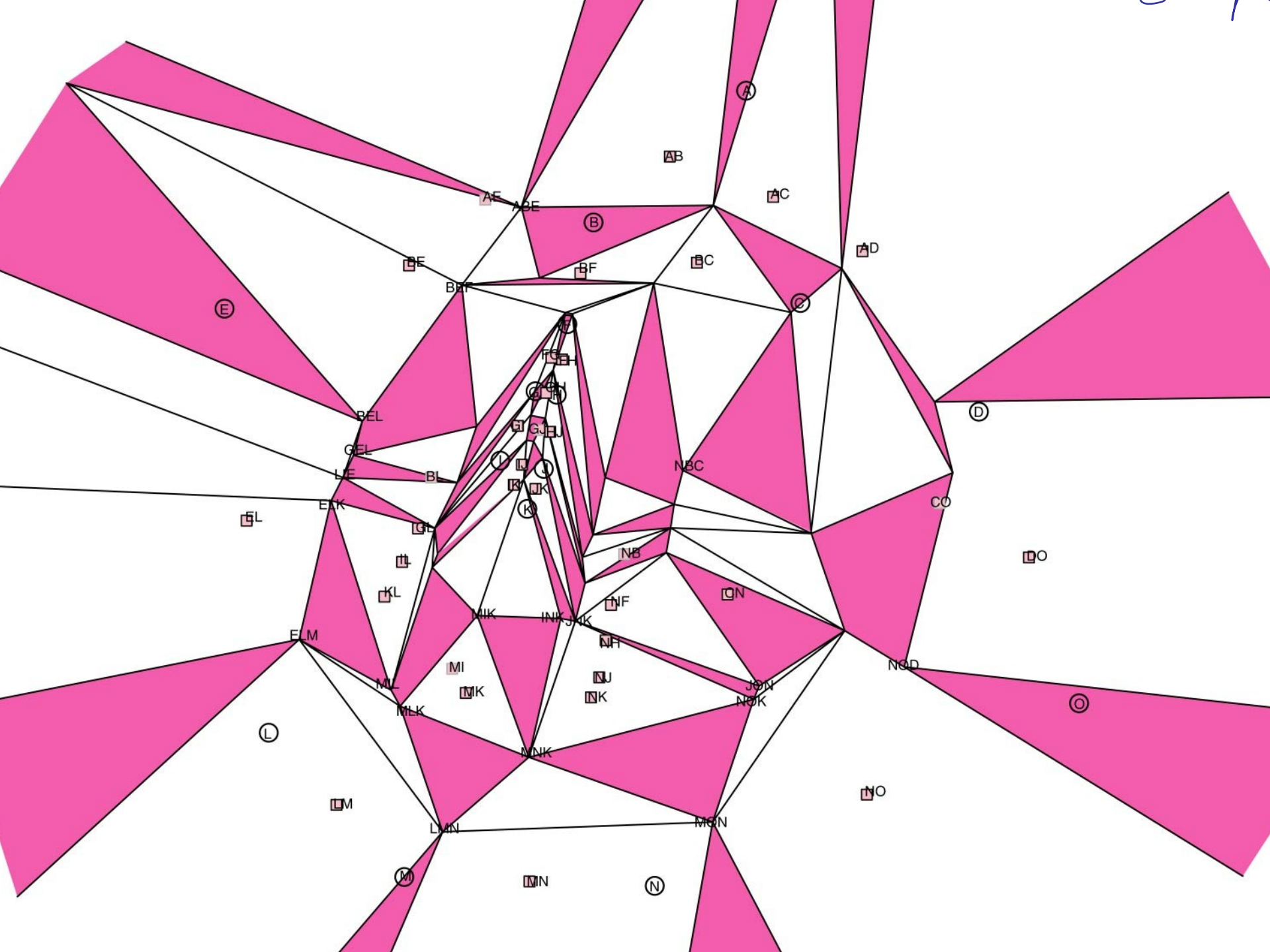
0. Begin w/ triangulation where every vertex has a unique label.  
Invariant: in our triangulation each vertex is the centroid of  $k$  points that make its label, and each edge joins vertices whose union has  $k+1$  labels.

1. For each edge, create the vertex that is its union.
2. For each triangle, join the three vertices created from its three edges
3. Discard original triangles & vertices
4. Complete to triangulation:

# Centroid triangulations

Dualizing Lee's algorithm to compute  
order  $k$  ~~Voronoi diagrams~~ centroid triangulations





# What are multivariate B-splines?

(multivariate) B-splines should define basis functions with no restriction on knot positions and have these properties of the classic B-splines:

- local support
- optimal smoothness
- partition of unity  $\sum B_i = 1$
- polynomial reproduction:  
for any degree  $k$  polynomial  $p$ ,  
with polar form  $P$ ,  
$$p = \sum P(S_{i+1} \dots S_{i+k}) B_i(\cdot | S_i \dots S_{i+k+1})$$

# Centroid triangulations

Delaunay configuration of degree  $k$ :  $(t, I)$ ,  
s.t. the circle through  $t$  contains exactly the  
 $k$  points of  $I$ .

Centroid triangle of order  $k$ :  $[A_{1..k}, B_{1..k}$  and  $C_{1..k}]$ , s.t.  
 $\#(A \cap B) = \#(B \cap C) = \#(A \cap C) = k-1$ .

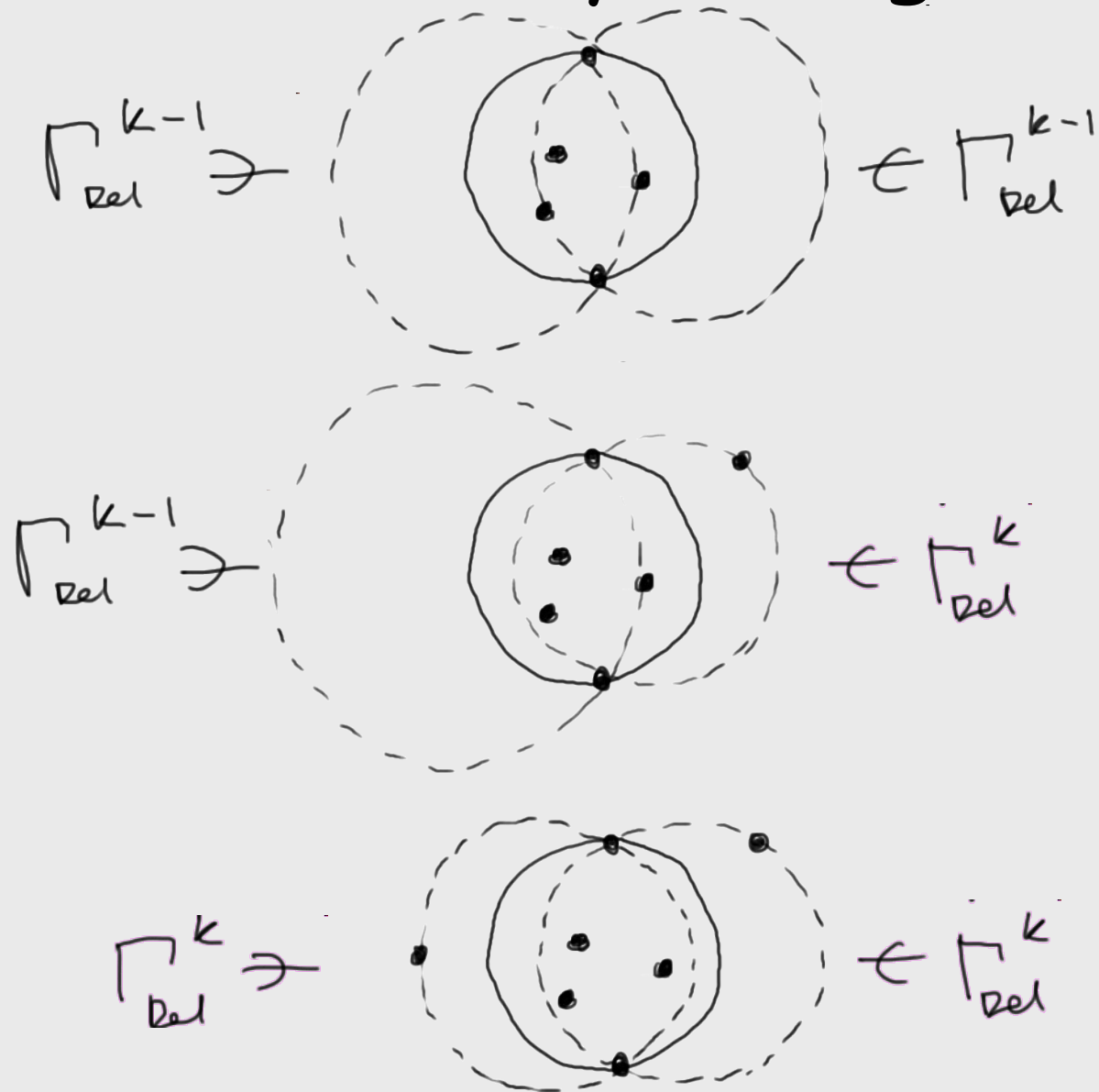
Map Delaunay configurations of deg.  $k-1, k-2$   
to centroid triangles of order  $k$ :

$(abc, J) \leftrightarrow [JU\{a\}, JU\{b\}, JU\{c\}]$   
deg.  $k-1$                       type 1

$(abc, I) \leftrightarrow [IU\{b,c\}, IU\{a,c\}, IU\{b,c\}]$   
deg.  $k-2$                       type 2

# Neamtu's use of Delaunay configs.

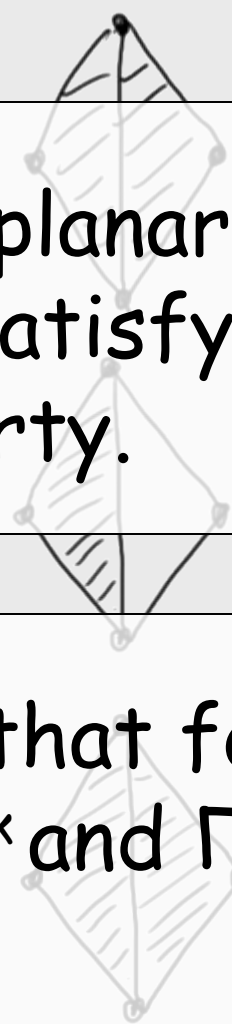
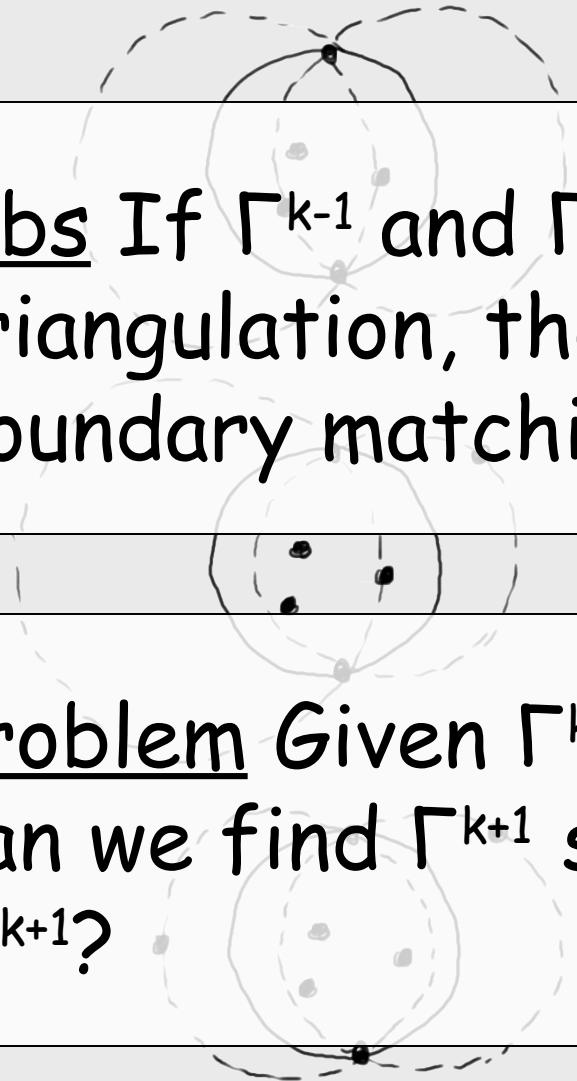
Key property  
for proof of  
polynomial  
repro. is  
*boundary  
matching:*





# Centroid Triangulations

relations of conf.  $\leftrightarrow$  triangle neighbors



Obs If  $\Gamma^{k-1}$  and  $\Gamma^k$  form a planar centroid triangulation, then they satisfy the boundary matching property.

Problem Given  $\Gamma^{k-1}$  and  $\Gamma^k$  that form  $\Delta^k$ , can we find  $\Gamma^{k+1}$  so that  $\Gamma^k$  and  $\Gamma^{k+1}$  form  $\Delta^{k+1}$ ?

# Centroid Triangulations

For a set of sites  $S$ , a *centroid triangle* of order  $k$  have vertices that are centroids of  $k$ -subsets of  $S$ , e.g.  $A_{1..k}$ ,  $B_{1..k}$  and  $C_{1..k}$ , and satisfy that

$$\#(A \cap B) = \#(B \cap C) = \#(A \cap C) = k-1.$$

Obs  $\#(A \cap B \cap C) = k-1$  or  $k-2$ .

type 1   type 2

Obs There is a 1-1 mapping between the centroid triangles of order  $k$  and conf of degree  $k-1$  and  $k-2$ .

$$(abc, J) \leftrightarrow [JU\{a\}, JU\{b\}, JU\{c\}],$$

$$(abc, I) \leftrightarrow [IU\{b,c\}, IU\{a,c\}, IU\{b,c\}]$$



# Centroid triangulations -> B-splines

## Theorem

Let  $\Gamma^0 \dots \Gamma^k$  be a sequence of configurations. If  $\Gamma^0$  is a triangulation, and  $\Gamma^{i-1}, \Gamma^i$  form a centroid triangulation for  $0 < i < k$ , then the simplex splines assoc. with  $\Gamma^k$  reproduce polynomials of deg.  $k$ .

for any deg.  $k$  polynomial  $p$ , with polar form  $P$

$$p = \sum_{(t, I) \in \Gamma^k} P(I) d(t) M(. | t \cup I)$$

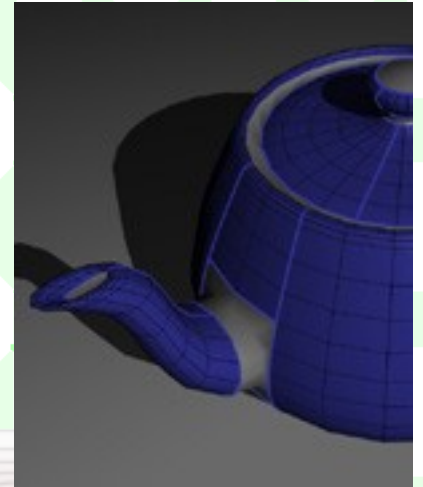
classic B-spline

for any deg.  $k$  polynomial  $p$ , with polar form  $P$

$$p = \sum P(S_{i+1} \dots S_{i+k}) B_i(. | S_i \dots S_{i+k+1})$$

# Reproducing the Zwart-Powell element

- An example of our claim:  
*Box splines are special centroid triangulation splines.*
- Theory motivation: Evidence that 'ct's provide general basis for bivariate splines.
- Practical motivation: Smooth blending of box spline patches.



# Polyhedral splines

*Polyhedron spline*  $M_{\Pi}(x | P)$

$P$ : a polyhedron in  $\mathbb{R}^n$

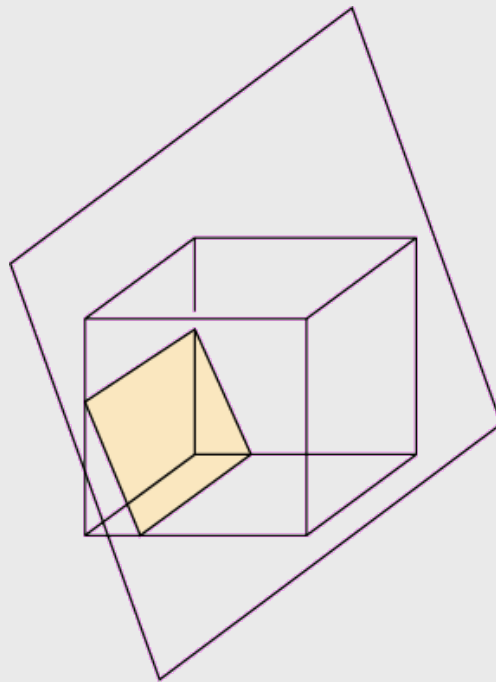
$\Pi$ : a projection matrix from  $\mathbb{R}^n$  to  $\mathbb{R}^m$

is an  $n$ -variate, degree  $(n-m)$  spline

$$M_{\Pi}(x | P) := \frac{\text{vol} \{ y \mid y \in P, \Pi y = x \}}{\text{vol } P}$$

*Box Spline*

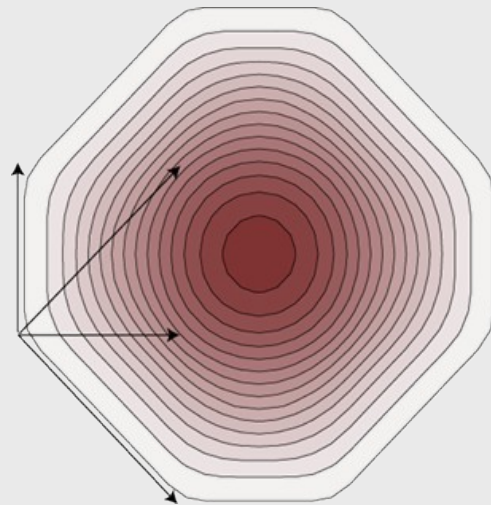
$$M_{\Pi}(x)$$



# Reproduction of Box splines

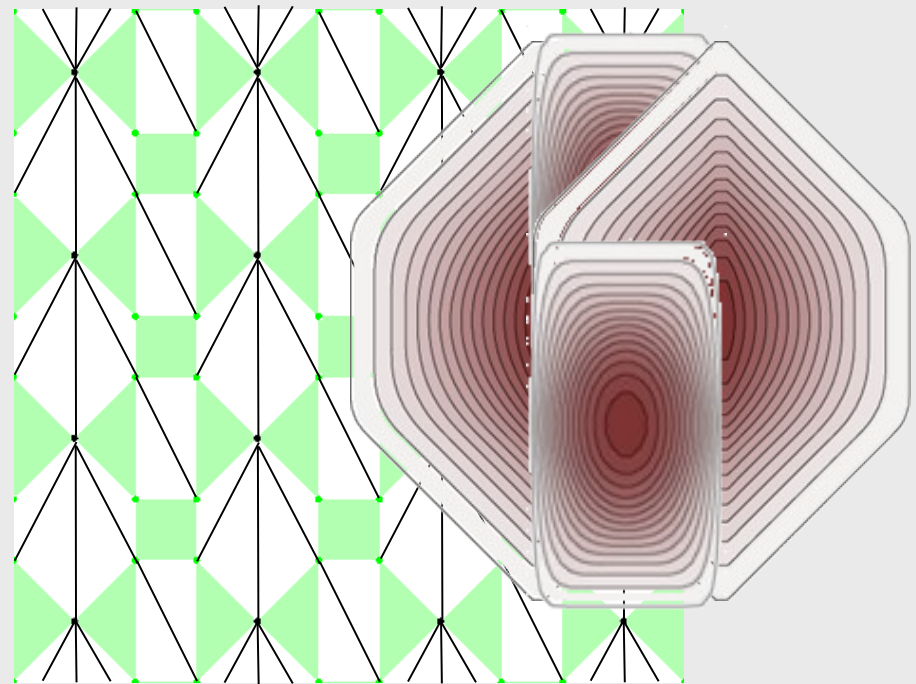
$$\text{span}_{v \in \mathbb{N} \times \mathbb{N}} \{ M_{\Pi}(x + v) \} \subset \text{span}_{X \in \Gamma^k} \{ M(x | X) \}$$

ZP element



$$\begin{aligned} \Pi &= \begin{pmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & -1 \end{pmatrix} \\ &= \{ \rightarrow \uparrow \nearrow \searrow \} \end{aligned}$$

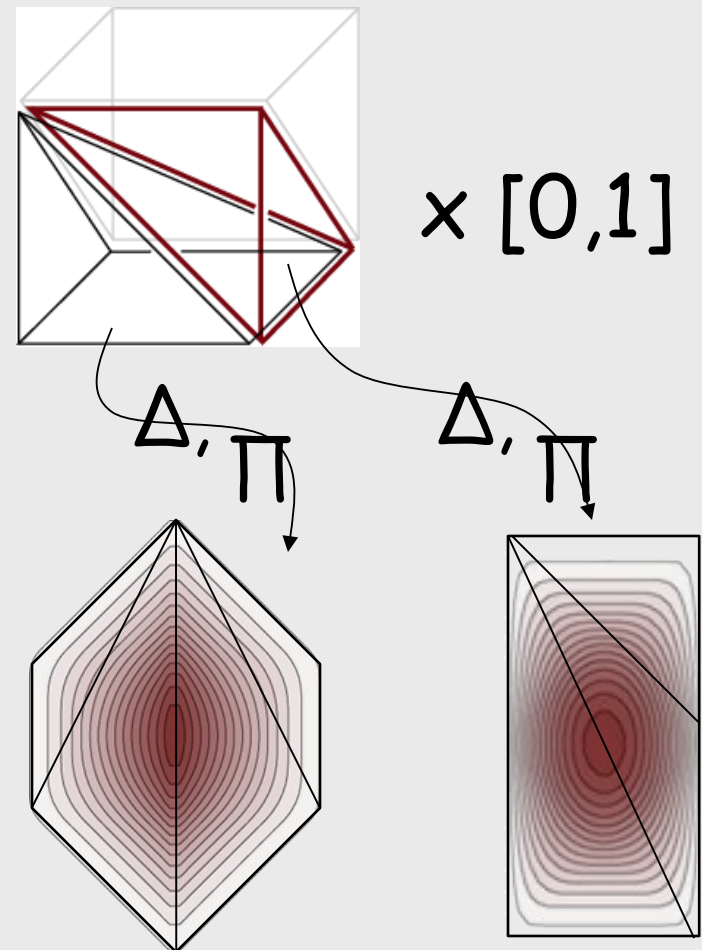
order 2 centroid triangulation



# Reproduction of Box splines

*Proof sketch* (reproducing a single ZP element)

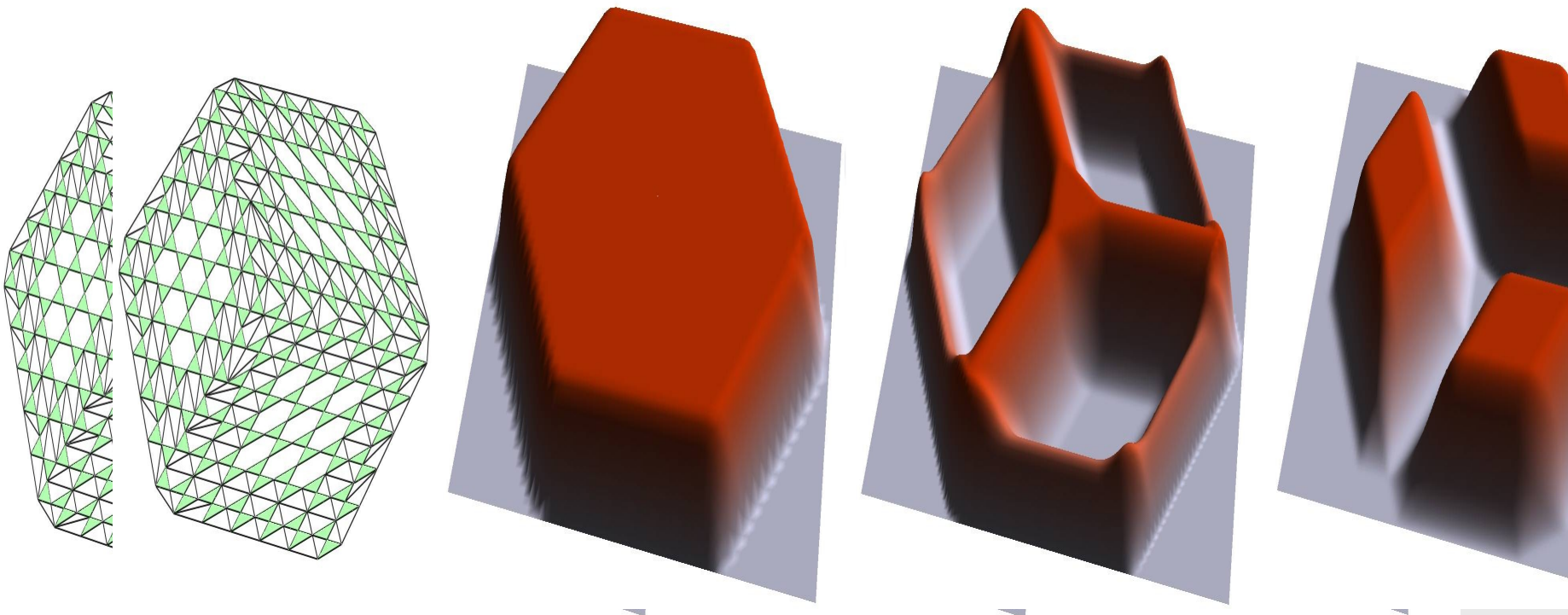
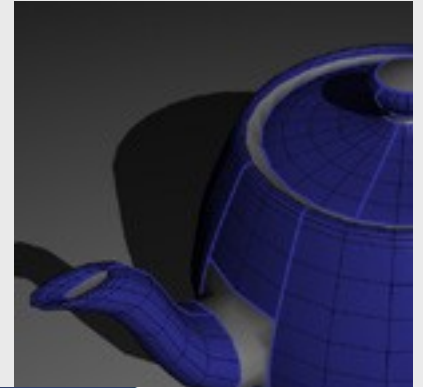
$\equiv$  ZP-element  
 $\equiv$  4-cube  
(partition)  
 $\rightarrow$  4-polytopes  
(triangulate)  
 $\rightarrow$  4-simplices  
 $\equiv$  simplex splines





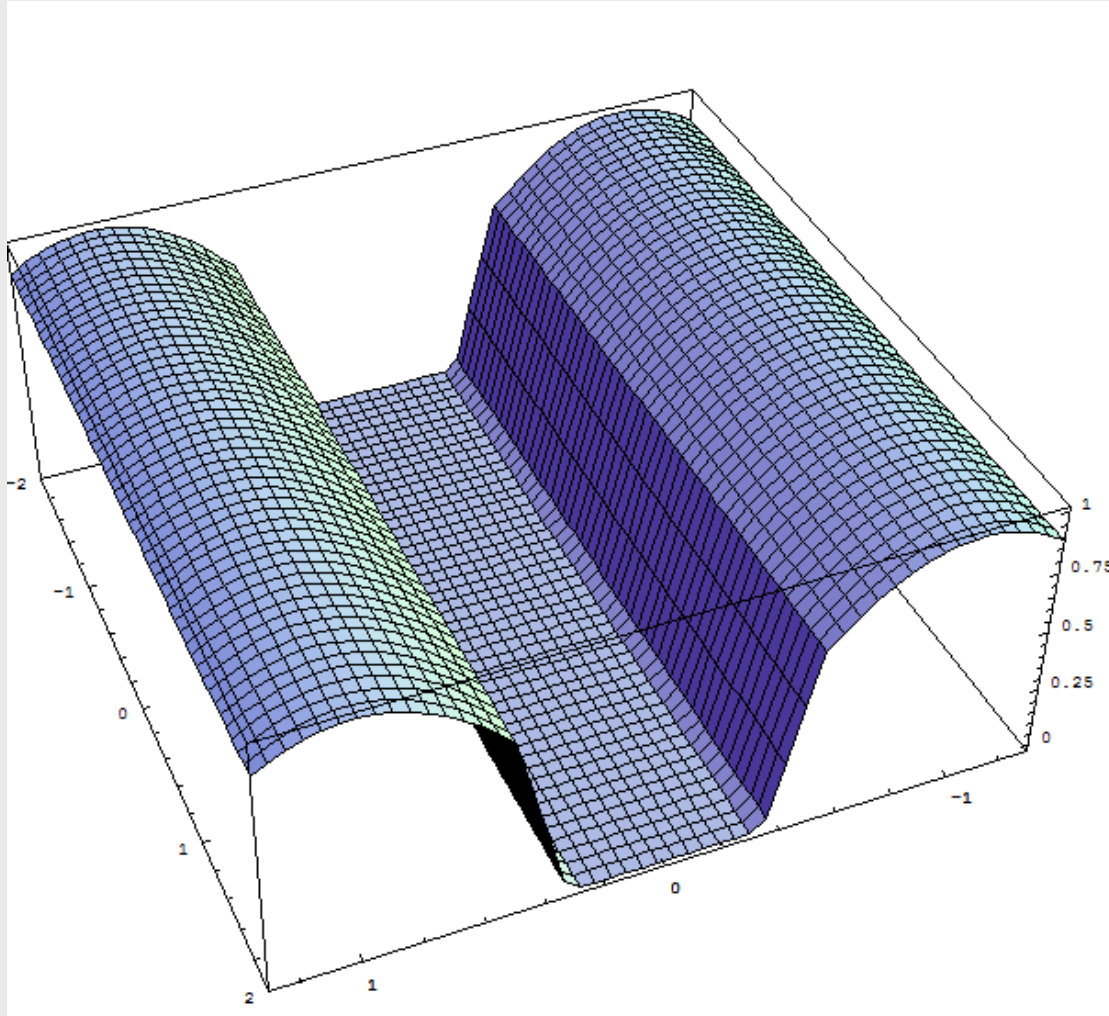
# Reproduction of Box-splines

- Patch blending  
E.g., a partition of unity by  
blending regular splines.



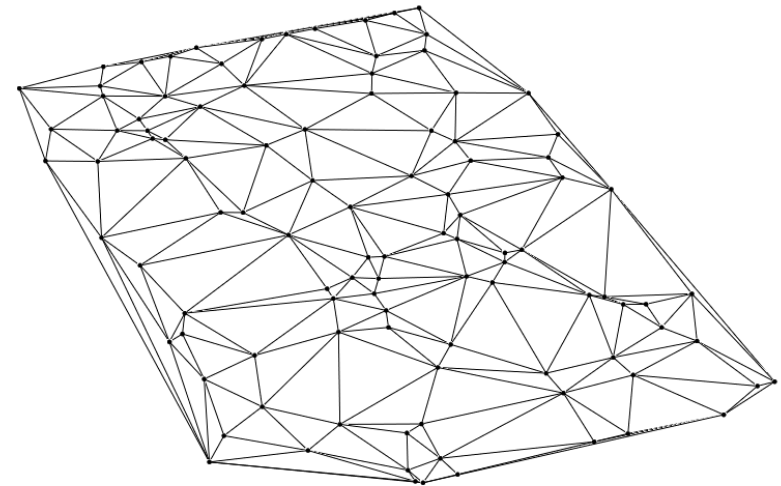
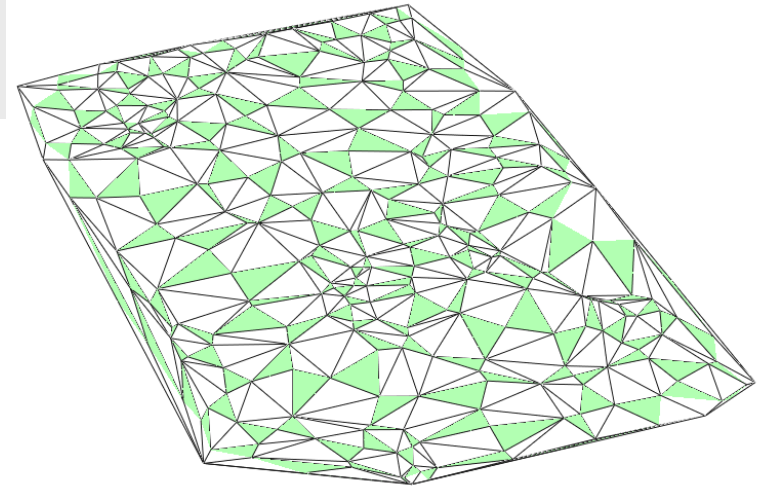
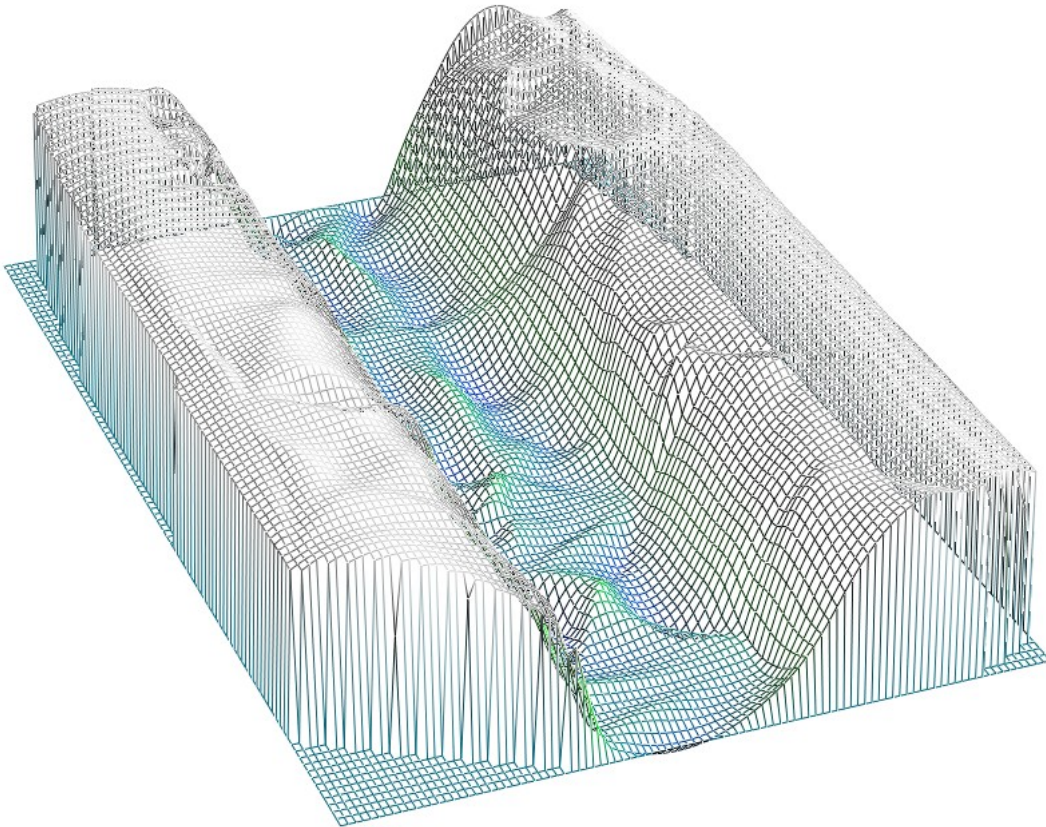
# Modeling sharp features

Data fitting with scattered data points  
and "break lines"



# Modeling sharp features

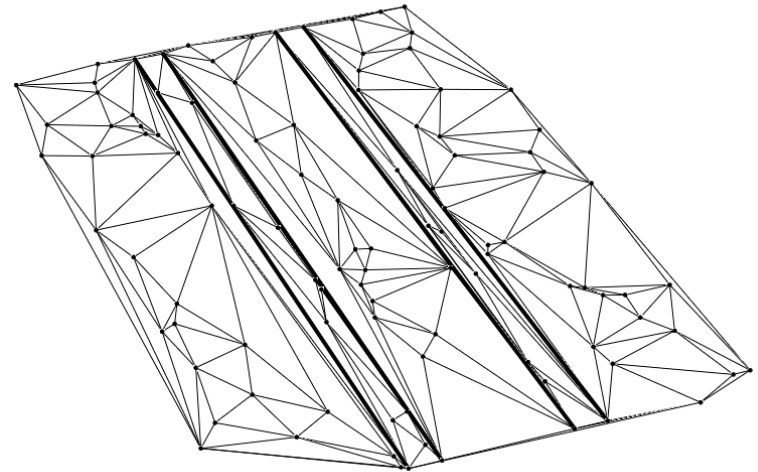
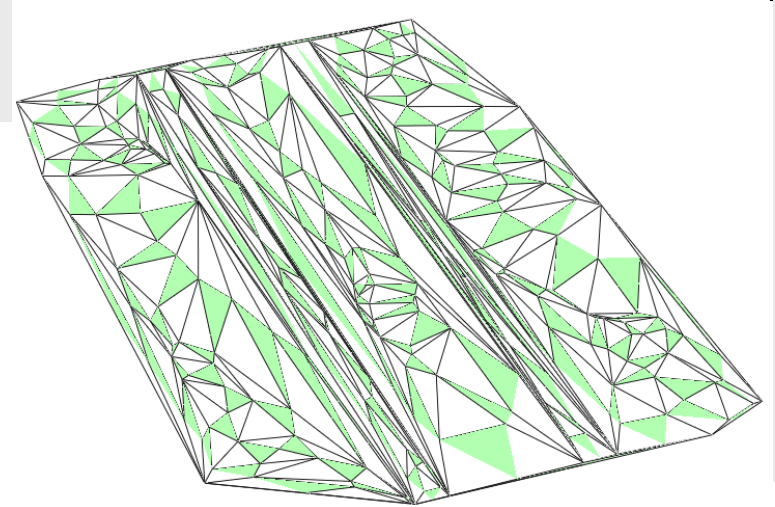
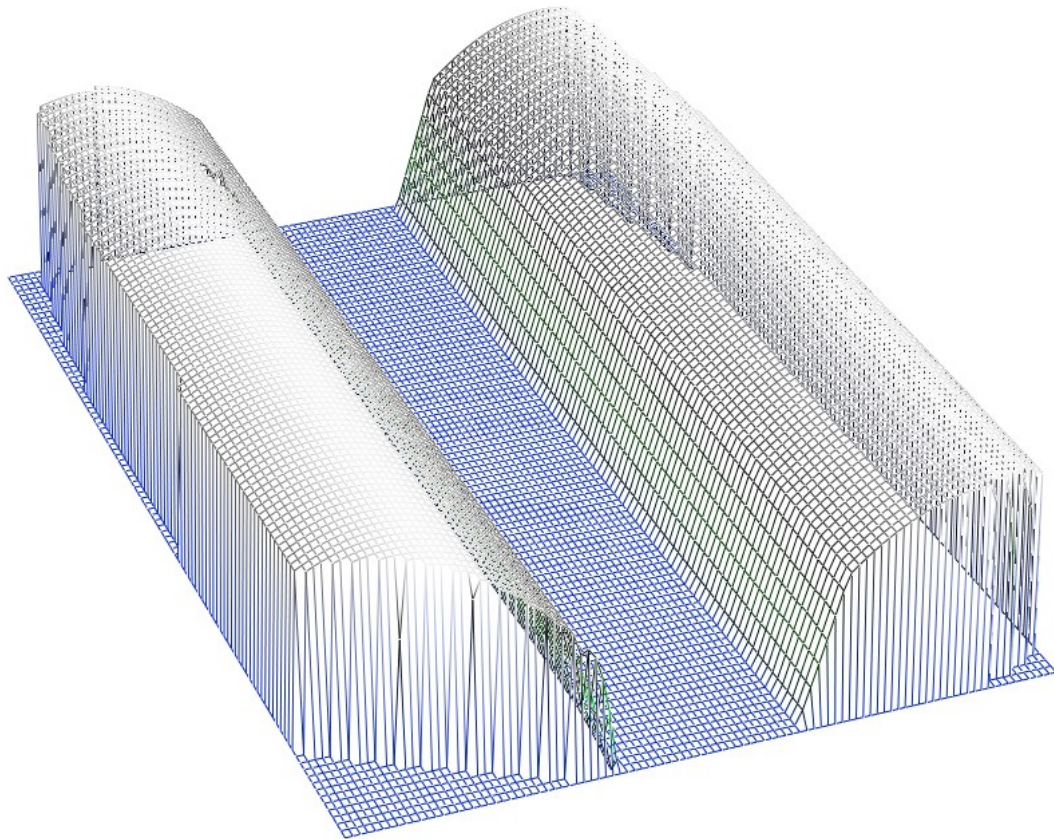
Data fitting with scattered data points  
and "break lines"





# Modeling sharp features

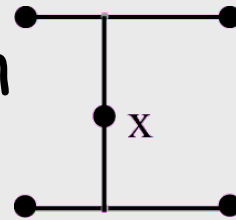
Data fitting with scattered data points  
and "break lines"



# Open Problems

- Prove no self-intersecting holes arise in the centroid triangulation algorithm
- Reproduce other box-splines by centroid triangulation

- bilinear interpolation  
(quadratic)



$$\Pi = \{\rightarrow \rightarrow \uparrow \uparrow\}$$

- loop subdivision  
(quartic)

$$\Pi = \{\rightarrow \rightarrow \uparrow \uparrow \nearrow \nearrow\}$$

