

Large gaps in sets of primes and other sequences

Kevin Ford

University of Illinois at Urbana-Champaign

October, 2018

Large gaps between primes

$G(x) := \max_{p_n \leq x} (p_n - p_{n-1})$, p_n is the n^{th} prime.

2, 3, 5, 7, ..., 109, 113, 127, 131, ..., 9547, 9551, 9587, 9601, ...

Upper bound: $G(x) \ll x^{0.525}$ (Baker-Harman-Pintz, 2001).

Large gaps between primes

$G(x) := \max_{p_n \leq x} (p_n - p_{n-1})$, p_n is the n^{th} prime.

2, 3, 5, 7, ..., 109, 113, 127, 131, ..., 9547, 9551, 9587, 9601, ...

Upper bound: $G(x) \ll x^{0.525}$ (Baker-Harman-Pintz, 2001).

Rankin, 1938. $G(x) \geq c(\log x) \frac{\log_2 x \log_4 x}{(\log_3 x)^2}$, $c = \frac{1}{3}$.

Improvements to the constant c until 1997 (Pintz, $c = 2e^\gamma$).

Large gaps between primes

$G(x) := \max_{p_n \leq x} (p_n - p_{n-1})$, p_n is the n^{th} prime.

2, 3, 5, 7, ..., 109, 113, 127, 131, ..., 9547, 9551, 9587, 9601, ...

Upper bound: $G(x) \ll x^{0.525}$ (Baker-Harman-Pintz, 2001).

Rankin, 1938. $G(x) \geq c(\log x) \frac{\log_2 x \log_4 x}{(\log_3 x)^2}$, $c = \frac{1}{3}$.

Improvements to the constant c until 1997 (Pintz, $c = 2e^\gamma$).

Erdős Conjecture. (\$10,000) Rankin's bound is true, for *any* $c > 0$.

Solved: Ford-Green-Konyagin-Tao and Maynard (arXiv, Aug-2014).

Theorem (Ford-Green-Konyagin-Maynard-Tao (2018))

For large x , $G(x) \gg \log x \frac{\log_2 x \log_4 x}{\log_3 x}$.

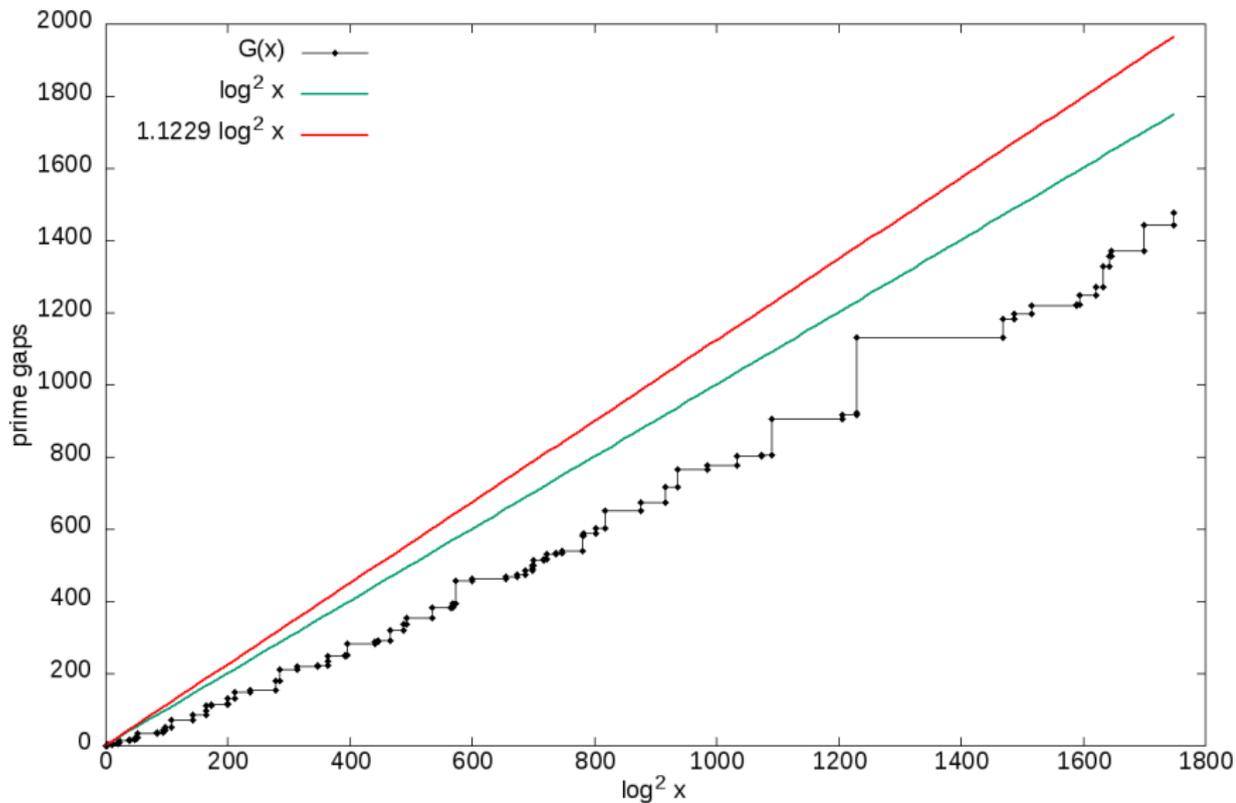
Conjectures

Cramér (1936):

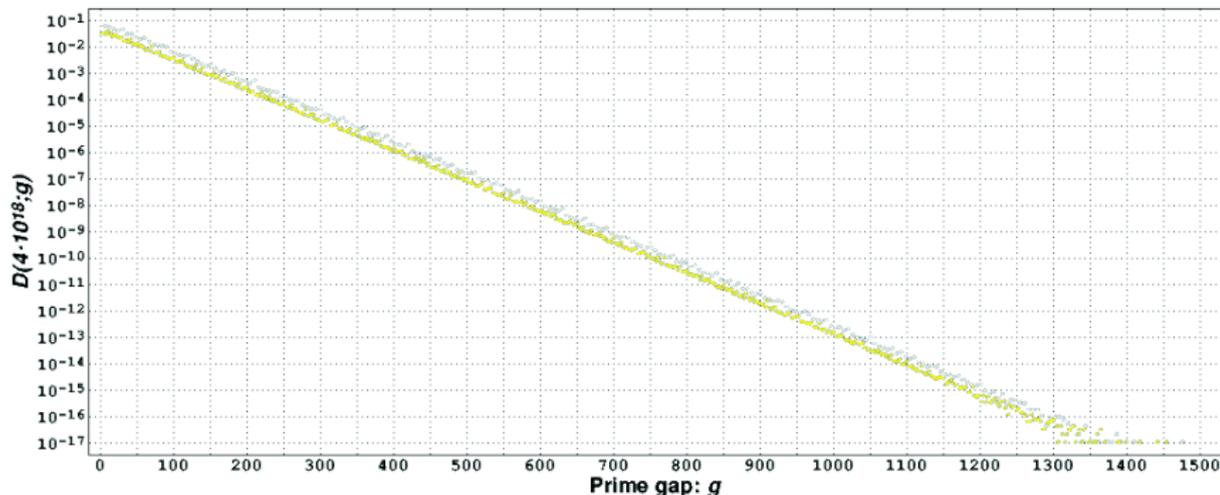
- $\limsup_{x \rightarrow \infty} \frac{G(x)}{\log^2 x} = 1.$
- $\frac{p_n - p_{n-1}}{\log p_n}$ has approximate exponential distribution.

Granville (1995): $\limsup_{x \rightarrow \infty} \frac{G(x)}{\log^2 x} \geq 2e^{-\gamma} = 1.1229 \dots$

Computational evidence, up to 10^{18}



Exponential distribution of gaps



Prime gap statistics, $p_n < 4 \cdot 10^{18}$

Gallagher, 1976. Prime k -tuples conjecture \Rightarrow exponential prime gap distribution

Cramér's model defects: small gaps

Cramér's model produces a set $\mathcal{P} \in \mathbb{N}$ of “random primes”:

$$\mathbb{P}(n \in \mathcal{P}) = 1/\log n \quad (n \geq 3).$$

Theorem. With probability 1, $\#\{n : n, n + 1 \in \mathcal{P}\} = \infty$

This does not hold for real primes!

Cramér's model defects: small gaps

Cramér's model produces a set $\mathcal{P} \in \mathbb{N}$ of “random primes”:

$$\mathbb{P}(n \in \mathcal{P}) = 1/\log n \quad (n \geq 3).$$

Theorem. With probability 1, $\#\{n : n, n + 1 \in \mathcal{P}\} = \infty$

This does not hold for real primes!

Theorem. With probability 1,

$$\#\{n \leq x : n, n + 2 \in \mathcal{P}\} \sim \frac{x}{\log^2 x}.$$

Conjecture (Hardy-Littlewood, 1923).

$$\#\{n \leq x : n, n + 2 \text{ prime}\} \sim C \frac{x}{\log^2 x},$$

where $C = 2 \prod_{p>2} (1 - 1/(p-1)^2) \approx 1.3203$

General Cramér type model

Theorem (classical, 1960s)

Choose N random points in $[0, x]$. With high probability, the max. gap is $\sim \frac{\log N}{N} x$.

Conjecture (Hardy-Littlewood; Bateman-Horn)

Let f_1, \dots, f_k be distinct, irreducible polynomials $f_i : \mathbb{Z} \rightarrow \mathbb{Z}$ with pos. leading coeff., and $f_1 \cdots f_k$ has no fixed prime factor. Then

$$\#\{n \leq x : f_1(n), \dots, f_k(n) \text{ all prime}\} \sim C \frac{x}{(\log x)^k},$$

where $C = C(f_1, \dots, f_k) > 0$ is constant.

General Cramér type model

Theorem (classical, 1960s)

Choose N random points in $[0, x]$. With high probability, the max. gap is $\sim \frac{\log N}{N} x$.

Conjecture (Hardy-Littlewood; Bateman-Horn)

Let f_1, \dots, f_k be distinct, irreducible polynomials $f_i : \mathbb{Z} \rightarrow \mathbb{Z}$ with pos. leading coeff., and $f_1 \cdots f_k$ has no fixed prime factor. Then

$$\#\{n \leq x : f_1(n), \dots, f_k(n) \text{ all prime}\} \sim C \frac{x}{(\log x)^k},$$

where $C = C(f_1, \dots, f_k) > 0$ is constant.

For $\{n \leq x : f_1(n), \dots, f_k(n) \text{ all prime}\}$, the model prediction:

$$\text{average gap} \sim \frac{(\log x)^k}{C}, \quad \text{maximal gap} \sim \frac{(\log x)^{k+1}}{C}.$$

Polynomial gaps

Bunyakovsky (1857) Conj: Infinitely many primes $p = n^2 + 1$.

Bateman-Horn Conj: $\#\{n \leq x : n^2 + 1 \text{ prime}\} \sim C \frac{x}{\log x}$.

Cramér type heuristic: max. gap $\sim \frac{(\log x)^2}{C}$.

Sieve methods: $\#\{n \leq x : n^2 + 1 \text{ prime}\} \ll \frac{x}{\log x}$.

Polynomial gaps

Bunyakovsky (1857) Conj: Infinitely many primes $p = n^2 + 1$.

Bateman-Horn Conj: $\#\{n \leq x : n^2 + 1 \text{ prime}\} \sim C \frac{x}{\log x}$.

Cramér type heuristic: max. gap $\sim \frac{(\log x)^2}{C}$.

Sieve methods: $\#\{n \leq x : n^2 + 1 \text{ prime}\} \ll \frac{x}{\log x}$.

Question: Can one *prove* that large strings of consecutive composite values of $n^2 + 1$ exist? i.e., strings longer than $O(\log x)$ below x .

Problem: Methods for prime gaps $G(x)$ **do not work!**

Polynomial gaps

Bunyakovsky (1857) Conj: Infinitely many primes $p = n^2 + 1$.

Bateman-Horn Conj: $\#\{n \leq x : n^2 + 1 \text{ prime}\} \sim C \frac{x}{\log x}$.

Cramér type heuristic: $\max. \text{ gap} \sim \frac{(\log x)^2}{C}$.

Sieve methods: $\#\{n \leq x : n^2 + 1 \text{ prime}\} \ll \frac{x}{\log x}$.

Question: Can one *prove* that large strings of consecutive composite values of $n^2 + 1$ exist? i.e., strings longer than $O(\log x)$ below x .

Problem: Methods for prime gaps $G(x)$ do not work!

Theorem (Ford-Konyagin-Maynard-Pomerance-Tao, 2018+)

Let $f : \mathbb{Z} \rightarrow \mathbb{Z}$ be a monic, irreducible polynomial with no fixed prime factor. Then there is a string of $\gg (\log x)(\log \log x)^c$ integers $n \leq x$ for which $f(n)$ is composite. Here c depends only on f .

Proving large prime gaps: Jacobsthal's function

$$\mathcal{S}_x = \{n \in \mathbb{Z} : (n, Q_x) = 1\}, \quad Q_x = \prod_{p \leq x} p.$$

(i.e., sieve of Eratosthenes using primes $p \leq x$)

Main goal: Find $J(x)$, **the largest gap in \mathcal{S}_x** ; long string of consecutive integers all with a small prime factor.

Cor: $G(2Q_x) \geq J(x)$; essentially $G(X) \gtrsim J(\log X)$.

Proving large prime gaps: Jacobsthal's function

$$\mathcal{S}_x = \{n \in \mathbb{Z} : (n, Q_x) = 1\}, \quad Q_x = \prod_{p \leq x} p.$$

(i.e., sieve of Eratosthenes using primes $p \leq x$)

Main goal: Find $J(x)$, the largest gap in \mathcal{S}_x ; long string of consecutive integers all with a small prime factor.

Cor: $G(2Q_x) \geq J(x)$; essentially $G(X) \gtrsim J(\log X)$.

Bounding $J(x)$:

- Average gap $Q_x / \phi(Q_x) \sim e^\gamma \log x$;
- (trivial) $J(x) \geq x - 2$ since $[2, x] \cap \mathcal{S}_x = \emptyset$;
- (FGKMT, 2018). $J(x) \gg x(\log x) \frac{\log_3 x}{\log_2 x}$.
- (Iwaniec, 1978). $J(x) \ll x^2(\log x)^2$.

Conjecture (Maier-Pomerance, 1990). $J(x) = x(\log x)^{2+o(1)}$.

Random Cramér type model: $J(x) \sim \frac{Q_x \log Q_x}{\phi(Q_x)} \sim e^\gamma x \log x$.

Least prime in an arithmetic progression

Let $p(k, l) = \min\{p \equiv l \pmod{k}, \text{prime}\}$, $M(k) = \max_{(l, k)=1} p(k, l)$.

Upper bounds

Linnik, 1944. $M(k) \ll k^L$. (**Xylouris** : $L = 5.18$).

ERH: $L = 2 + \varepsilon$; Chowla conjecture: $L = 1 + \varepsilon$.

Lower bounds

Trivial: $M(k) \gg \phi(k) \log k$.

Prachar; Schinzel - 1961/62. For infinitely many k ,

$$M(k) \gg \phi(k) \log k \frac{\log_2 k \log_4 k}{(\log_3 k)^2}. \quad (1)$$

Wagstaff (1978) - (1) holds for all prime k .

Pomerance (1980) - (1) holds for almost all k , in fact all k with at most $(\log k)^{c/\log_3 k}$ prime factors.

Least prime in an arithmetic progression, II

Pomerance: $M(k) \gg \phi(k) \log k \frac{\log_2 k \log_4 k}{(\log_3 k)^2}$ for almost all k .

Lemma (Pomerance): Let $j(m)$ be the maximal gap between numbers coprime to m . If $0 < m \leq k/j(k)$ and $(m, k) = 1$ then $M(k) > kj(m)$.

Take $m = \prod_{\substack{p \leq (1-\delta) \log k \\ p \nmid k}} p$ need a lower bound on $j(m)$.

Corollary (FGKMT, 2018). If k has no prime factor $\leq \log k$, then

$$M(k) \gg \phi(k) \log k \frac{\log_2 k \log_4 k}{\log_3 k}. \quad (2)$$

Theorem (Junxian Li, Kyle Pratt and George Shakan, 2018)

Inequality (2) holds for almost all k ; in fact, for all k with at most $\exp\{(1/2 - \varepsilon) \frac{\log_2 k \log_4 k}{\log_3 k}\}$ prime factors.

Least Prime in an A.P. – conjectures

Conjecture (folklore): $M(k) \ll k \log^{2+\varepsilon} k$.

Conjecture (Wagstaff, 1979): $M(k) \sim \phi(k) \log^2 k$ for “most k ”

Conjecture (Li-Pratt-Shakan, 2018)

$$\liminf_{k \rightarrow \infty} \frac{M(k)}{\phi(k) \log^2 k} = 1, \quad \limsup_{k \rightarrow \infty} \frac{M(k)}{\phi(k) \log^2 k} = 2.$$

Heuristic: coupon collectors problem.

Least prime in an AP: data

Conjecture (Li-Pratt-Shakan, 2018)

$$\liminf_{k \rightarrow \infty} \frac{M(k)}{\phi(k) \log^2 k} = 1, \quad \limsup_{k \rightarrow \infty} \frac{M(k)}{\phi(k) \log^2 k} = 2.$$

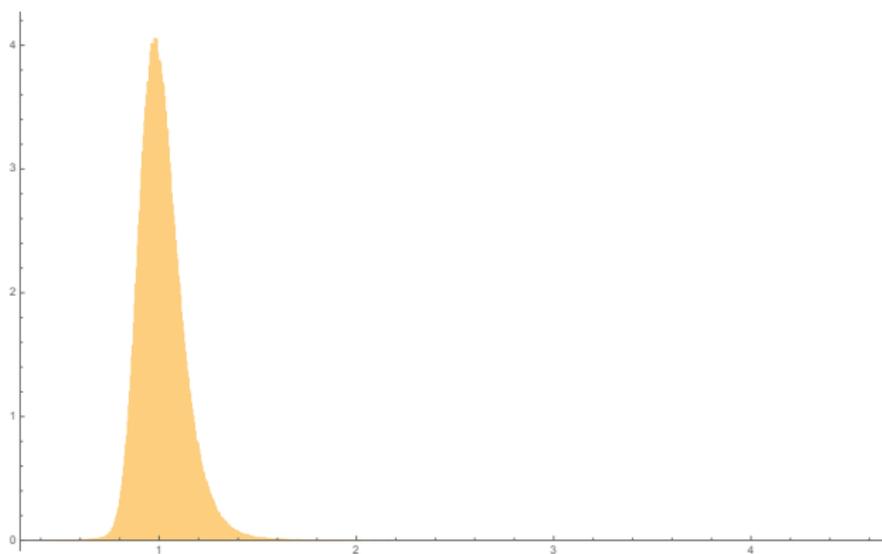


Figure: Histogram for $\frac{M(k)}{\phi(k) \log(\phi(k)) \log k}$ for $k \leq 10^6$

Covering the gap

Covering: $J(x)$ is the largest y so that there are a_2, a_3, a_5, \dots with

$$\{a_p \pmod p : p \leq x\} \supseteq [0, y]$$

Proof: If $[n, n + y]$ is a gap in S_x , $y = J(x)$, define a_p for $p \leq x$ by

$$(-n \pmod{Q_x}) = \bigcap_{p \leq x} (a_p \pmod p).$$

Goal: succeed with y a bit larger than x .

Finding large gaps in \mathcal{S}_x

Need $\bigcap_{p \leq x} (a_p \bmod p) \supseteq [0, y]$, with y a bit larger than x .



$$y = cx \frac{\log x \log_3 x}{(\log_2 x)^2}, \quad z = x^c \frac{\log_3 x}{\log_2 x} \quad \text{Want } \{a_p \bmod p : p \leq x\} \supseteq [0, y]$$

Classical 3-stage-process (Westzynthius-Erdős-Rankin)

- 1 (Key!!) Take $a_p = 0$ for $p \in (z, \frac{x}{2}] \cap [2, \frac{2y}{x}]$. Left uncovered: z -smooth numbers (few for appropriate z) and primes;
 $\sim \frac{y}{\log y}$ numbers uncovered.

A typical choice of a_p leaves $\sim y \frac{\log z}{\log y}$ uncovered numbers

- 2 Greedy choice for $a_p, p \in (2y/x, z]$
- 3 (trivial) for $p \in (\frac{x}{2}, x]$, choose a_p to cover one uncovered element from step 2. Success if $\leq \pi(x) - \pi(x/2) \sim \frac{x}{2 \log x}$ such elements.

New bounds on $J(x)$ [FKMPT, 2018]



$$y = cx \frac{\log x \log_3 x}{\log_2 x}, z = x^c \frac{\log_3 x}{\log_2 x} \quad \text{Want } \{a_p \bmod p : p \leq x\} \supseteq [0, y]$$

- 1 $a_p = 0$ for $p \in (z, x/4] \cap [2, \log^{10} x]$. Uncovered: z -smooth numbers and primes;
- 2 Random, uniform choice of a_p , $\log^{10} x < p \leq z$.
- 3 Strategic choice of a_p , $x/4 < p \leq x/2$ to cover many remaining elements; some AP modulo p has many primes in $[0, y]$.
- 4 (trivial) Use single a_p for each $x/2 < p \leq x$ to cover each remaining uncovered element.

Tools: Maynard sieve, efficient hypergraph covering

Analog of $J(x)$ for polynomials

$$\mathcal{S}_f(x) = \{n \in \mathbb{Z} : (f(n), Q_x) = 1\}, \quad Q_x = \prod_{p \leq x} p.$$

Gaps: Let $J_f(x)$ be the largest gap in $\mathcal{S}_f(x)$.

Analog of $J(x)$ for polynomials

$$\mathcal{S}_f(x) = \{n \in \mathbb{Z} : (f(n), Q_x) = 1\}, \quad Q_x = \prod_{p \leq x} p.$$

Gaps: Let $J_f(x)$ be the largest gap in $\mathcal{S}_f(x)$.

Covering problem: Let $I_p = \{n \bmod p : f(n) \equiv 0 \pmod{p}\}$.
 $J_f(x)$ is the largest y so that $[0, y]$ is covered by

$$\{b_p + \nu_p \bmod p : p \leq x, \nu_p \in I_p\}$$

for some residues $b_p \bmod p$.

Analog of $J(x)$ for polynomials

$$\mathcal{S}_f(x) = \{n \in \mathbb{Z} : (f(n), Q_x) = 1\}, \quad Q_x = \prod_{p \leq x} p.$$

Gaps: Let $J_f(x)$ be the largest gap in $\mathcal{S}_f(x)$.

Covering problem: Let $I_p = \{n \bmod p : f(n) \equiv 0 \pmod{p}\}$.
 $J_f(x)$ is the largest y so that $[0, y]$ is covered by

$$\{b_p + \nu_p \bmod p : p \leq x, \nu_p \in I_p\}$$

for some residues $b_p \bmod p$.

Difficulty: For a set p of positive density, $I_p = \emptyset$ (unused primes).
For $f(n) = n^2 + 1$, $I_p = \emptyset$ for $p \equiv 3 \pmod{4}$.

This means that Step 1 in the usual method for large prime gaps (the smooth number estimate) cannot be used. Without it, the other steps give only the trivial bound $J_f(x) \gg x$.

New estimate for $J_f(x)$

Theorem (FKMPT, 2018+)

Let $f : \mathbb{Z} \rightarrow \mathbb{Z}$ be a monic, irreducible polynomial with no fixed prime factor. Then $J_f(x) \gg x(\log x)^c$, where c depends on f .

New estimate for $J_f(x)$

Theorem (FKMPT, 2018+)

Let $f : \mathbb{Z} \rightarrow \mathbb{Z}$ be a monic, irreducible polynomial with no fixed prime factor. Then $J_f(x) \gg x(\log x)^c$, where c depends on f .

Corollary (FKMPT)

Let $f : \mathbb{Z} \rightarrow \mathbb{Z}$ be a non-constant polynomial. Then $\exists G_f \geq 2$ such that for any $k \geq G_f$ there are infinitely many integers $n \geq 0$ so that none of $f(n+1), \dots, f(n+k)$ is coprime to all the others.

Previously, this was known only for degree ≤ 3 .

New estimate for $J_f(x)$

Theorem (FKMPT, 2018+)

Let $f : \mathbb{Z} \rightarrow \mathbb{Z}$ be a monic, irreducible polynomial with no fixed prime factor. Then $J_f(x) \gg x(\log x)^c$, where c depends on f .

Corollary (FKMPT)

Let $f : \mathbb{Z} \rightarrow \mathbb{Z}$ be a non-constant polynomial. Then $\exists G_f \geq 2$ such that for any $k \geq G_f$ there are infinitely many integers $n \geq 0$ so that none of $f(n+1), \dots, f(n+k)$ is coprime to all the others.

Previously, this was known only for degree ≤ 3 .

Proof of Corollary. WLOG $f \in \mathbb{Z}[x]$, irreducible. If $x \in \mathbb{N}$ is large, then $J_f(x) \geq 2x + 1$. Let $k = 2x$ or $k = 2x + 1$. Then \mathbb{N} has infinitely many strings of k consecutive numbers, each having $p|f(n)$ for some $p \leq x$. But $p \leq k/2$, so $p|f(n \pm p)$ also, and one of $n \pm p$ is in the same interval.

Conjectures for $J_f(x)$

Theorem (FKMPT, 2018+)

Let $f : \mathbb{Z} \rightarrow \mathbb{Z}$ be a monic, irreducible polynomial with no fixed prime factor. Then $J_f(x) \gg x(\log x)^c$, where c depends on f .

Conjecture: $J_f(x) = x(\log x)^{1+o(1)}$.

(based on considering $S_f(x)$ as a random subset of $[1, Q_x]$)

Conjectures for $J_f(x)$

Theorem (FKMPT, 2018+)

Let $f : \mathbb{Z} \rightarrow \mathbb{Z}$ be a monic, irreducible polynomial with no fixed prime factor. Then $J_f(x) \gg x(\log x)^c$, where c depends on f .

Conjecture: $J_f(x) = x(\log x)^{1+o(1)}$.

(based on considering $S_f(x)$ as a random subset of $[1, Q_x]$)

Recall:

Conjecture (Maier-Pomerance, 1990). $J(x) = x(\log x)^{2+o(1)}$.

Why the difference? The smooth number bound gives an “arithmetic boost” to $J(x)$, but not to $J_f(x)$.

Method for showing that $J_f(x)$ is large, I

Main input: $|I_p|$ is 1 on average (Prime Ideal Theorem).



$$y = x(\log x)^c, z = x/(\log x)^{1/100}$$

Want $\{b_p + \nu_p \bmod p : p \leq x, \nu_p \in I_p\} \supseteq [0, y]$

- 1 (Random) Pick b_p at random, $p \leq z$ (uniformly, independently)
- 2 (Random-Greedy) For $z < p \leq \frac{x}{2}$, choose b_p at random, but only from “rich” residue classes (those covering many uncovered numbers from Step 1. Dependent on Step 1, non-uniform.
- 3 (Trivial) Same as prime case. Use $b_p \bmod p$ for $\frac{x}{2} < p \leq x, |I_p| \geq 1$ to cover anything left over ($\gg \frac{x}{\log x}$ such primes).

Method for showing that $J_f(x)$ is large, II

$$y = x(\log x)^c, \quad z = \frac{x}{(\log x)^{1/100}}$$

- 1 (Random) Pick b_p at random, $p \leq z$ (uniformly, independently)
- 2 (Random-Greedy) For $z < p \leq \frac{x}{2}$, take b_p from rich classes

Heuristic for Step 2:

- For fixed $q \in (z, x/2]$, let

$$S_1(r, q) = [0, y] \cap (r \pmod q) \setminus \bigcup_{p \leq (y/q)^{100}} \bigcup_{\nu_p \in I_p} (b_p + \nu_p \pmod p),$$

$$S_2(r, q) = [0, y] \cap (r \pmod q) \setminus \bigcup_{p \leq z} \bigcup_{\nu_p \in I_p} (b_p + \nu_p \pmod p).$$

There are many r for which $S_2(r, q) = S_1(r, q)$;
("rich" residue classes.)

- Sieving by primes $< (y/q)^{100}$ always leaves $\asymp \frac{y}{q \log(y/q)}$ elements.

Open Problems

I. Select a residue $a_p \in \mathbb{Z}/p\mathbb{Z}$ for each $p \leq x$, let

$$\mathcal{S} = [0, x] \setminus \bigcup_{p \leq x} (a_p \pmod p).$$

$\mathcal{S} = \emptyset$ possible: $a_2 = 1$, $a_p = 0$ ($3 \leq p \leq \frac{x}{2}$), a_p for $\frac{x}{2} < p \leq x$ cover $\{1, 2, 2^2, \dots\}$

Problem: What is the largest possible $|\mathcal{S}|$?

- A random choice yields $|\mathcal{S}| \sim e^{-\gamma} \frac{x}{\log x}$.
 - Any choice leaves $|\mathcal{S}| \ll \frac{x}{\log x}$ (sieve).
-

Open Problems

I. Select a residue $a_p \in \mathbb{Z}/p\mathbb{Z}$ for each $p \leq x$, let

$$S = [0, x] \setminus \bigcup_{p \leq x} (a_p \pmod p).$$

$S = \emptyset$ possible: $a_2 = 1$, $a_p = 0$ ($3 \leq p \leq \frac{x}{2}$), a_p for $\frac{x}{2} < p \leq x$ cover $\{1, 2, 2^2, \dots\}$

Problem: What is the largest possible $|S|$?

- A random choice yields $|S| \sim e^{-\gamma} \frac{x}{\log x}$.
 - Any choice leaves $|S| \ll \frac{x}{\log x}$ (sieve).
-

II. For each prime $p \leq \sqrt{x}$, choose a residue $a_p \pmod p$, and let

$$S = [0, x] \setminus \bigcup_{p \leq \sqrt{x}} (a_p \pmod p).$$

I. When $a_p = 0$ for all p , $|S| \sim x / \log x$.

II. A random choice yields $|S| \sim x(2e^{-\gamma} / \log x)$.

Question. Are these the extreme cases?