

# MAT 200 Handout 1

March 3, 2016

## SOME COMMON TAUTOLOGIES

This is not the full list, only the most common ones.

- (1) Modus ponens  $(P \wedge (P \Rightarrow Q)) \Rightarrow Q$
- (2)  $((P \Rightarrow Q) \wedge (\neg Q)) \Rightarrow \neg P$
- (3)  $((P \Rightarrow Q) \wedge (Q \Rightarrow R)) \Rightarrow (P \Rightarrow R)$
- (4) De Morgan's laws:  $(\neg(P \vee Q)) \iff (\neg P) \wedge (\neg Q)$
- (5)  $(\neg(P \wedge Q)) \iff (\neg P) \vee (\neg Q)$
- (6) Contrapositive:  $(P \Rightarrow Q) \iff (\neg Q \Rightarrow \neg P)$
- (7)  $(P \vee Q) \wedge \neg Q \Rightarrow P$
- (8)  $\neg(P \Rightarrow Q) \iff (P \wedge \neg Q)$

## SOME METHODS OF PROOF

**Direct proof:** To prove  $P \Rightarrow Q$ , assume  $P$ ; derive  $Q$  (you are allowed to use that  $P$  is true when doing this). Then you can conclude that  $P \Rightarrow Q$  is true (without any assumptions!).

**Indirect proof, also known as proof by contradiction:** To prove that  $P$  is true, assume  $\neg P$ ; derive from this a contradiction (you are allowed to use that  $P$  is false when doing this). Then you can conclude that  $P$  is true (without any assumptions!).

Important special case: to prove  $P \Rightarrow Q$ , you can assume  $\neg(P \Rightarrow Q)$  (which, by (8) above, is the same as  $P \wedge \neg Q$ ) and get a contradiction

**Proof by cases:** If  $P_1 \vee P_2 \vee \dots$  is true, and you have proved that  $P_1 \Rightarrow Q$ ,  $P_2 \Rightarrow Q$ ,  $\dots$ , then you can conclude that  $Q$  is true.

## INDUCTION

**Proofs by induction:** if  $P(n)$  is some statement depending on an integer number  $n$ , and we have checked that

- (Induction base) For  $n = 1$ ,  $P(n)$  is true
- (Induction step) For any  $n$ , if  $P(n)$  is true then  $P(n + 1)$  is also true.

Then  $P(n)$  is true for all integer  $n \geq 1$ .

**Strong induction** if  $P(n)$  is some statement depending on an integer number  $n$ , and we have checked that

- (Induction base) For  $n = 1$ ,  $P(n)$  is true
- (Induction step) For any  $n$ , if  $P(k)$  is true for  $k = 1, \dots, n$ , then  $P(n + 1)$  is also true.

Then  $P(n)$  is true for all integer  $n \geq 1$ .

## PROOFS WITH QUANTIFIERS

- (1) De Morgans laws:  $\neg(\exists x P(x)) \iff \forall x \neg P(x)$
- (2)  $\neg(\forall x P(x)) \iff \exists x \neg P(x)$
- (3) Given that  $\forall x \in M P(x)$  (where  $P(x)$  is some statement) and that  $c \in M$ , we can conclude  $P(c)$ .
- (4) If you know that  $\exists x P(x)$  is true, you can say "let us choose an  $a$  such that  $P(a)$  is true". Warning: to denote this chosen value, you should use a variable which was not used so far in your arguments!
- (5) To prove  $\exists x P(x)$ , it suffices to give one example of  $x$  for which  $P(x)$  is true.

- <sup>2</sup> (6) To prove  $\forall x P(x)$ , you have to give a proof of  $P(x)$  which would work for all possible values of  $x$ . Proving it in one (or two, or five....) special cases is not enough. Thus, a typical proof of  $\forall x \in M P(x)$  begins: "Let  $x$  be an arbitrary element of  $M$ ..."
- To **disprove**  $\forall x P(x)$ , it suffices to give one example of  $x$  for which  $P(x)$  is false.