Trees in Dynamics and Probability

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What do we see?



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The Julia set $\mathcal{J}(P)$ of $P(z) = z^2 + i$ (= set of points with bounded orbit under iteration).

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• a *dendrite*, i.e., a locally connected continuum with empty interior that does not separate the plane.

Follows from non-trivial, but standard facts in complex dynamics, because *P* is *postcritically-finite* and has no finite periodic critical points:

$$0 \longrightarrow i \longrightarrow -1 + \overbrace{i}^{i} - i$$

 a continuum tree, i.e., a locally connected, connected, compact metric space s.t. any two points can be joined by a unique arc.

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(Continuum) trees appear in various contexts:

- as Julia sets,
- as attractors of iterated function systems (e.g., the CSST=continuum self-similar tree),
- in probabilistic models (e.g., the CRT=continuum random tree).

CSST T (=continuum self-similar tree)





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$\mathsf{CSST}\ \mathbb{T}$



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∽ ९ (२) 10 / 26 The CSST ${\mathbb T}$ is:

- a geodesic continuum tree (as an abstract metric space).
- an attractor of an iterated function system (as a subset of the plane). Define

$$f_1(z) = rac{1}{2}z - rac{i}{2}, \quad f_2(z) = -rac{1}{2}ar{z} + rac{i}{2}, \quad f_3(z) = rac{i}{2}ar{z} + rac{1}{2},$$

Then $\mathbb{T}\subseteq\mathbb{C}$ is the unique non-empty compact set satisfying

$$\mathbb{T} = f_1(\mathbb{T}) \cup f_2(\mathbb{T}) \cup f_3(\mathbb{T}).$$

So \mathbb{T} is the attractor of the iterated function system $\{f_1, f_2, f_3\}$ in the plane.

Theorem (B.-Huy Tran 2018; folklore)

A continuum tree T is homeomorphic to the CCST \mathbb{T} iff all branch points of T have order 3 and they are dense in T.

Proof: \Rightarrow : Looks obvious, but is somewhat involved if one defines \mathbb{T} as an attractor of an iterated function system.

 $\Leftarrow: For each level <math>n \in \mathbb{N}$ carefully cut T into 3^n pieces. Label pieces by words in a finite alphabet to align with pieces of \mathbb{T} . Use general lemma to obtain a homeomorphism.

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Lemma providing a homeomorphism

Let (X, d_X) and (Y, d_Y) be compact metric spaces. Suppose that for each level $n \in \mathbb{N}$, the space X admits a decomposition $X = \bigcup_{i=1}^{M_n} X_{n,i}$ as a finite union of non-empty compact subsets $X_{n,i}$ with the following properties:

(i) Each set $X_{n+1,j}$ is the subset of some set $X_{n,j}$.

(ii) Each set $X_{n,i}$ is equal to the union of some of the sets $X_{n+1,j}$. (iii) $\max_{1 \le i \le M_n} \operatorname{diam}(X_{n,i}) \to 0$ as $n \to \infty$.

Suppose that for $n \in \mathbb{N}$ the space Y admits similar decompositions $Y = \bigcup_{i=1}^{M_n} Y_{n,i}$ with properties analogous to (i)–(iii) such that

$$X_{n+1,j} \subseteq X_{n,i}$$
 if and only if $Y_{n+1,j} \subseteq Y_{n,i}$ (1)

and

$$X_{n,i} \cap X_{n,j} \neq \emptyset$$
 if and only if $Y_{n,i} \cap Y_{n,j} \neq \emptyset$ (2)

for all *n*, *i*, *j*.

Then there exists a unique homeomorphism $f: X \to Y$ such that $f(X_{n,i}) = Y_{n,i}$ for all n and i. In particular, the spaces X and Y are homeomorphic.

Each end-piece X is cut into three children by using a branch point $p \in X$ of largest weight w(p) in T.

Each arc piece X is cut into three children by using a branch point $p \in \alpha$ that is contained in the distinguished arc α of X and has largest weight w(p) in T among these branch points.

 $w(p) = \text{diameter of third largest component of } T \setminus \{p\}.$

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The CRT is a random geodesic continuum tree $T = T(\omega)$ constructed from Brownian excursion.

If $e = e(\omega) \colon [0,1] \to [0,\infty)$ is a sample of Brownian excursion, define

$$d(s,t) = e(s) + e(t) - 2\min_{u \in [s,t]} e(u).$$

Then $T = [0, 1] / \sim$, where

$$s \sim t : \Leftrightarrow d(s, t) = 0$$

 $\Leftrightarrow e(s) = e(t) = \min_{u \in [s,t]} e(u).$

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Are two independent samples of the CRT almost surely homeomorphic?

Theorem (B.-Tran 2018)

The CRT is almost surely homeomorphic to the CSST.

Actually, this theorem was essentially proved earlier by Croyden and Hambly in 2008.

For the previous and other results on the topology of trees see: M. Bonk and Huy Tran, *The continuum self-similar tree*, Preprint, March 2018, on arXiv.

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The map f is quasisymmetric (=qs) iff there exists a distortion function $\eta: [0, \infty) \rightarrow [0, \infty)$ s.t.

$$\frac{d_{Y}(f(x), f(y))}{d_{Y}(f(x), f(z))} \leq \eta \left(\frac{d_{X}(x, y)}{d_{X}(x, z)}\right),$$

whenever x, y, z are distinct points in X. If there exists a quasisymmetry $f \colon X \to Y$, then X and Y are said to be *qs-equivalent*. Let (X, d_X) and (Y, d_Y) be metric spaces, and $f: X \to Y$ a homeomorphism.

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Geometry of a quasisymmetric map



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- *f* is quasisymmetric if it maps balls to "roundish" sets of uniformly controlled "eccentricity".
- Quasisymmetry global version of quasiconformality.
- bi-Lipschitz \Rightarrow qs \Rightarrow qc.
- In ℝⁿ, n ≥ 2: qs ⇔ qc.
 Also true for "Loewner spaces" (Heinonen-Koskela).

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- When can one promote a homeomorphism between trees to a quasisymmetry?
- Is there a characterization of the CSST up to qs-equivalence?
- What can one say about the qc-geometry of the CRT or Julia sets of postcritically-finite polynomials?
- Are there some canonical models for certain classes of trees up to qs-equivalence (uniformization problem)?

Theorem (Tukia-Väisälä 1980)

Let α be a metric arc. Then α is qs-equivalent to [0,1] iff α is doubling and of bounded turning.

A metric space X is *doubling* is there exists $N \in \mathbb{N}$ such that every ball in X can be covered by N (or fewer) balls of half the radius.

A metric space (X, d) is of *bounded turning* if there exists K > 0such that for all points $x, y \in X$ there exists a continuum γ with $x, y \in \gamma$ and

diam $\gamma \leq Kd(x, y)$.

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Let T be a quasi-tree, i.e., a tree that is doubling and of bounded turning. Then T is qs-equivalent to a geodesic tree.

A tree (T, d) is *geodesic* if for all $x, y \in T$,

 $\operatorname{length}[x,y] = d(x,y).$

Note: in a quasi-tree we have diam $[x, y] \simeq d(x, y)$.

Proof of Theorem: Decompose the metric space (T, d) into pieces, and carefully redefine metric d by assigning new diameters to these pieces to obtain a geodesic metric ϱ on T. With suitable choices, the identity map $(T, d) \rightarrow (T, \varrho)$ is a quasisymmetry.

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Proof: Carefully cut T into pieces with good geometric control and align with pieces of CCST. Geometric control comes from conformal elevator techniques.

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Proof: Carefully cut T into pieces with good geometric control and align with pieces of CCST. Geometric control comes from conformal elevator techniques.

A better proof would follow from a qs-characterization of the CSST (open problem!).

Let P be a postcritically-finite polynomial. If its Julia set $\mathcal{J}(P)$ is homeomorphic to the CSST, then it is qs-equivalent to the CSST.

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Proof: Same ideas.

Quasiconformal geometry of the CRT

Question

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Open Problem (very hard!)

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General theme: Geometric uniqueness of probabilistic models. Known (up to quasi-isometric equivalence) for Bernoulli percolation on \mathbb{Z} , Poisson point process on \mathbb{R} (Basu-Sly).