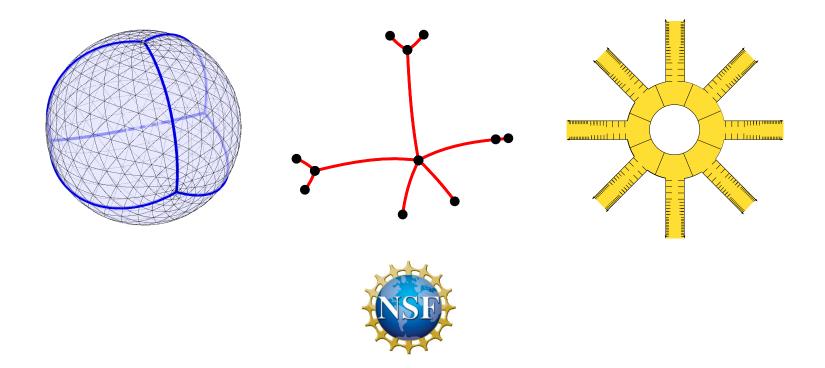
DESSINS AND DYNAMICS Christopher Bishop, Stony Brook

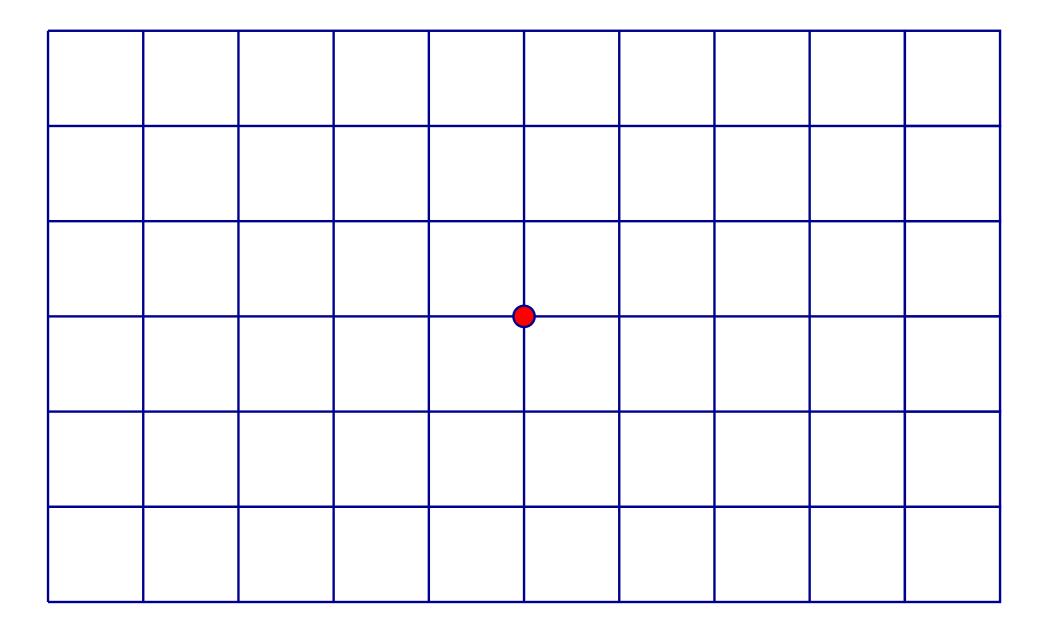
Groups and Dynamics Seminar, Texas A&M 3-4pm, Wednesday February 23, 2022

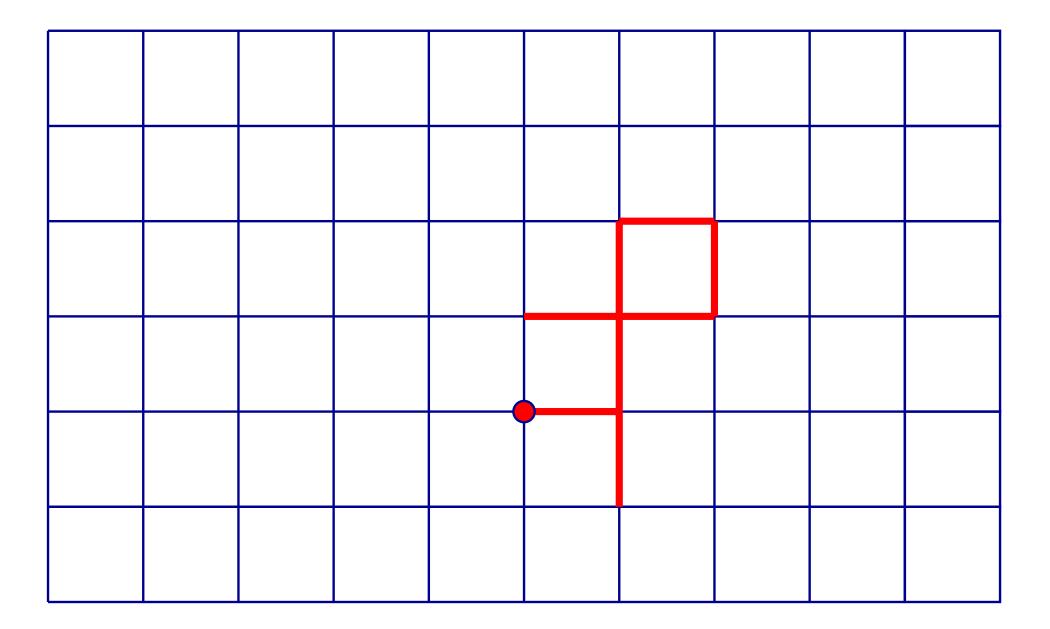
www.math.sunysb.edu/~bishop/lectures

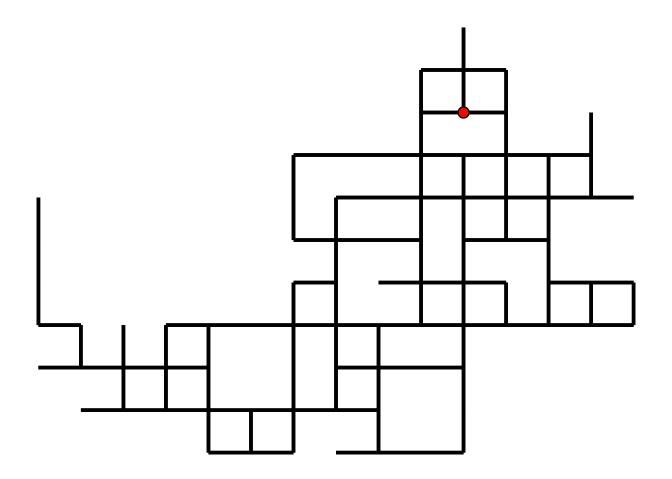


THE PLAN

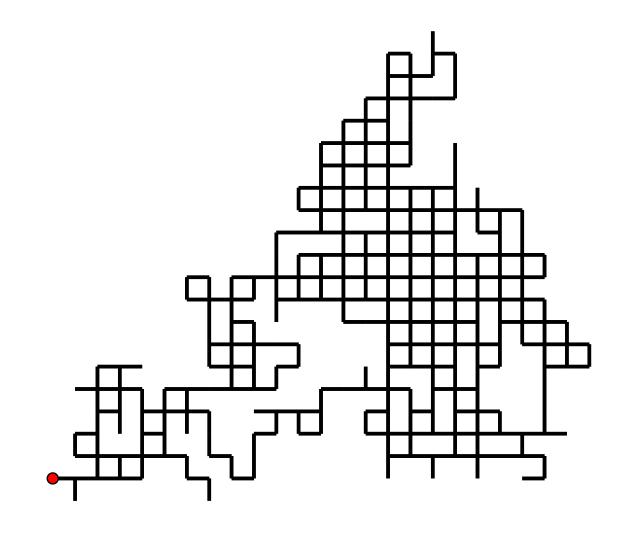
- Harmonic measure
- Finite trees and Shabat polynomials
- Infinite trees and entire functions
- Wandering domains, dimensions of Julia sets
- Equilateral triangulations and critical orbits



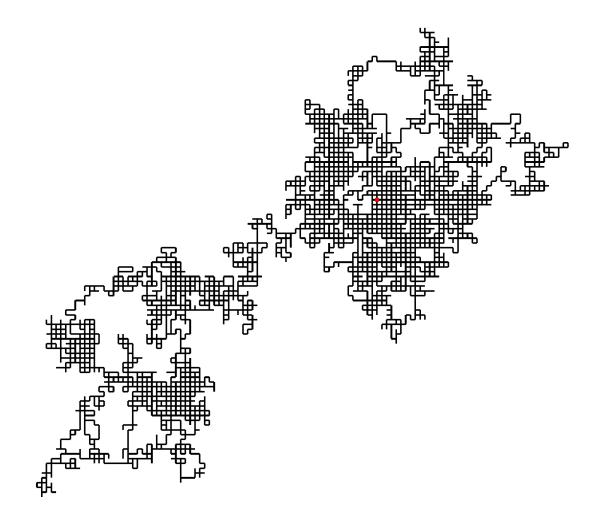




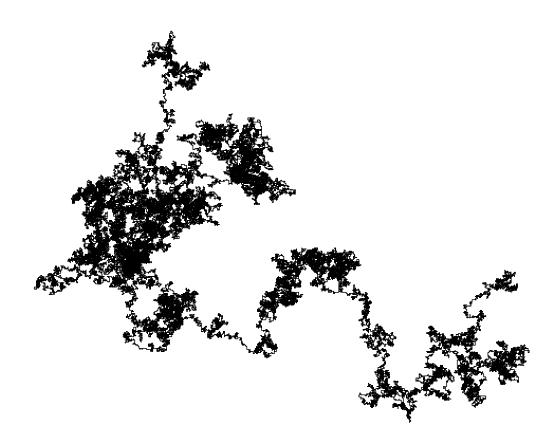
200 step random walk.



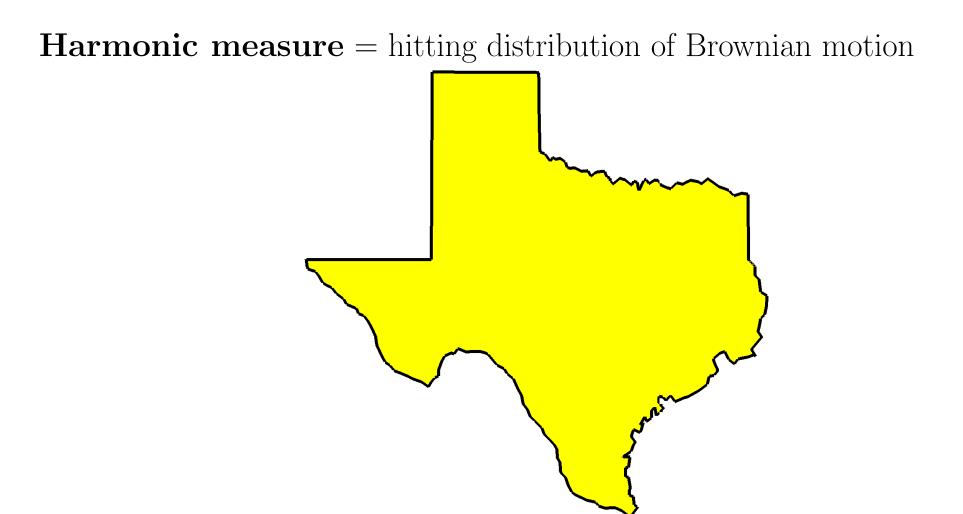
1000 step random walk.



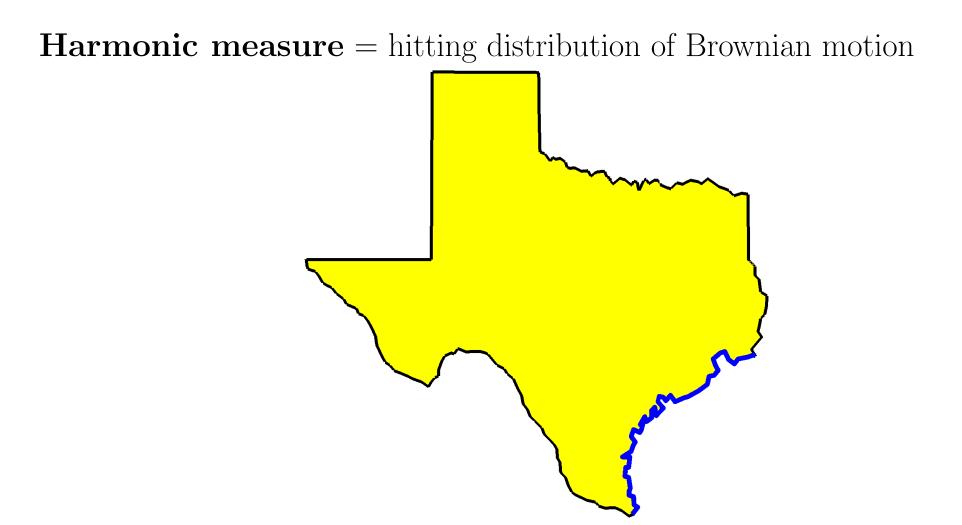
10,000 step random walk.



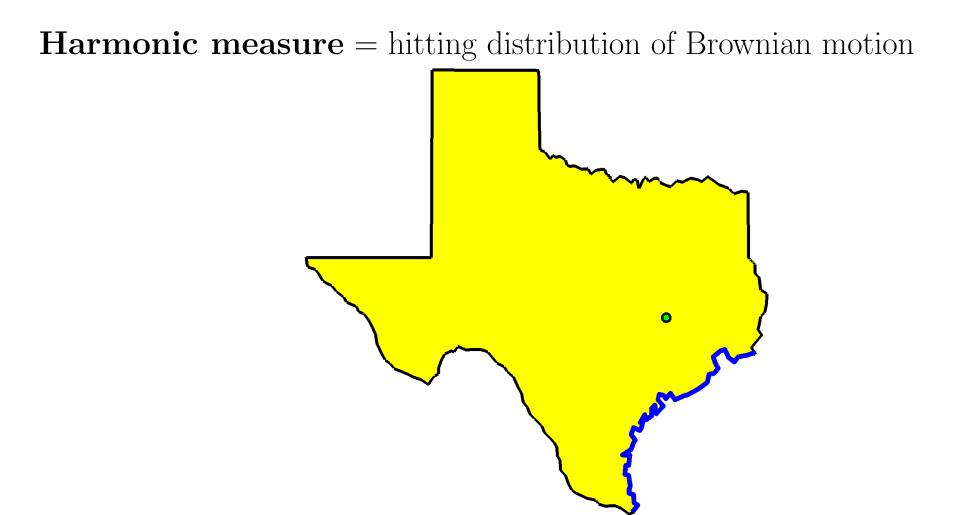
100,000 step random walk.



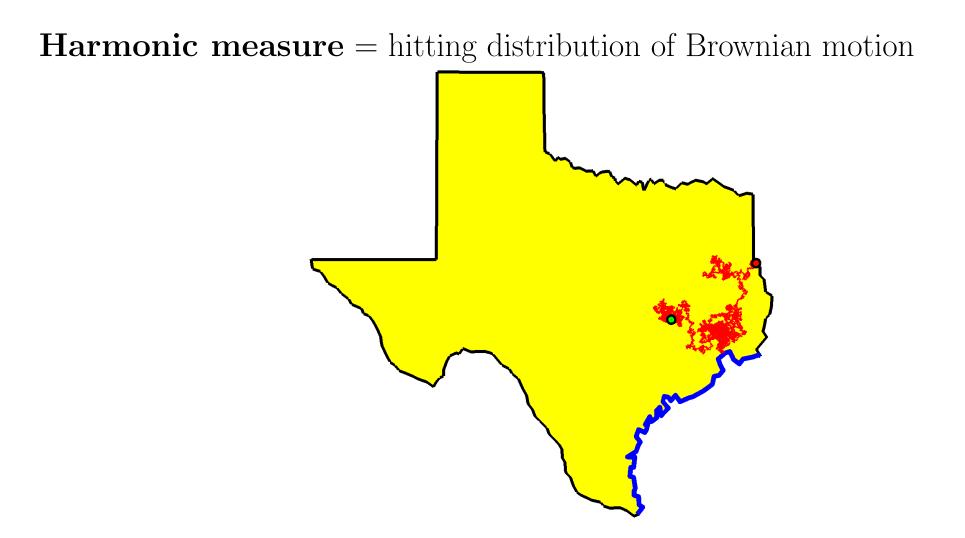
Suppose Ω is a planar Jordan domain.

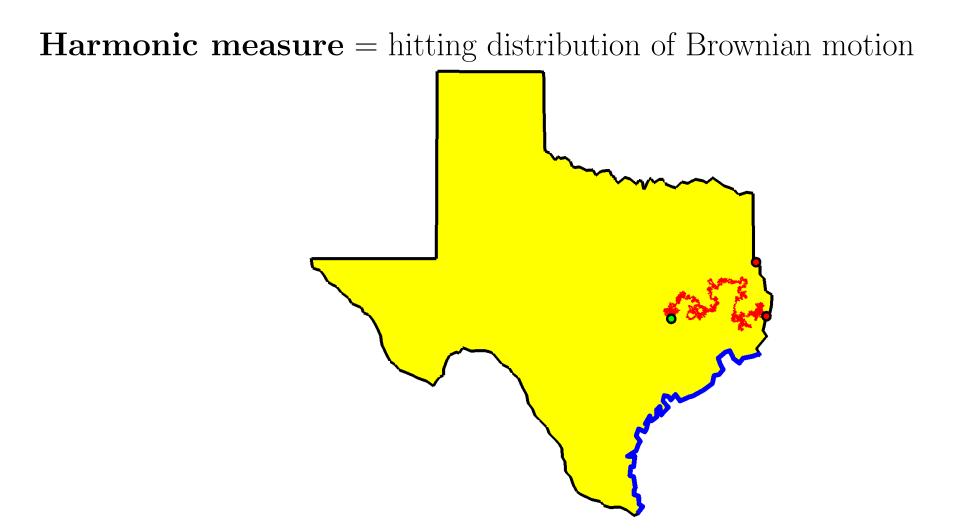


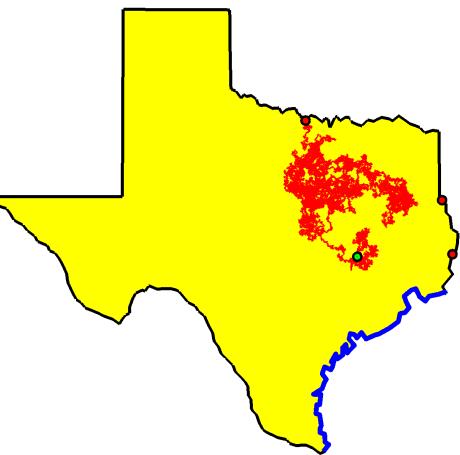
Let E be a subset of the boundary, $\partial \Omega$.

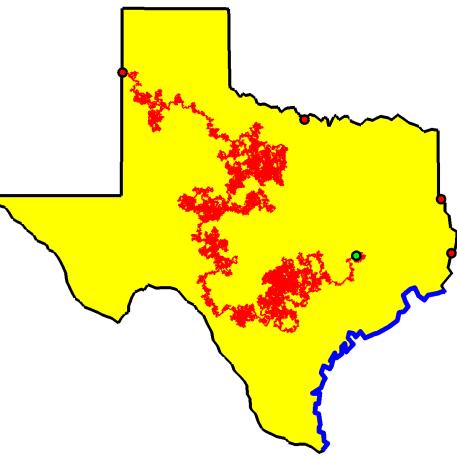


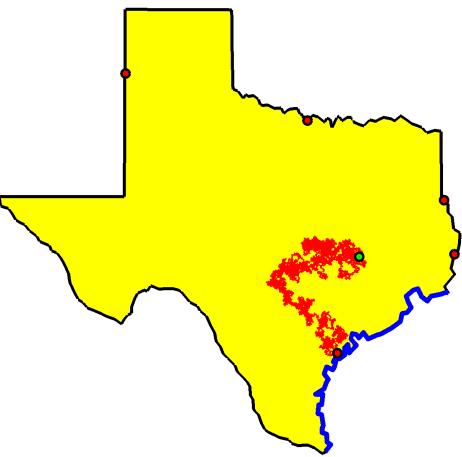
Choose a point z inside Ω .

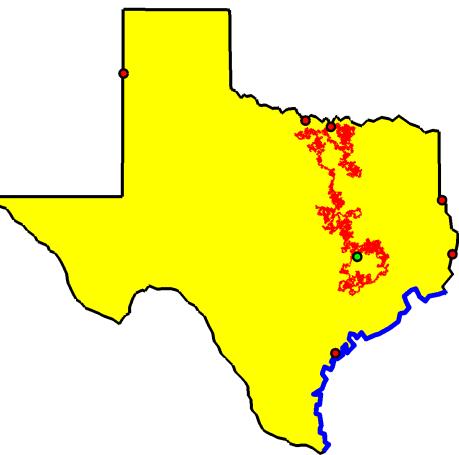


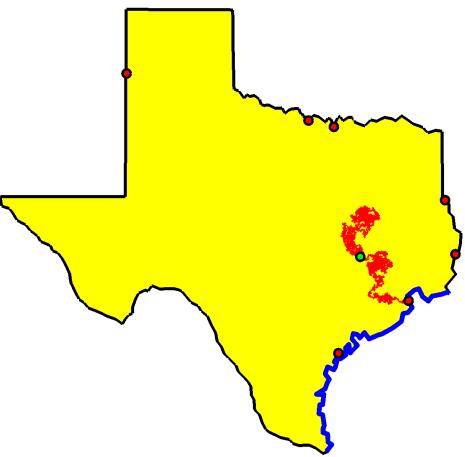




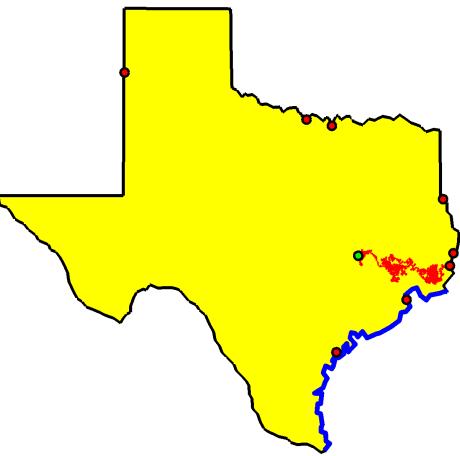


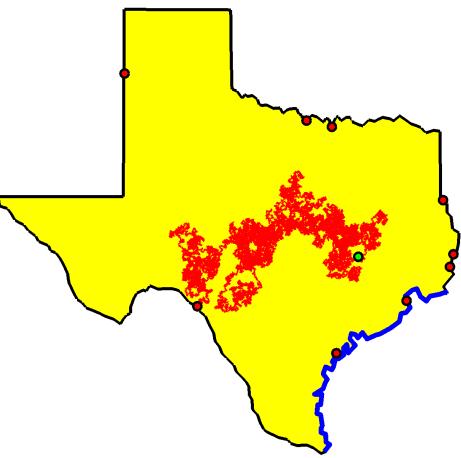


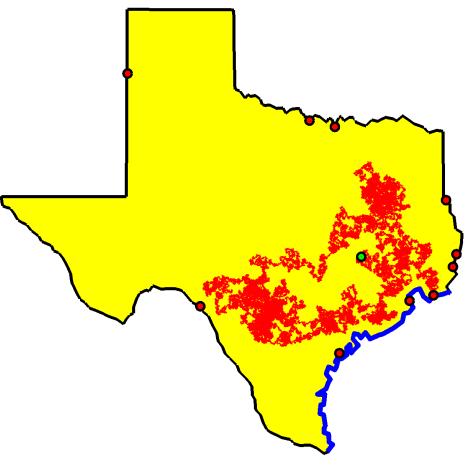








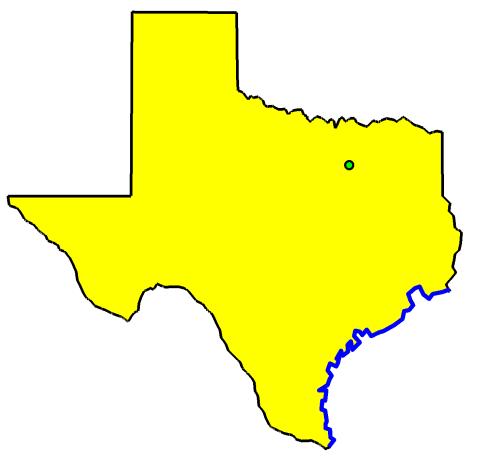




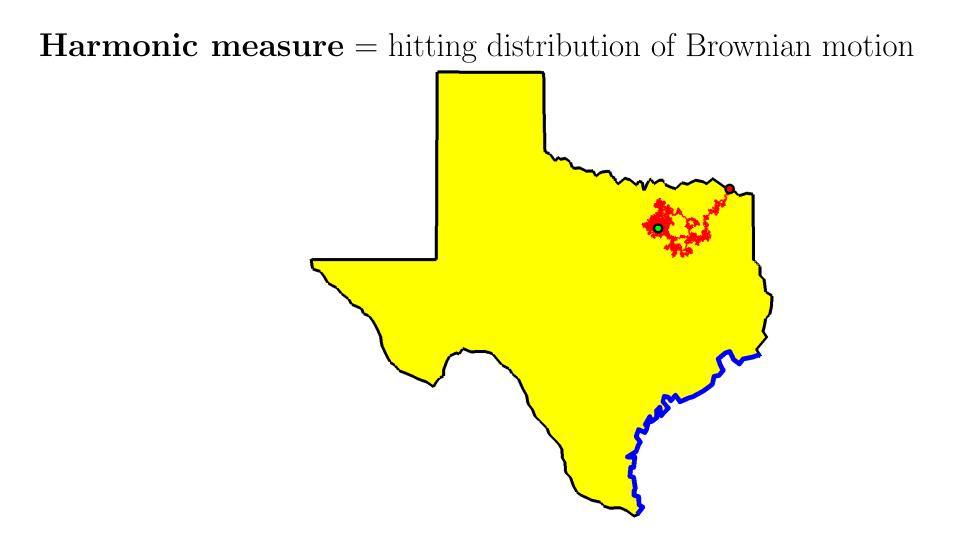
 $\omega(z,E,\Omega) \thickapprox 3/10$

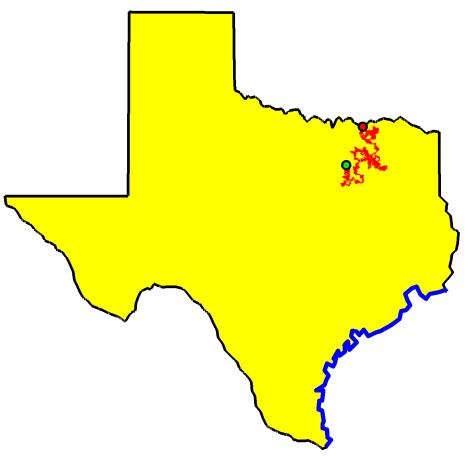


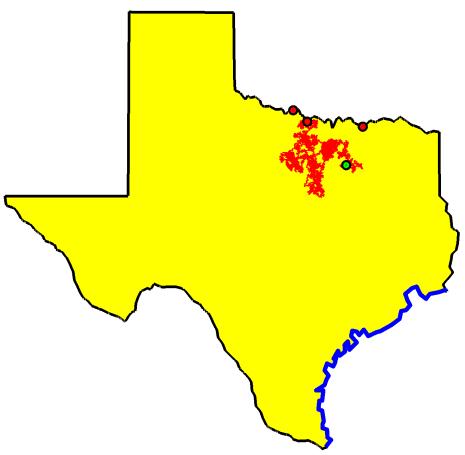
 $\omega(z, E, \Omega) \approx 42/100.$

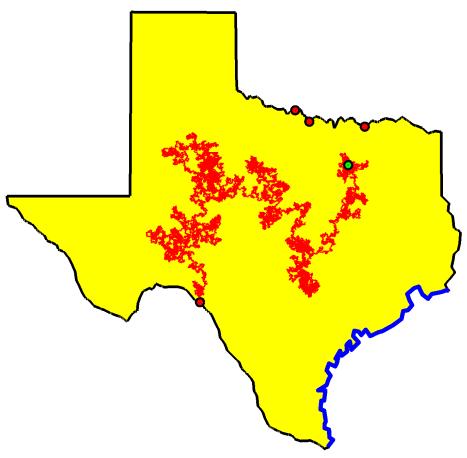


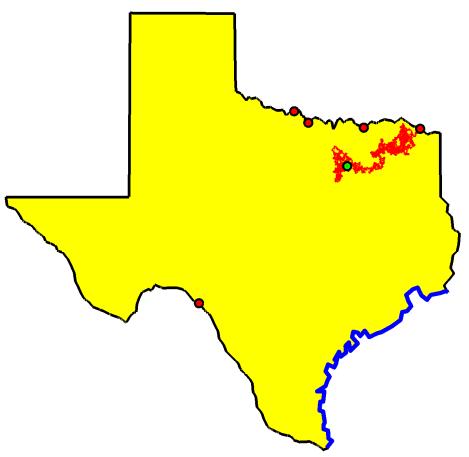
Choose a new starting point $z \in \Omega$.

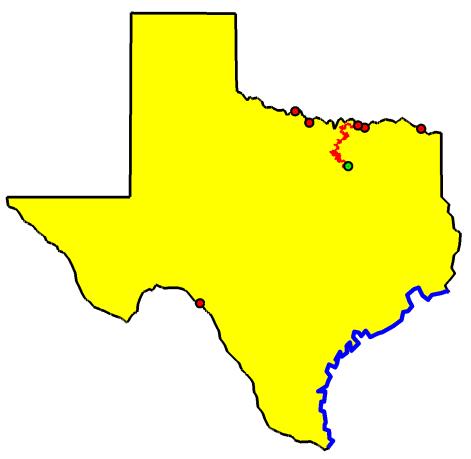


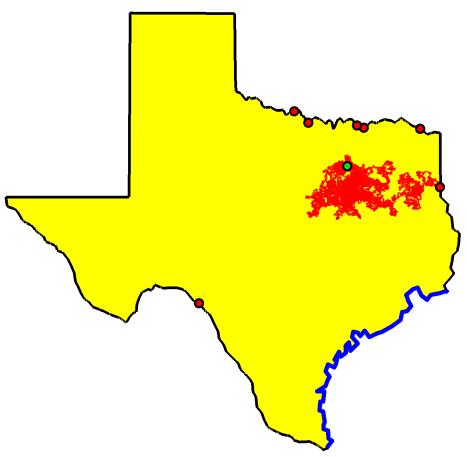


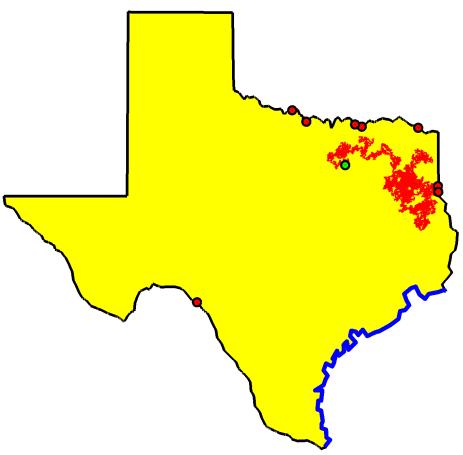


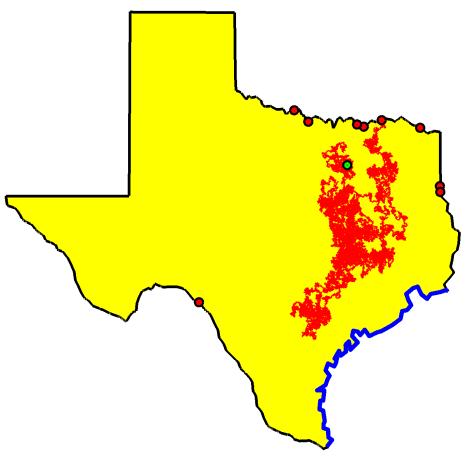


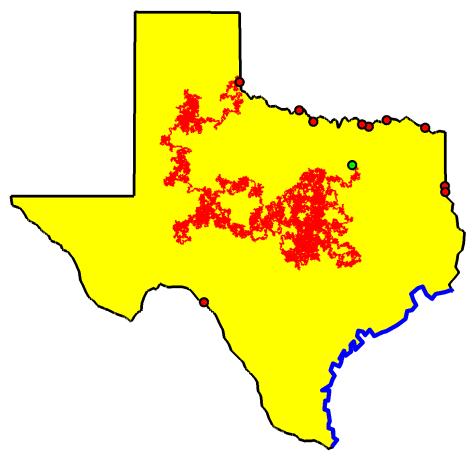








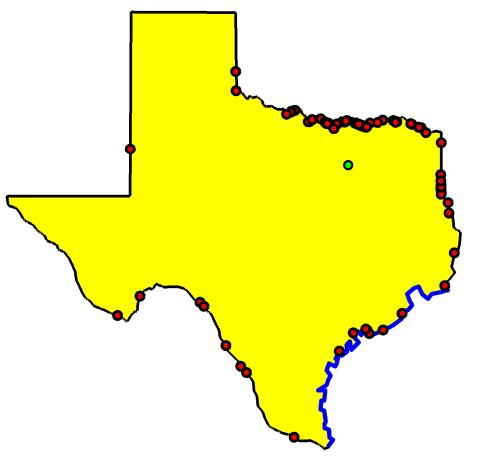




 $\omega(z,E,\Omega) \thickapprox 0/10$

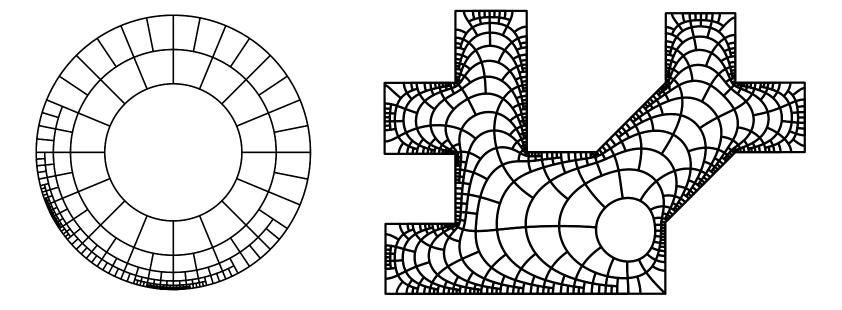


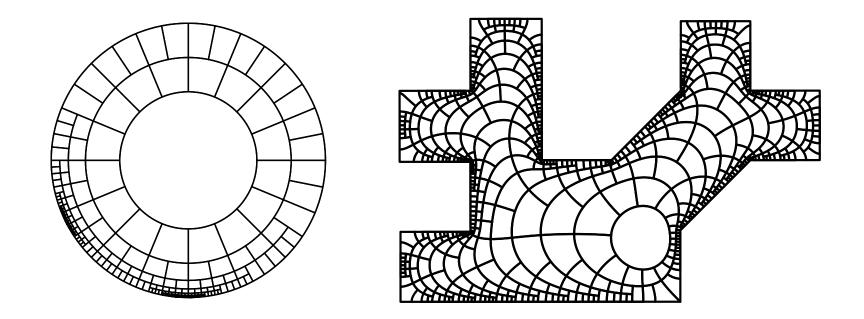
 $\omega(z, E, \Omega) \approx 8/100.$



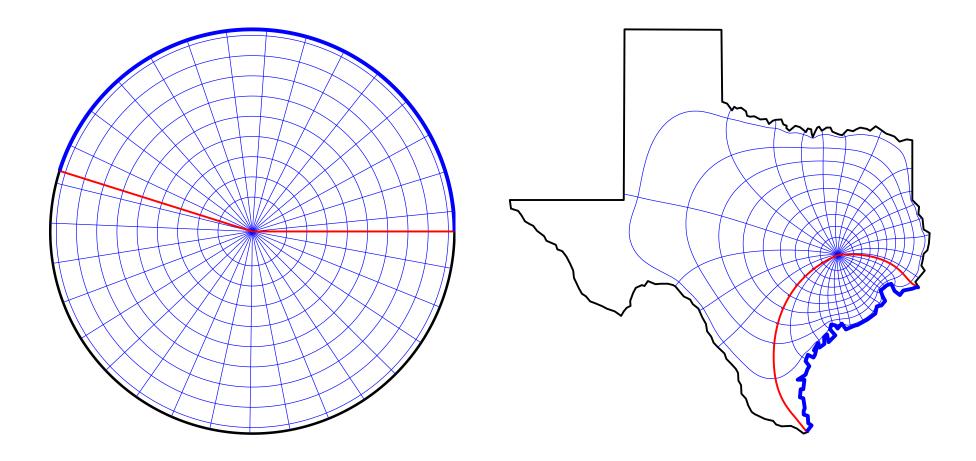
 $\omega(z, E, \Omega)$ is a harmonic function of z on Ω . Boundary values are 1 on E and 0 elsewhere. Easy to solve this PDE on the disk: Poisson kernel. **Riemann Mapping Theorem:** If $\Omega \subsetneq \mathbb{R}^2$ is simply connected and $z \in \Omega$, then there is a conformal map $f : \mathbb{D} \to \Omega$ with f(0) = z.

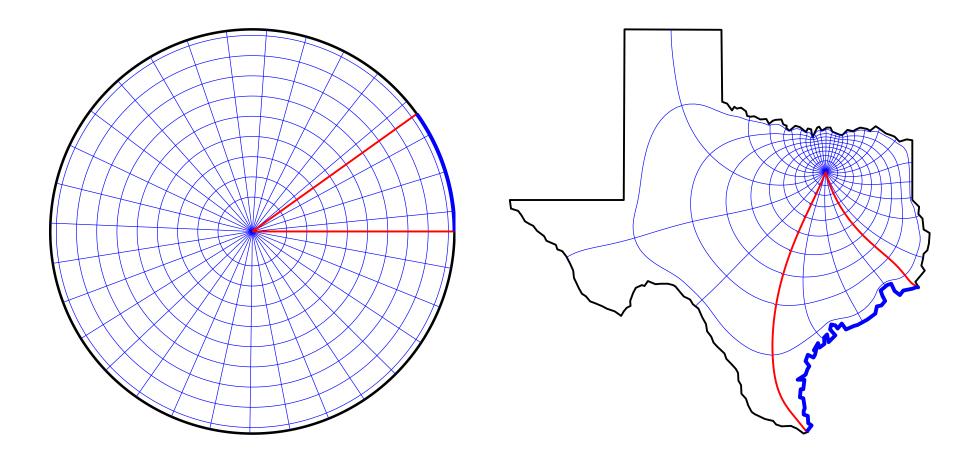
conformal = 1-1, holomorphic = angle and orientation preserving

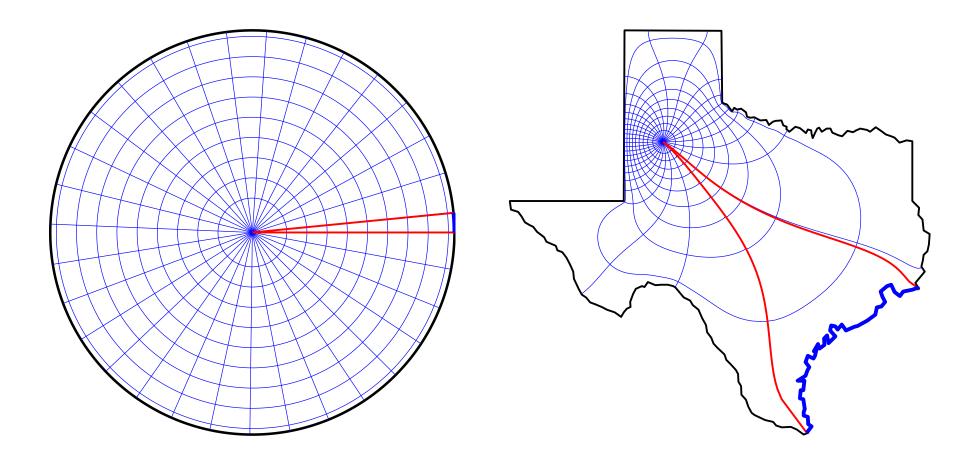


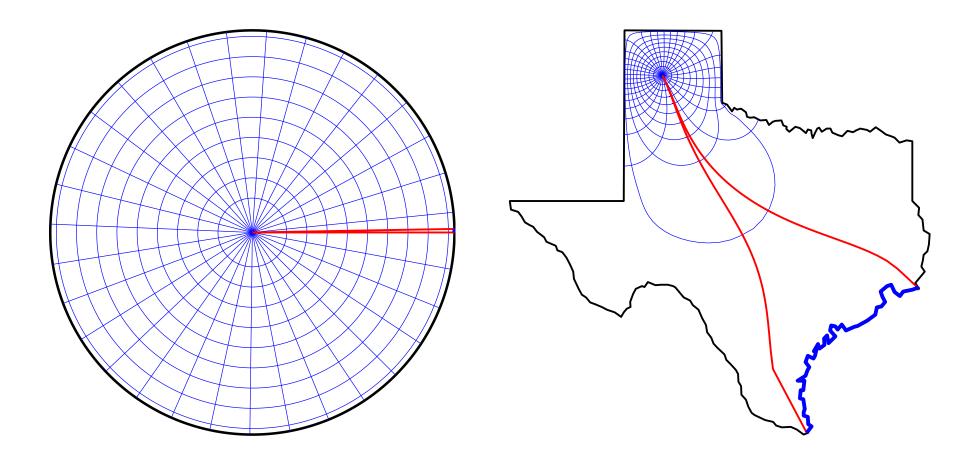


If ω is harmonic on Ω , then $\omega \circ f$ is harmonic on \mathbb{D} . $\Rightarrow \omega(z, E, \Omega) = \omega(0, f^{-1}(E), \mathbb{D}) = |E|/2\pi.$ $\Rightarrow \omega$ on $\partial\Omega$ = conformal image of (normalized) length on \mathbb{T}





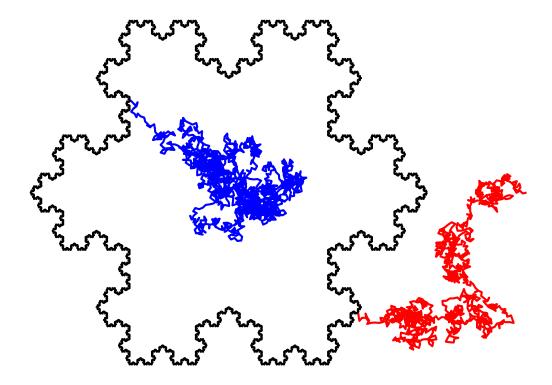




- different base points $z \in \Omega$ give different distributions on $\partial \Omega$,
- but all points in Ω give mutually continuous distributions.

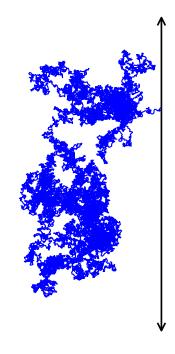
Both can fail for points on different sides of $\partial \Omega$.

For a fractal curve, inside and outside harmonic measures are singular.



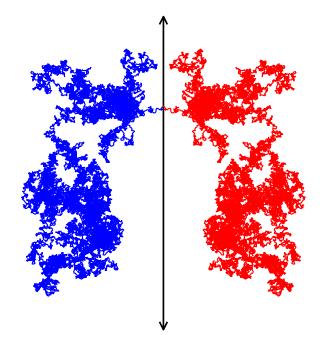
 $\omega_1 \perp \omega_2$ iff tangents points have zero length. (B. 1987, based on results of N.G. Makarov)

For which curves is $\omega_1 = \omega_2$? (inside measure = outside measure)



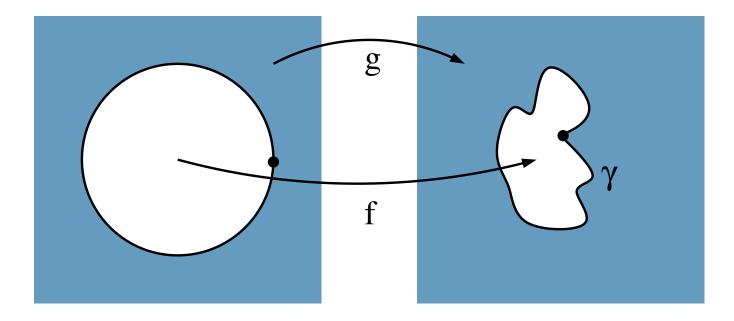
True for lines and circles.

For which curves is $\omega_1 = \omega_2$? (inside measure = outside measure)



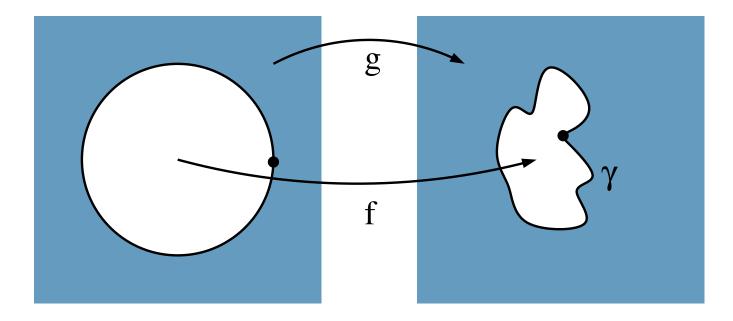
True for lines and circles.

Converse is also true.



Suppose $\omega_1 = \omega_2$ for a curve γ .

Conformally map two sides of circle to two sides of γ so f(1) = g(1). $\omega_1 = \omega_2$ implies maps agree on whole boundary.



Suppose $\omega_1 = \omega_2$ for a curve γ .

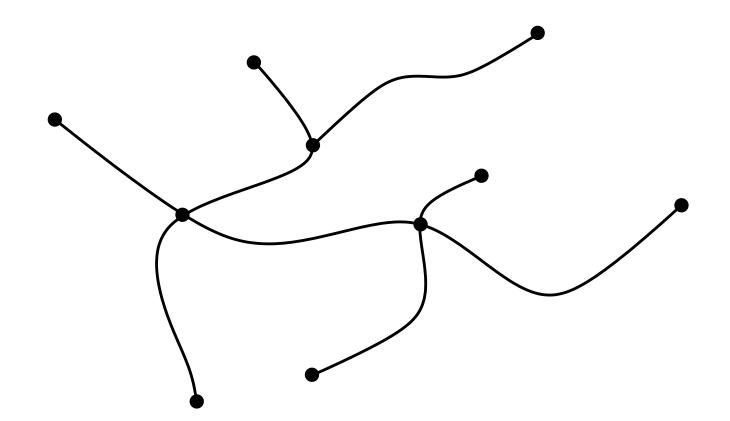
Conformally map two sides of circle to two sides of γ so f(1) = g(1). $\omega_1 = \omega_2$ implies maps agree on whole boundary.

So f, g define homeomorphism h of plane holomorphic off circle. Then h is entire by Morera's theorem.

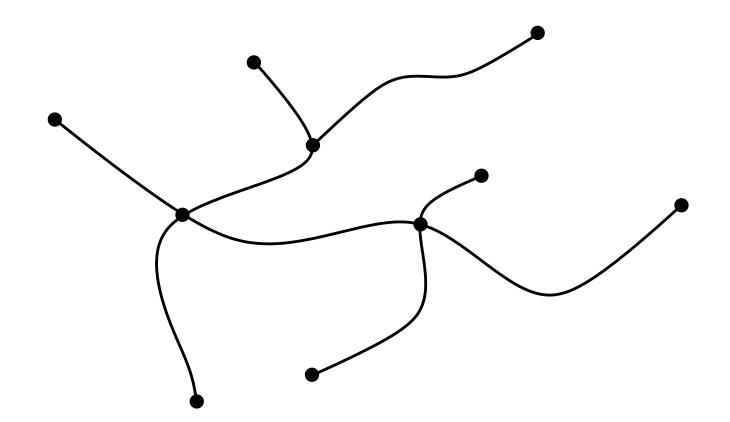
Entire and 1-1 implies h is linear (Liouville's thm), so γ is a circle.

Circles are only closed curves with equal harmonic measures on both sides.

What about other kinds of 2-sided objects?



A planar graph is a finite set of points connected by non-crossing edges. It is a tree if there are no closed loops.



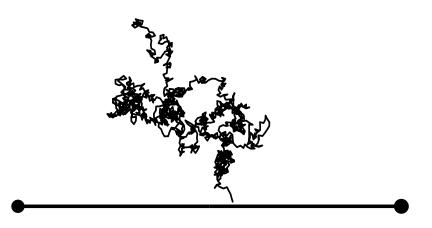
A planar tree is **conformally balanced** if

- \bullet every edge has equal harmonic measure from ∞
- edge subsets have same measure from both sides

This is also called a "**true tree**". A line segment is an example.

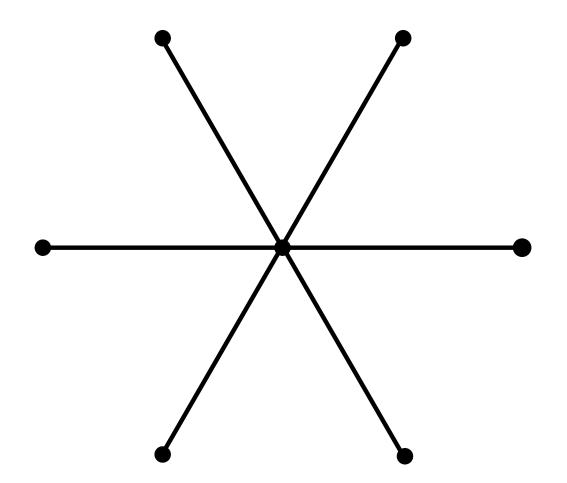


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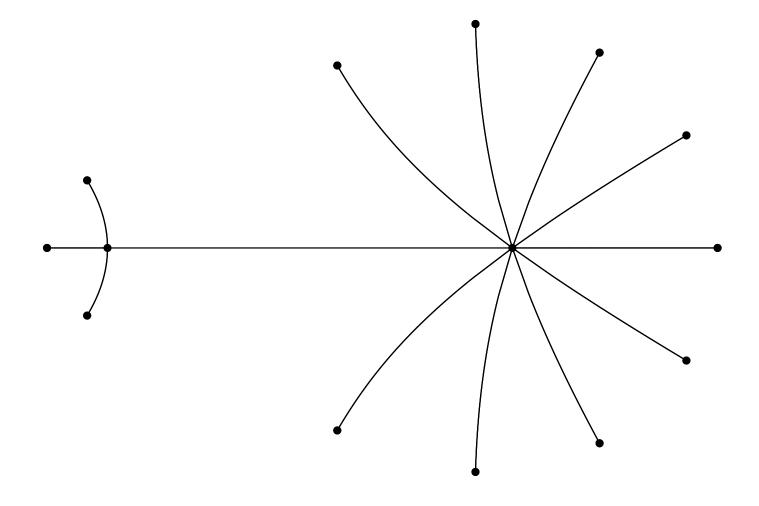


This is also called a "**true tree**". A line segment is an example.





Trivially true by symmetry



Non-obvious true tree

Definition of critical value: if p = polynomial, then

$$CV(p) = \{p(z) : p'(z) = 0\} = critical values$$

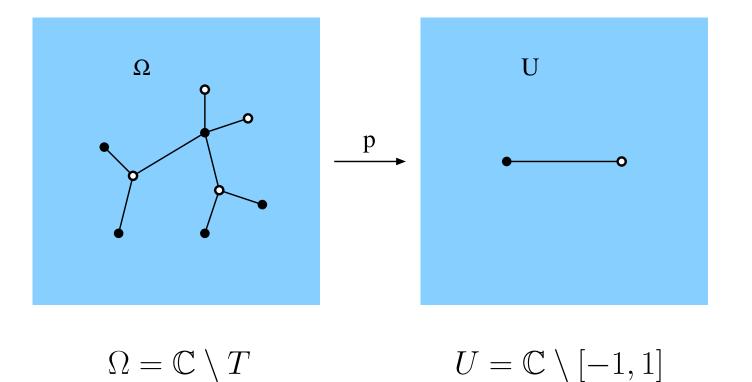
If $CV(p) = \pm 1$, p is called **generalized Chebyshev** or **Shabat**.

Definition of critical value: if p = polynomial, then

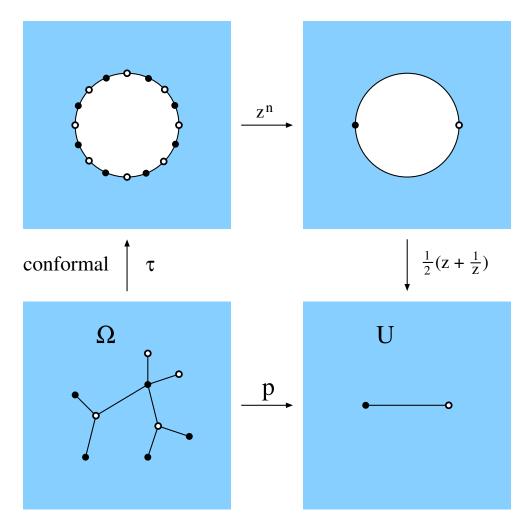
$$CV(p) = \{p(z) : p'(z) = 0\} = critical values$$

If $CV(p) = \pm 1$, p is called **generalized Chebyshev** or **Shabat**.

Thm: T is balanced iff $T = p^{-1}([-1, 1]), p =$ Shabat.



T conformally balanced $\Leftrightarrow p$ Shabat.

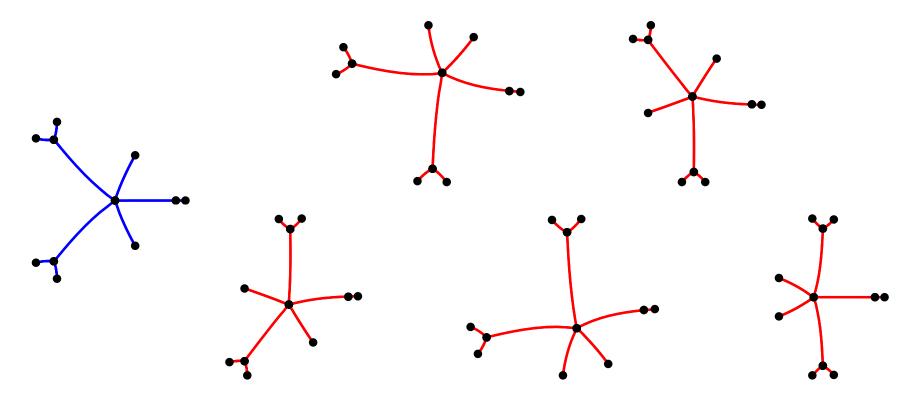


p is entire and n-to-1 $\Leftrightarrow p =$ polynomial. $CV(p) \notin U \leftrightarrow p : \Omega \to U$ is covering map.

Algebraic aside:

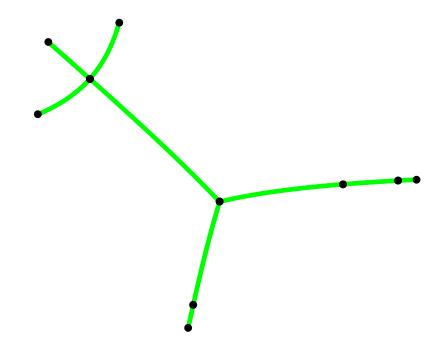
True trees are examples of Grothendieck's *dessins d'enfants* on sphere.

Normalized polynomials are algebraic, so trees correspond to number fields. Computing number field from tree is difficult.



Six graphs of type 5 1 1 1 1 1 - 3 3 2 1 1, two orbits.

Kochetkov (2009, 2014): cataloged all trees with 9 and 10 edges.



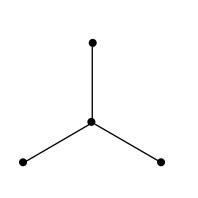
For example, the polynomial for this 9-edge tree is

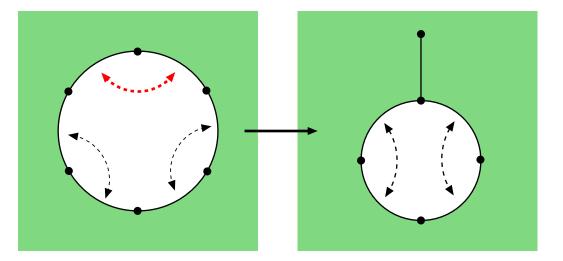
$$p(z) = z^4(z^2 + az + b)^2(z - 1),$$

where a is a root of ...

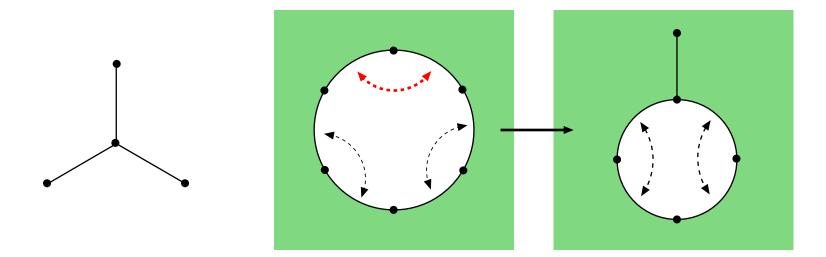
I did **not** use this to draw tree on previous slide.

Don Marshall's **ZIPPER** uses conformal mapping to draw true trees.





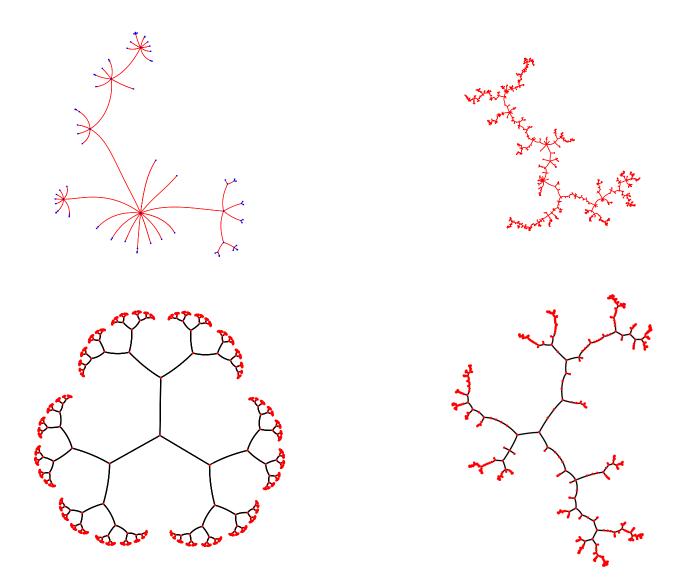
Don Marshall's **ZIPPER** uses conformal mapping to draw true trees.



Conformal mapping can handle trees with thousands of edges.

Can obtain polynomial roots α with thousands of digits of accuracy.

Marshall and Rohde compute number field by finding integer relation between $1, \alpha, \alpha^2, \ldots$ using Ferguson's PSLQ algorithm.



Some true trees, courtesy of Marshall and Rohde

Which planar trees have a true form?

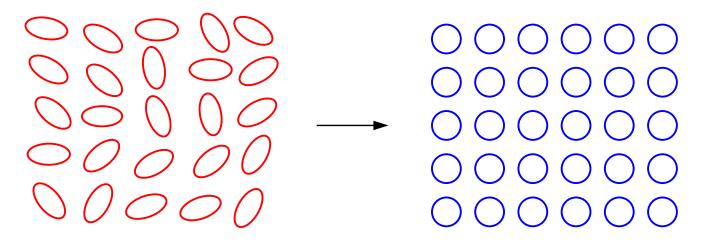
What can that true form look like?

Theorem: Every finite planar tree has a true form.

Unique up to similarities (Morera + Liouville).

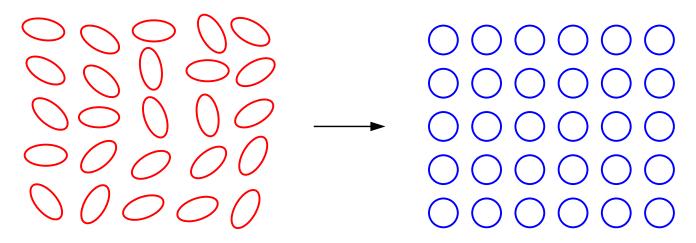
Standard proof uses the uniformization theorem.

I will describe an alternate proof using quasiconformal maps.



Eccentricity = ratio of major to minor axis of ellipse.

K-quasiconformal = ellipses have eccentricity $\leq K$ almost everywhere

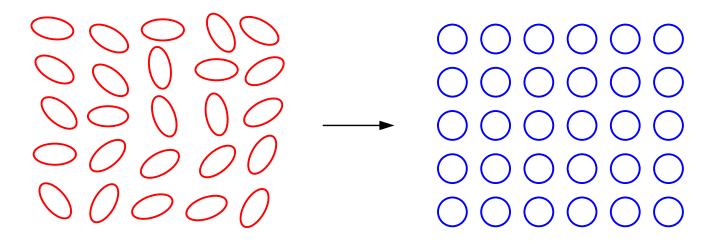


Eccentricity = ratio of major to minor axis of ellipse.

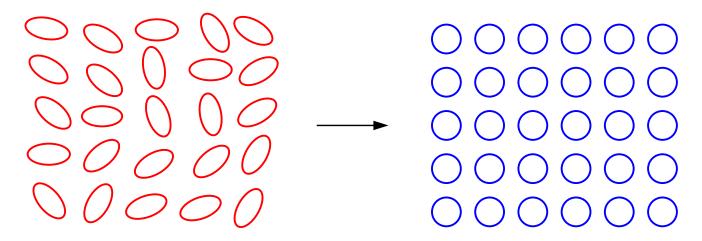
K-quasiconformal = ellipses have eccentricity $\leq K$ almost everywhere Ellipses determined a.e. by measurable dilatation $\mu = f_{\overline{z}}/f_z$ with K-1

$$|\mu| \le \frac{K - 1}{K + 1} < 1.$$

Conversely, \ldots

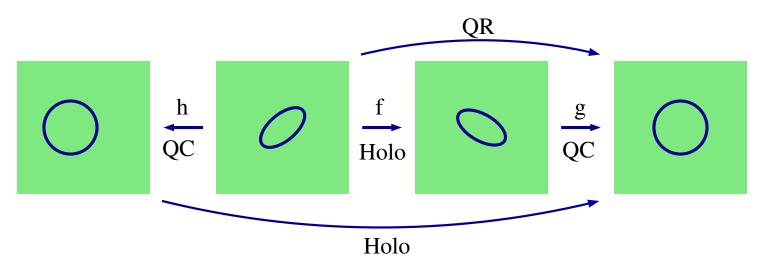


Mapping theorem: any such μ comes from some QC map f.



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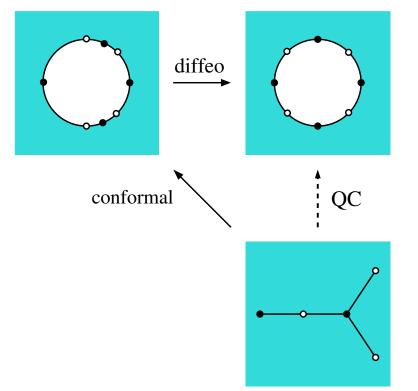
Cor: If f is holomorphic and g is QC, then there is a QC map h so that $F = g \circ f \circ h^{-1}$ is also holomorphic. $(g \circ f$ is called QR = quasiregular)



QC proof that every finite tree has a true form:

Map $\Omega = \mathbb{C} \setminus T$ to $\{|z| > 1\}$ conformally.

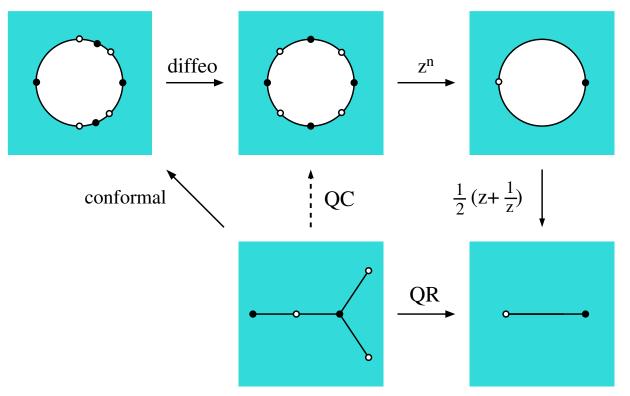
"Equalize intervals" by diffeomorphism. Composition is quasiconformal.



Dilatation of QC map depends degree of "imbalance" of harmonic measure.

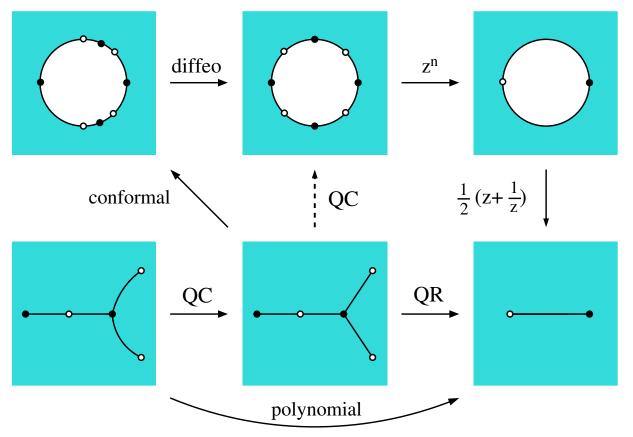
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"Equalize intervals" by diffeomorphism. Composition is quasiconformal.



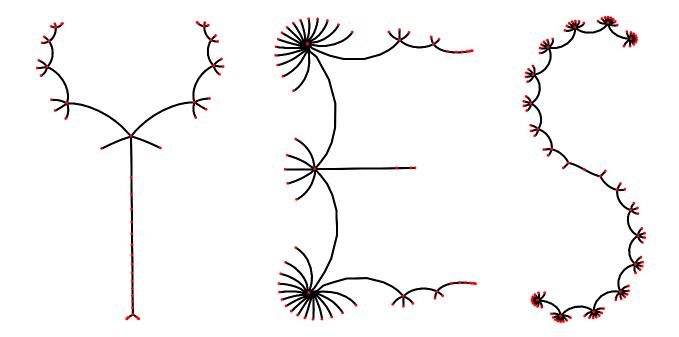
Mapping theorem implies there is a QC φ so $p = q \circ \varphi$ is a polynomial.

Thus true trees can have any combinatorics.

Alex Eremenko: can they have any shape?

Thus true trees can have any combinatorics.

Alex Eremenko: can they have any shape?



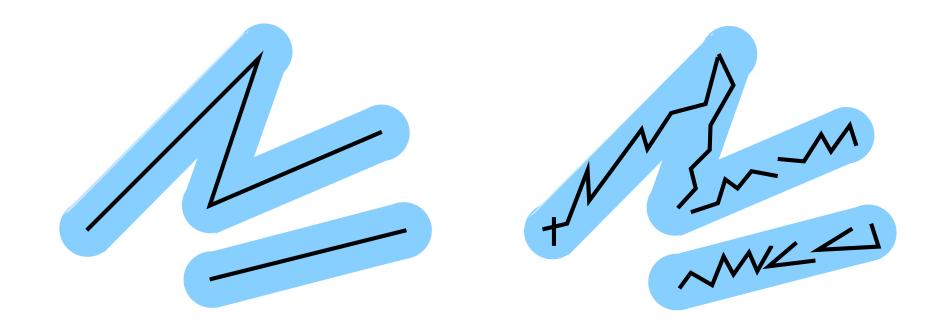
Thm: Every planar continuum is Hausdorff limit of true trees.

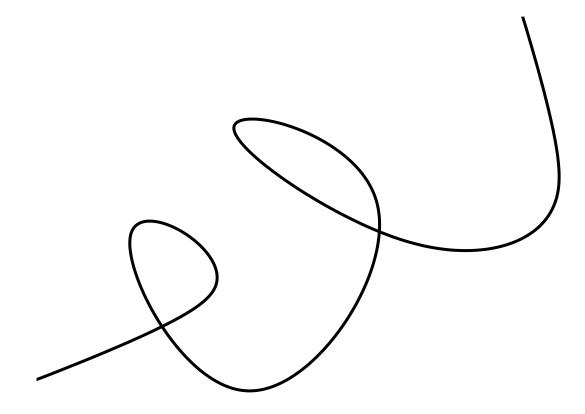
Cor (B.-Pilgrim): Every planar continuum is Hausdorff limit of Julia sets of PCF polynomials (PCF = all critical points have finite orbits).

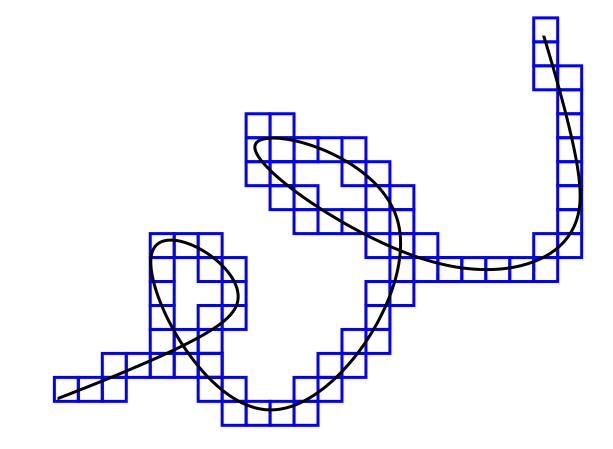
For a set E and $\epsilon > 0$ define ϵ -neighborhood of E: $E(\epsilon) = \{z : \operatorname{dist}(z, E) < \epsilon\}.$

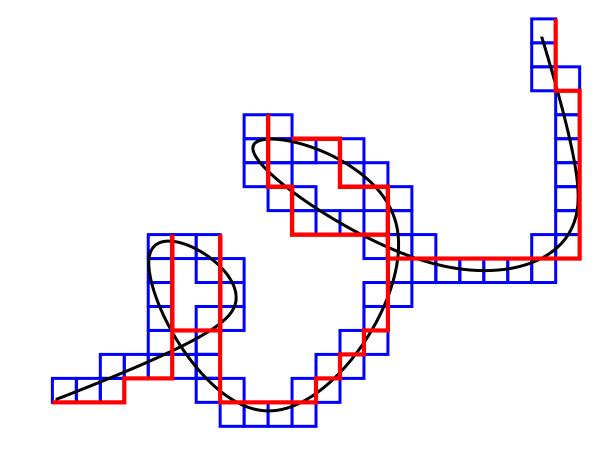
Defn: Hausdorff distance between sets is

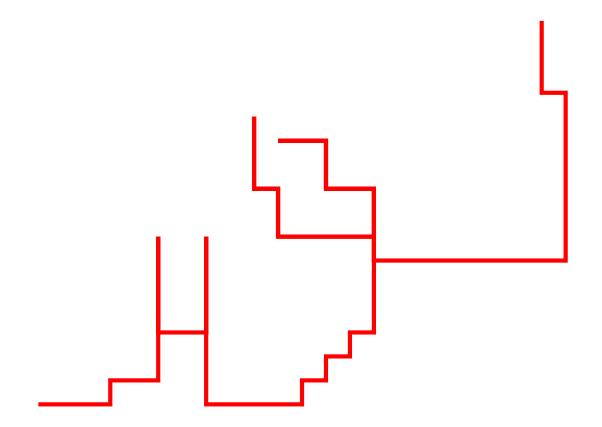
$$d(E, F) = \inf\{\epsilon > 0 : E \subset F(\epsilon), F \subset E(\epsilon)\}.$$





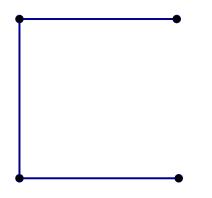






Theorem: Every planar continuum is a limit of true trees.

Idea of Proof: reduce harmonic measure ratio by adding edges.

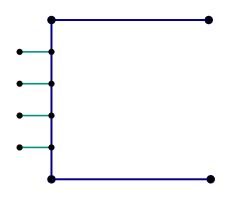


Vertical side has much larger harmonic measure from left.

Add edges (\Rightarrow change combinatorics) to "balance" harmonic measure.

Theorem: Every planar continuum is a limit of true trees.

Idea of Proof: reduce harmonic measure ratio by adding edges.

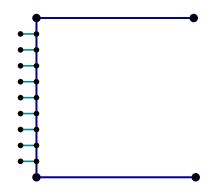


"Left" harmonic measure is reduced (roughly 3-to-1).

New edges are approximately balanced (universal constant).

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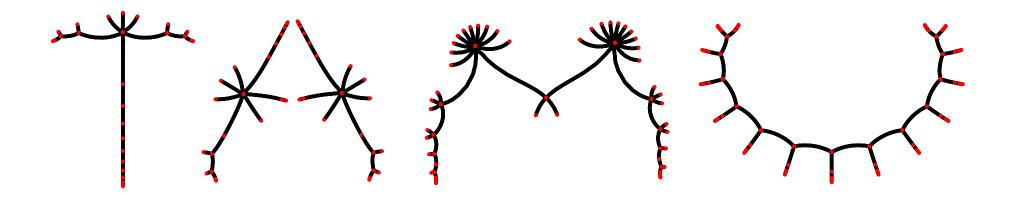
"Left" harmonic measure is reduced (roughly 3-to-1).

New edges are approximately balanced (universal constant).

Mapping theorem gives exactly balanced.

QC correction map is near identity if "spikes" are short.

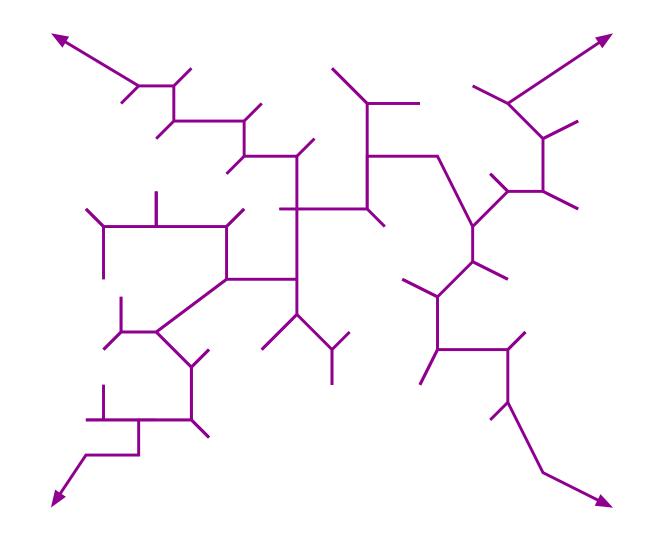
New tree approximates shape of old tree; different combinatorics.



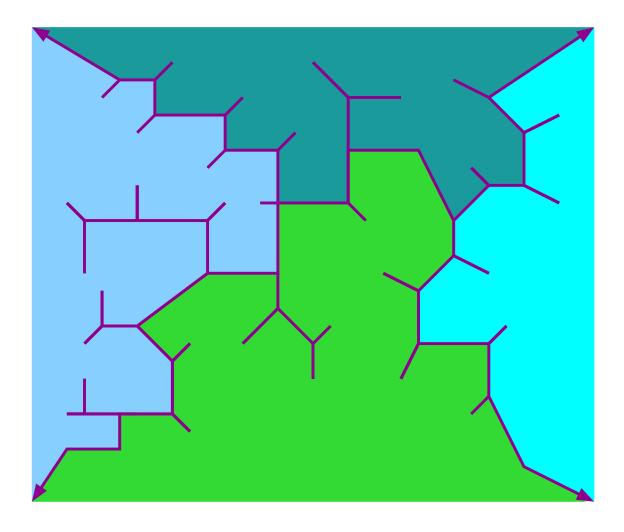
Dessins d'adolescents

Finite planar trees \Leftrightarrow polynomials with 2 critical values.

What about infinite planar trees? Shabat entire functions?



Do infinite trees correspond to entire functions with 2 critical values?

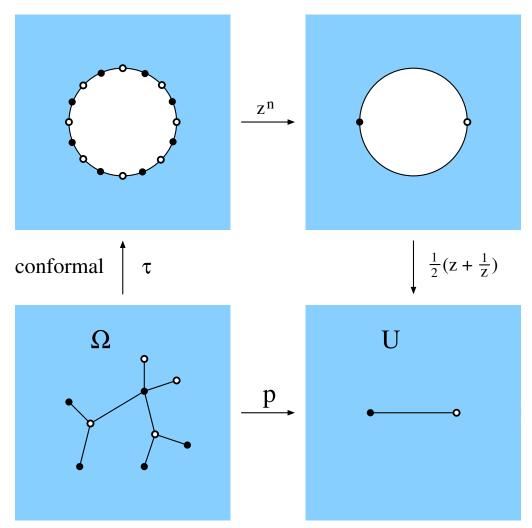


Main difference:

 $\mathbb{C}\setminus$ finite tree = one topological annulus

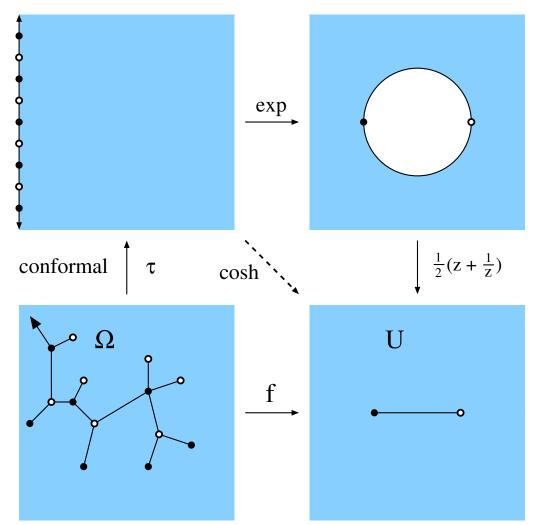
 $\mathbb{C}\setminus$ infinite tree = many simply connected components

Recall finite case



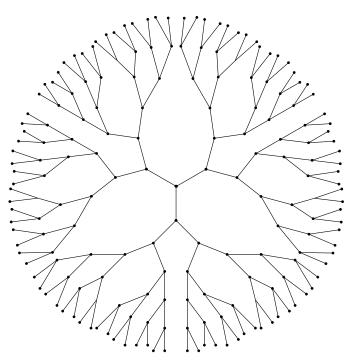
T is true tree $\Leftrightarrow p = \frac{1}{2}(\tau^n + 1/\tau^n)$ is continuous across T.

Infinite case



Infinite balanced tree $\Leftrightarrow f = \cosh \circ \tau$ is continuous across T.

Infinite 3-regular tree does not have a true form.



Sketch: Let G = ambient automorphisms of tree.

True form exists + Morera Thm + Liouville Thm \Rightarrow automorphisms extend to isometries of plane $\Rightarrow G$ is discrete subgroup of isometries of \mathbb{R}^2 G has exponential growth, isometry group of \mathbb{R}^2 doesn't. $\Rightarrow \Leftarrow$ Are there enough "infinite true trees" to approximate any shape?

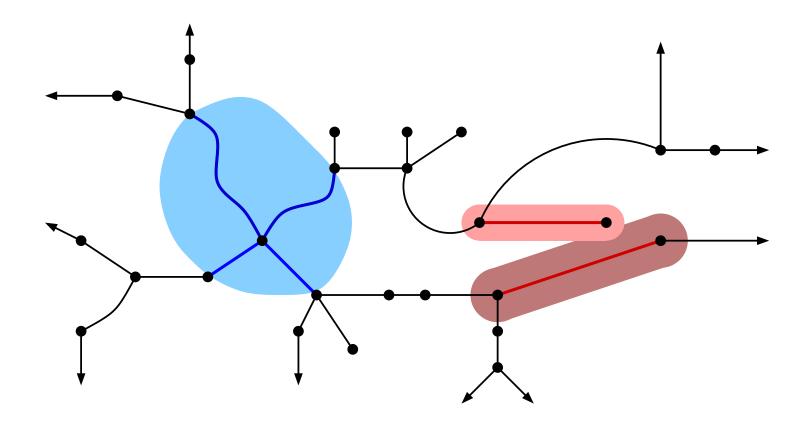
Thm: every "nice" infinite planar tree T is approximated by $f^{-1}([-1, 1])$ for some entire function f with $CV = \{\pm 1\}$.

"nice" = two conditions that are automatic for finite trees.

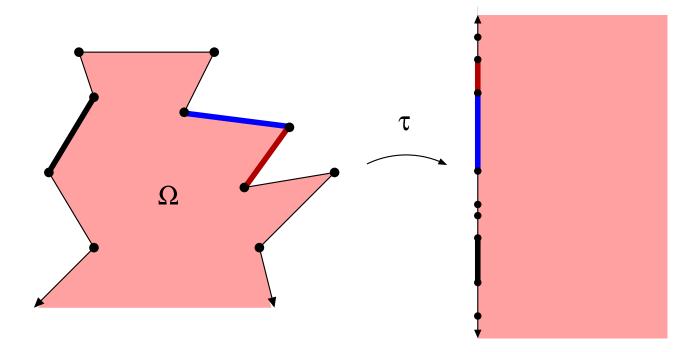
(1) Bounded Geometry (local condition; easy to verify):

- edges are uniformly smooth.
- adjacent edges form bi-Lipschitz image of a star = $\{z^n \in [0, r]\}$
- non-adjacent edges are well separated,

 $\operatorname{dist}(e, f) \geq \epsilon \cdot \min(\operatorname{diam}(e), \operatorname{diam}(f)).$



(2) *τ*-Lower Bound (global condition; harder to check):
Complementary components of tree are simply connected.
Each can be conformally mapped to right half-plane. Call map *τ*.



We assume all images have length $\geq \pi$.

Need positive lower bound; actual value usually not important. Components are "thinner" than half-plane near ∞ .

If e is an edge of T and r > 0 let

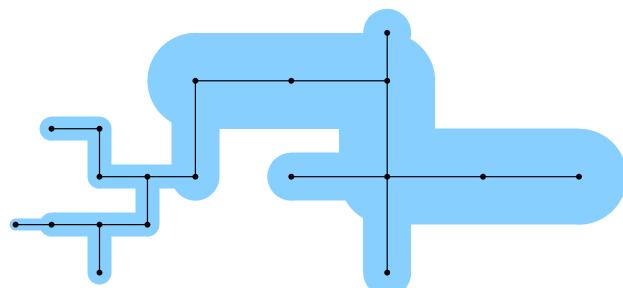
$$e(r) = \{z : \operatorname{dist}(z, e) \le r \cdot \operatorname{diam}(e)\}$$



If e is an edge of T and r > 0 let

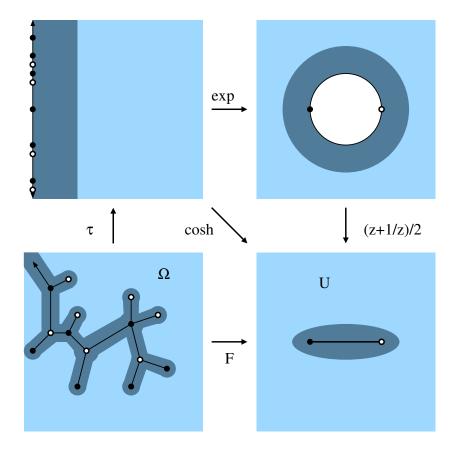
$$e(r) = \{z : \operatorname{dist}(z, e) \le r \cdot \operatorname{diam}(e)\}$$

Define neighborhood of $T: T(r) = \bigcup \{ e(r) : e \in T \}.$



T(r) for infinite tree replaces Hausdorff metric in finite case.

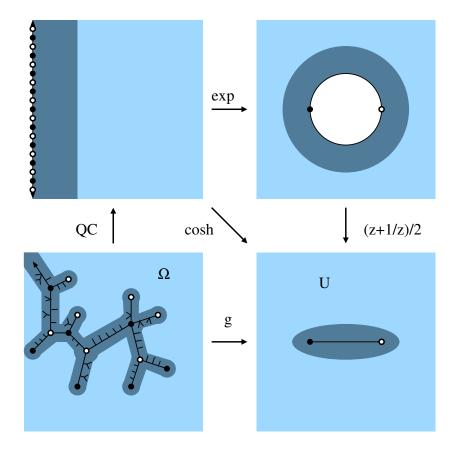
Thm: If T has bounded geometry and satisfies a τ -lower bound, then there is a QR g and r > 0 with $g = \cosh \circ \tau$ off T(r) and $CV(g) = \pm 1$.



This is the "QC-Folding Theorem".

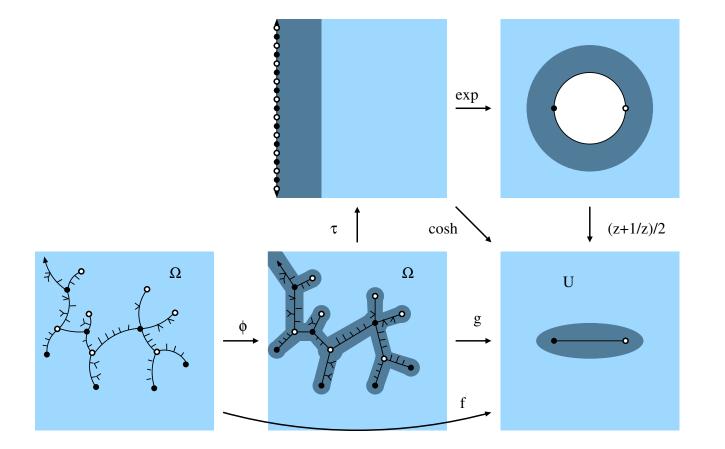
 $F = \cosh \circ \tau$ is usually discontinuous across T.

Thm: If T has bounded geometry and satisfies a τ -lower bound, then there is a QR g and r > 0 with $g = \cosh \circ \tau$ off T(r) and $CV(g) = \pm 1$.

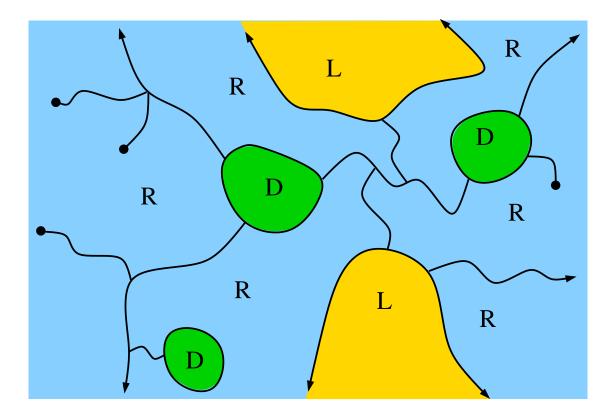


But g is continuous everywhere. Tree combinatorics are changed.

Thm: If T has bounded geometry and satisfies a τ -lower bound, then there is a QR g and r > 0 with $g = \cosh \circ \tau$ off T(r) and $CV(g) = \pm 1$.



By MRMT there is QC ϕ so $f = g \circ \phi$ is entire. Often can prove $\phi(z) \approx z$, so $T \approx f^{-1}([-1, 1])$ More generally: replace tree by graph:



R = unbounded domains ($F = e^{\tau(z)}$, previous case)

D = bounded Jordan domains $(F = (z - a)^n)$, high degree critical points) L = unbounded Jordan domains $(F = e^{-\tau(z)})$, finite asymptotic values) Can specify singular values in last two cases **exactly**. Transcendental = entire, not polynomial

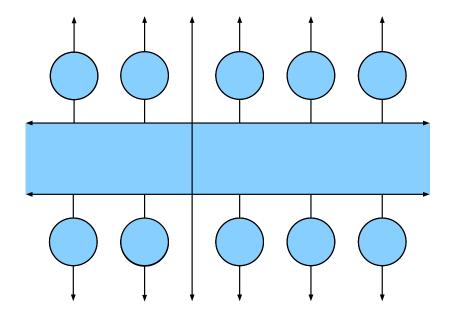
Singular set = closure of critical values and finite asymptotic values = smallest set so that f is a covering map onto $\mathbb{C} \setminus S$

Eremenko-Lyubich class = bounded singular set = \mathcal{B}

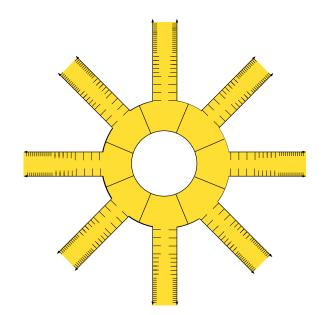
Speiser class = finite singular set = $\mathcal{S} \subset \mathcal{B}$

Generalized folding produces "all" functions in both classes.

Two applications of the folding theorem:



Entire functions with wandering domains. Transcendental Julia sets with small dimension

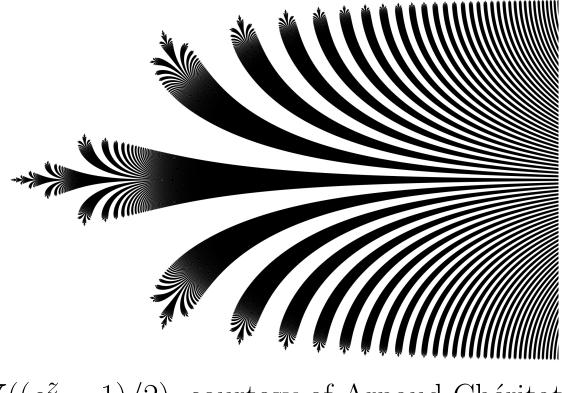


Dimensions of transcendental Julia sets:

Given an entire function f,

Fatou set $= \mathcal{F}(f) =$ open set where iterates are normal family. Julia set $= \mathcal{J}(f) =$ complement of Fatou set.

Julia set is usually fractal. What is its (Hausdorff) dimension?



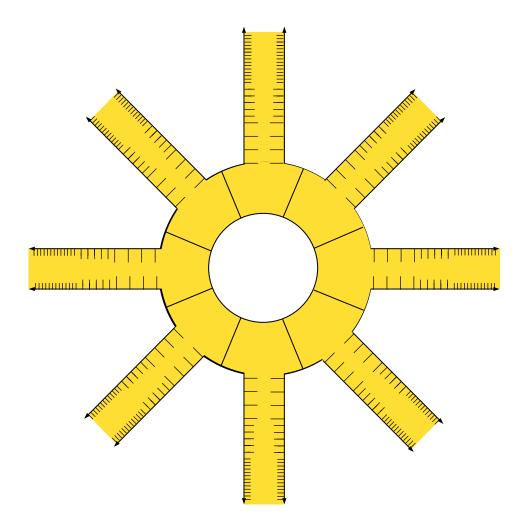
 $\mathcal{J}((e^z - 1)/2)$, courtesy of Arnaud Chéritat

A short (and incomplete) history:

- Baker (1975): f transcendental $\Rightarrow \mathcal{J}$ contains a continuum $\Rightarrow \dim \geq 1$.
- Misiurewicz (1981): $\mathcal{J}(e^x) = \mathbb{C}$.
- McMullen (1987): H-dim = 2, zero area.
- Stallard (1997, 2000): $\{\dim(\mathcal{J}(f)) : f \in \mathcal{B}\} = (1, 2].$
- Albrecht-B (2020): $1 < H-\dim < 2$ in Speiser class (folding thm).
- B (2018): H-dim = 1 example exists (folding thm).

Transcendentals: easy to get dim = 2, hard to get dim ≈ 1 .

Polynomials: dim = 2 is hard (Shishikura), area > 0 (Buff-Cheritat).



Speiser class Julia sets with dimension < 2. Build explicit QR map with small Julia set. Prove QC correction map is bi-Lipschitz near Julia set. Given an entire function f,

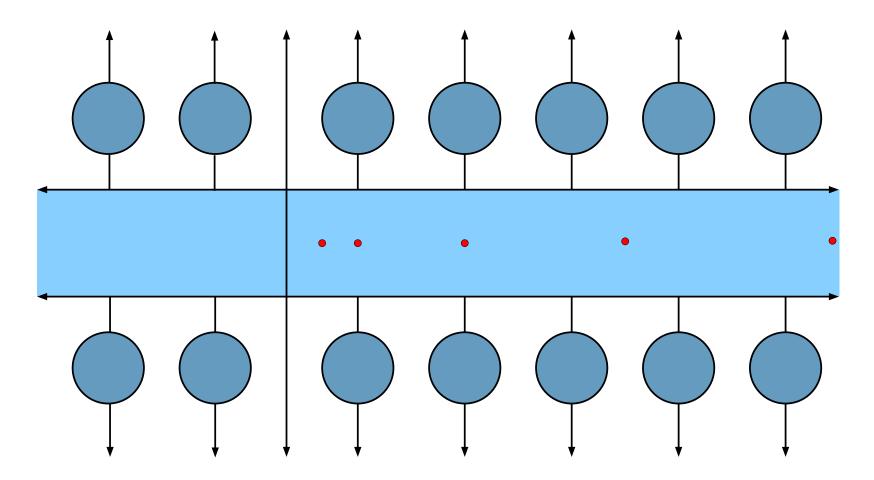
Fatou set $= \mathcal{F}(f) =$ open set where iterates are normal family. Julia set $= \mathcal{J}(f) =$ complement of Fatou set.

f maps Fatou components into other Fatou components.

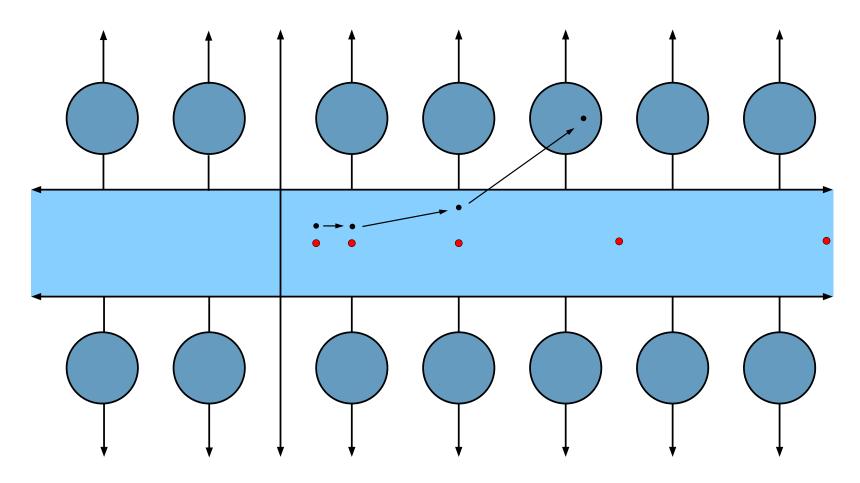
Wandering domain = Fatou component with infinite orbit.

- Entire functions can have wandering domains (Baker 1975).
- No wandering domains for rational functions (Sullivan 1985).
- Also none in Speiser class (Eremenko-Lyubich, Goldberg-Keen).

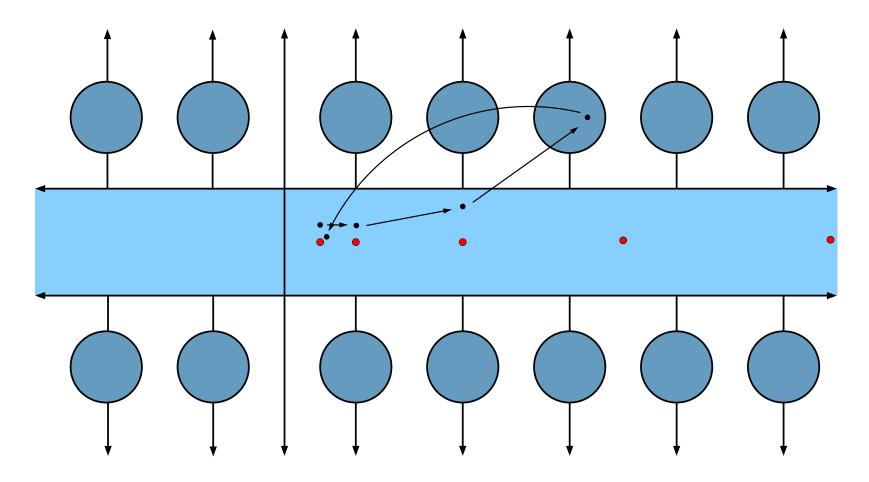
Wandering domains in Eremenko-Lyubich class? Open since 1985.



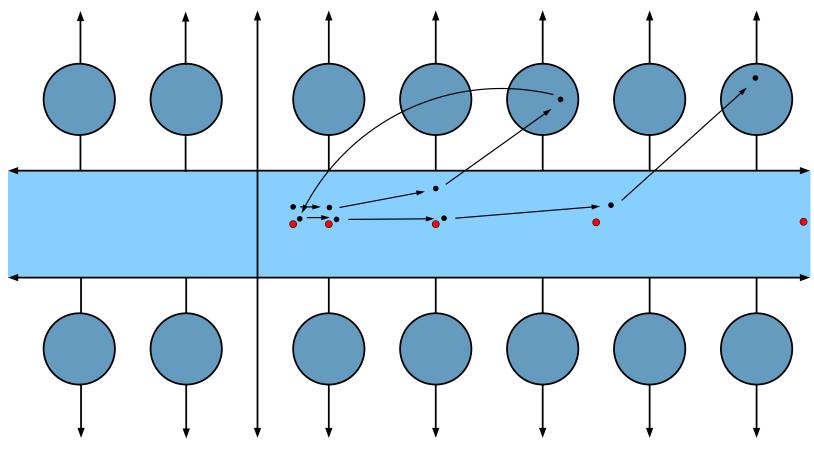
Graph giving EL wandering domain via folding (B 2015). Symmetry $\Rightarrow x = \frac{1}{2}$ iterates to $+\infty$ on \mathbb{R} .



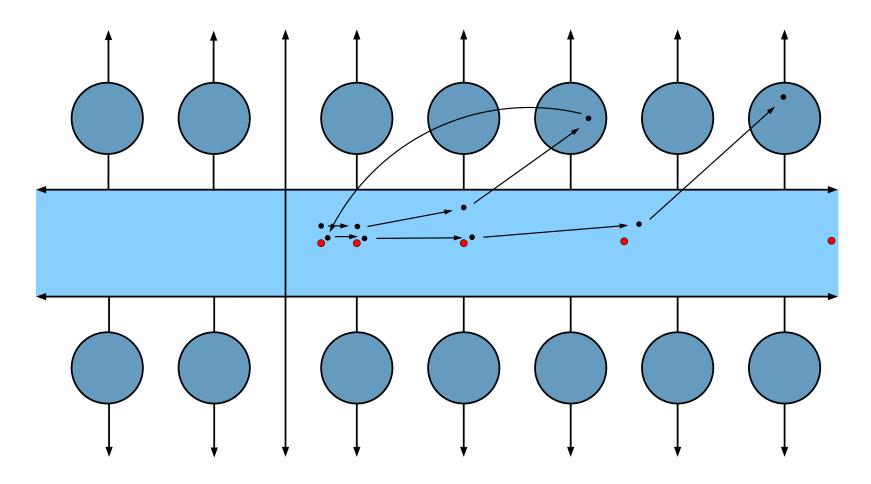
Iterates of $\frac{1}{2} + i\epsilon$ follow orbit of $\frac{1}{2}$ for a time, eventually diverge.



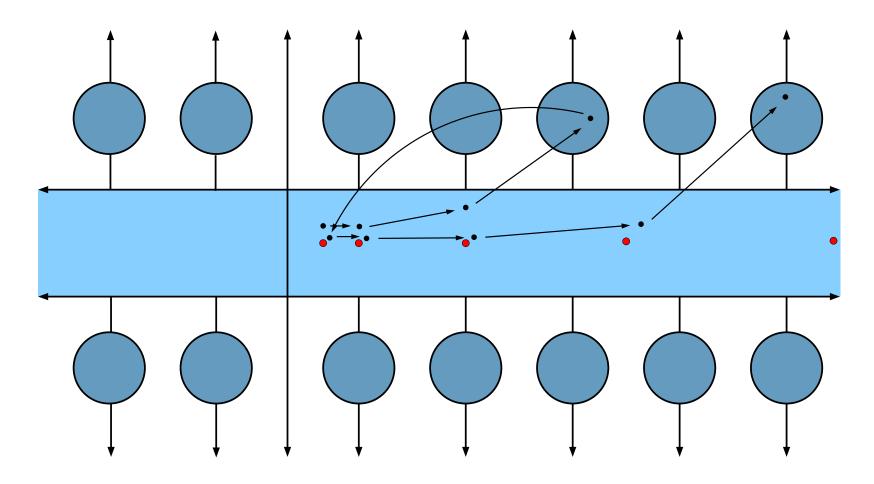
Orbit lands in D-component. Critical point compresses. Can specify that orbit lands very near $\frac{1}{2}$. Closer than before.



Follows $\frac{1}{2}$ -orbit longer, then diverges and returns near $\frac{1}{2}$.



Compression \Rightarrow orbit in Fatou component. Oscillation between $\frac{1}{2}$ and $\infty \Rightarrow$ wandering domain. (Schwarz lemma: iteration decreases hyperbolic metric.)

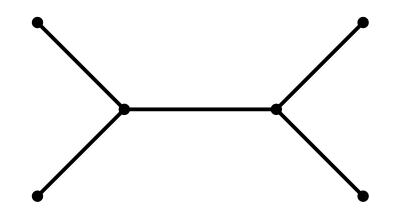


Alternate proof by Marti-Pete and Shishikura. Variations by Lazebnik, Fagella-Godillon-Jarque, Osborne-Sixsmith.

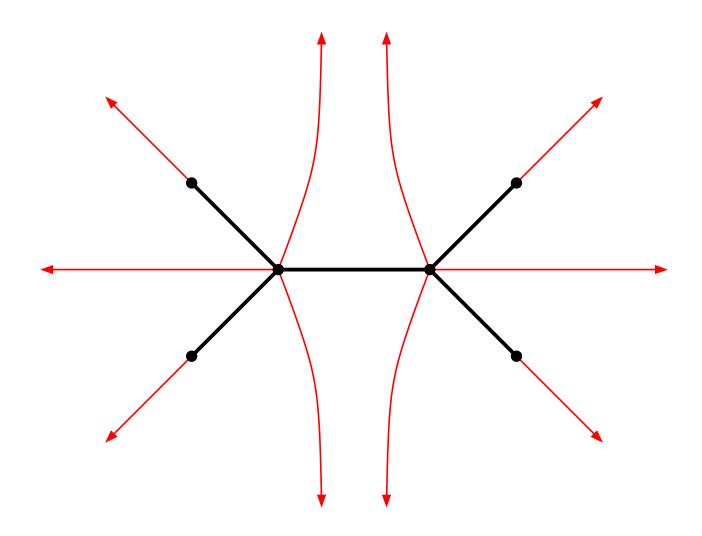
Dessins and equilateral triangulations

A graph (children's drawing) on a surface induces a conformal structure.

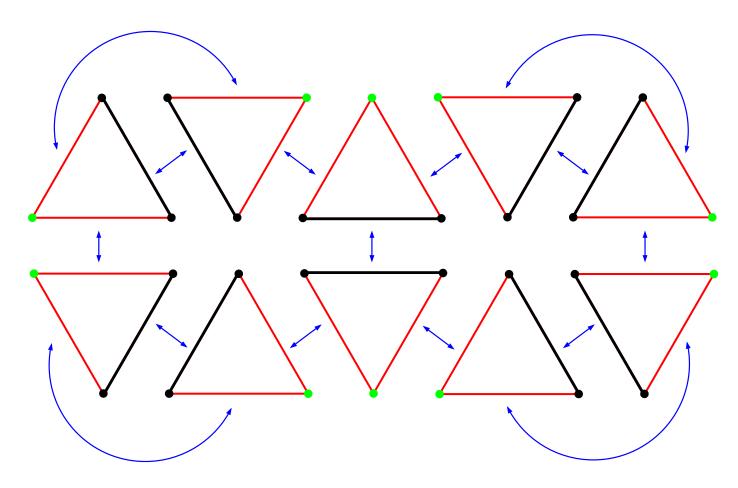
Consider trees on the sphere.



• start with a finite tree.



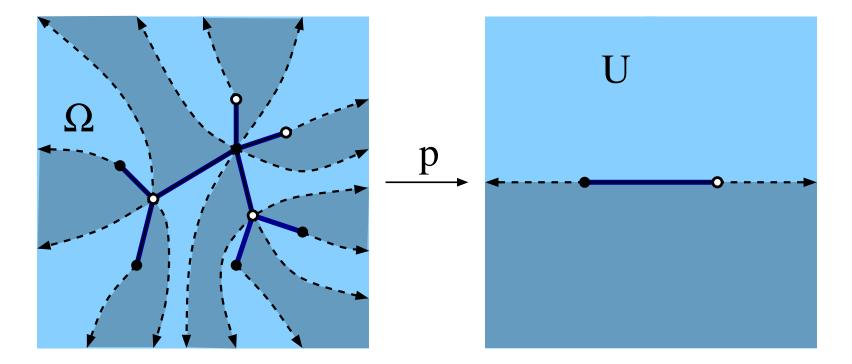
- connect vertices of T to infinity; gives finite triangulation of sphere.
- Defines adjacencies between triangles.



- Glue equilateral triangles using adjacencies: get a conformal 2-sphere.
- works for graphs on compact surfaces
- graph on surface \Rightarrow triangulation of surface \Rightarrow conformal structure

Only countably many ways to glue together finitely many triangles. Which compact Riemann surfaces occur in this way? **Belyi function** = holomorphic map from Riemann surface with 3 critical values. Shabat polys are examples on sphere.

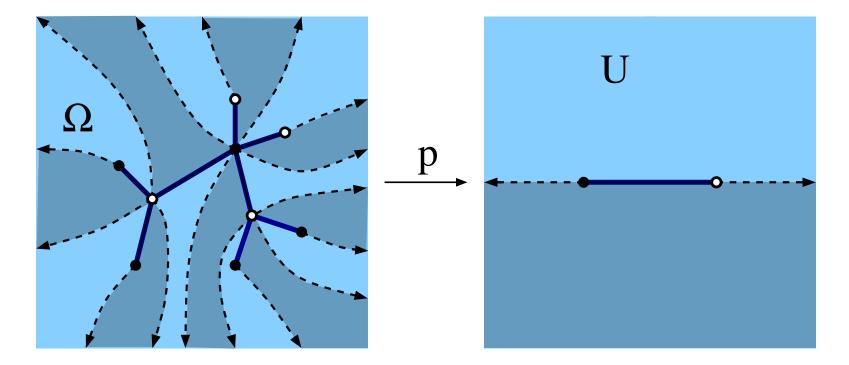
Theorem (Voevodsky-Shabat): A Riemann surface can constructed from equilateral triangles iff it has a Belyi function.



Triangles are inverse images of upper and lower half-planes.

Belyi function = holomorphic map from Riemann surface with 3 critical values. Shabat polys are examples on sphere.

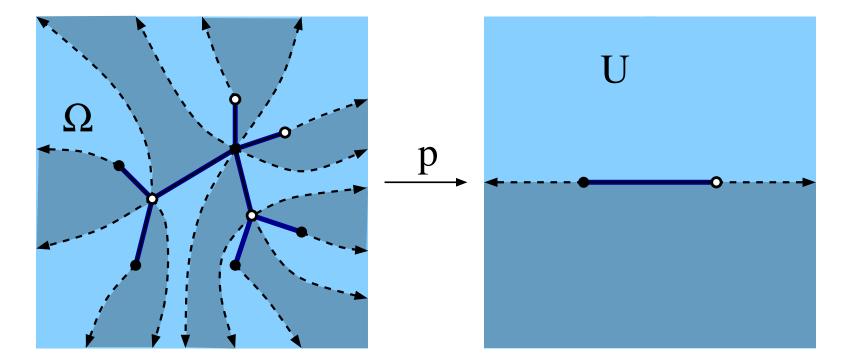
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Equilateral iff anti-holomorphic reflection between adjacent triangles. \Rightarrow whole triangulation is determined by a single triangle.

Belyi function = holomorphic map from Riemann surface with 3 critical values. Shabat polys are examples on sphere.

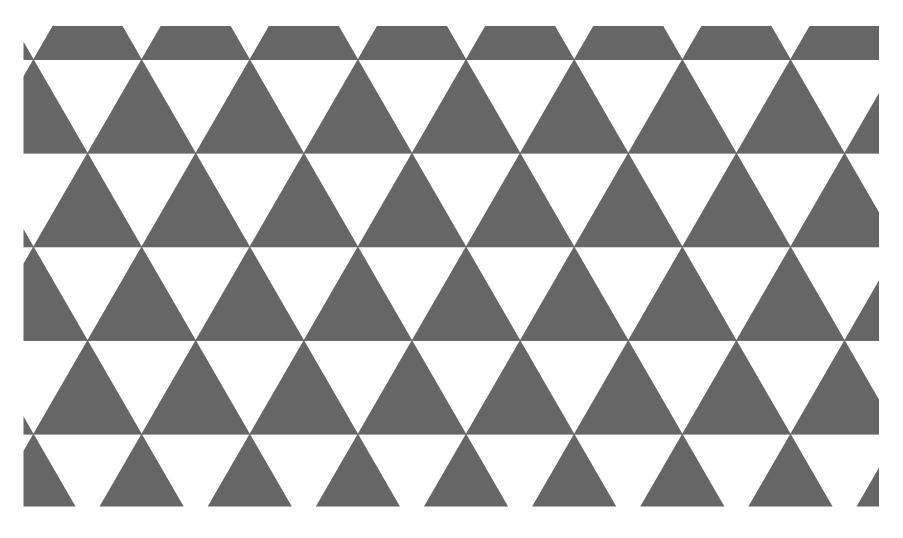
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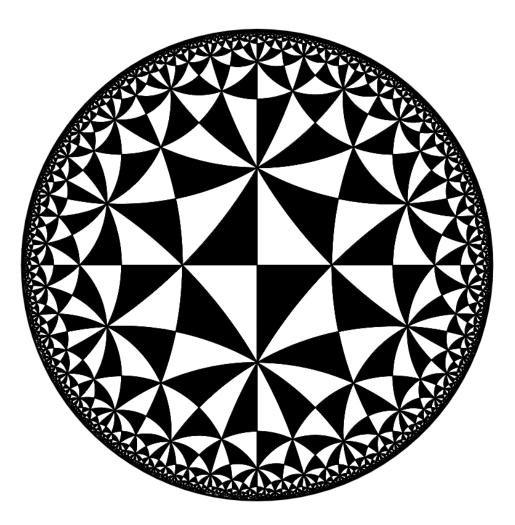
Belyi's Thm: A compact surface has a Belyi function iff it is algebraic. (Zero set of P(z, w) with algebraic integers as coefficients). What about non-compact Riemann surfaces?

Now we can use countably many triangles.

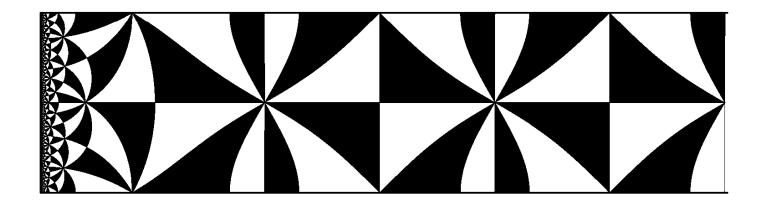
 \Rightarrow uncountably many way to glue them together.



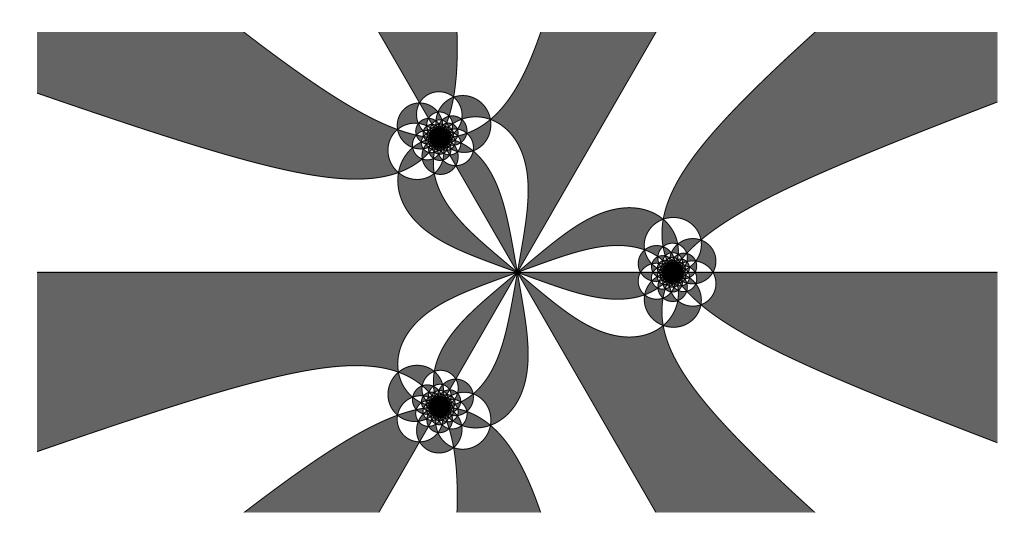
The plane



The disk



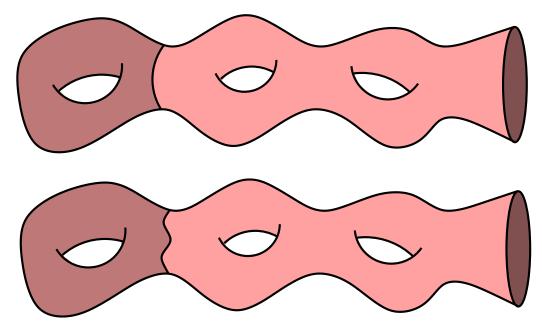
The punctured disk (identify top and bottom sides)



The trice punctured sphere

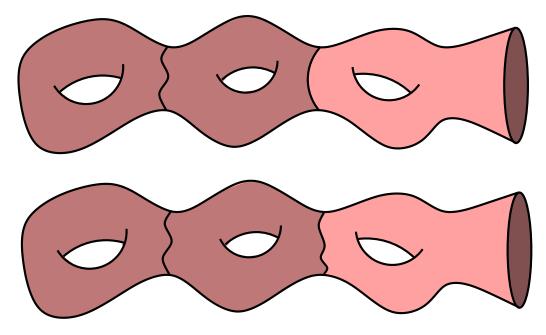
Idea of proof:

- Translate "true trees are dense" proof to Riemann surface S.
- Compact pieces can be approximated by triangulated surfaces.
- Conformal structure is changed, but as little as we wish.
- Key fact: small perturbation \Rightarrow triangulated pieces re-embed in S.
- Triangulate a compact exhaustion of S. Take limit.



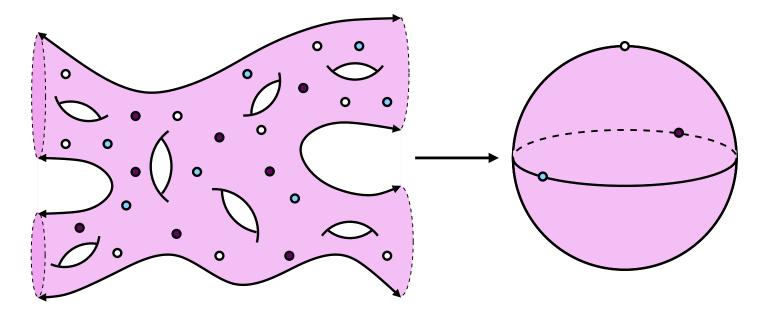
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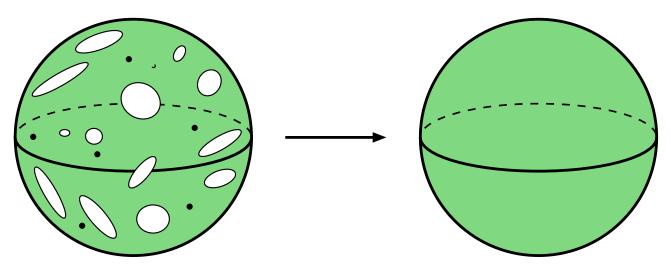


Corollary: Every Riemann surface is a branched cover of the sphere, branched over finitely many points.

- For compact surfaces, this is Riemann-Roch.
- Compact, genus g sometimes needs 3g branch points.
- 3 branch points suffice for all non-compact surfaces.



Corollary: Any open $U \subset \mathbb{C}$ is 3-branched cover of the sphere.



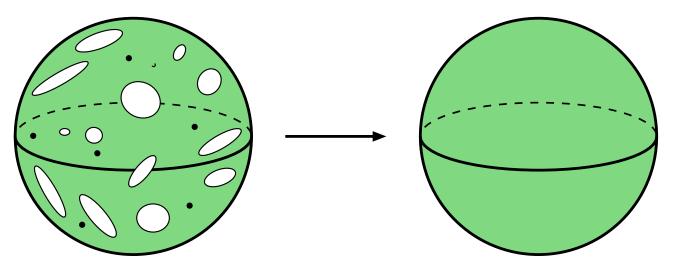
Adam Epstein: a holomorphic map $X \to Y$ is finite type if

- Y is compact,
- f is open with no isolated removable singularities,
- the set of singular values is finite.

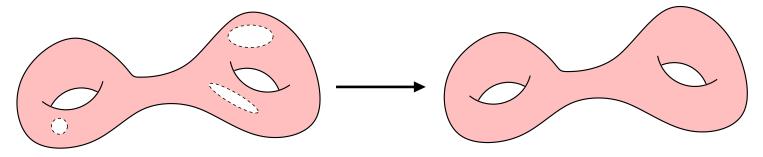
If $X \subset Y$, map can be iterated while orbits stay in X, e.g.,

- ▶ Rational maps on X = Y = 2-sphere
- ► Speiser class on $X = \mathbb{C}$, Y = 2-sphere

Corollary: Any open $U \subset \mathbb{C}$ is 3-branched cover of the sphere.



Gives many new dynamical systems of finite type.



Can have $f: X \to Y$ where $X \subset Y$ open, Y compact, genus > 0. Unknown which surfaces Y work. Belyi functions have finite singular sets.

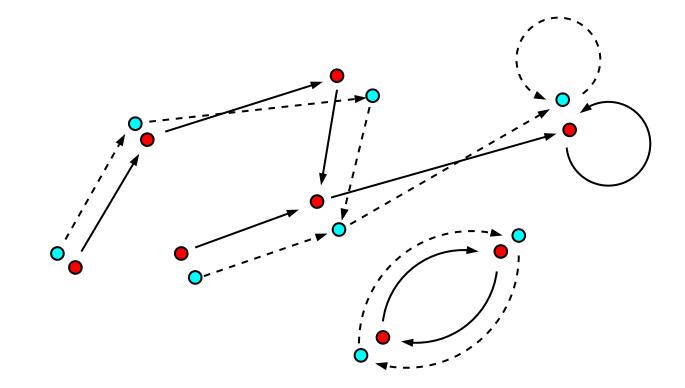
What about finite singular orbits?

Let P(f) = union of orbits of singular points.

Mentioned earlier that Julia sets with P(f) finite can take any "shape". Do all possible "combinatorics" occur? **Thm (DeMarco-Koch-McMullen):** For any finite $S \subset \mathbb{C}$ and $h : S \to S$, there is a rational f and a bijection $\psi : P(f) \to S$ so that $f = \psi^{-1} \circ h \circ \psi$ and $|\psi(z) - z| \leq \epsilon$.

"Any finite post-critical dynamics can occur for rational maps."

Their proof uses iteration on Teichmüller space.



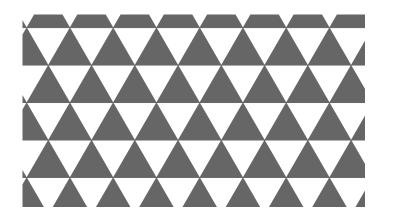
Lazebnik, Urbanski and I proved:

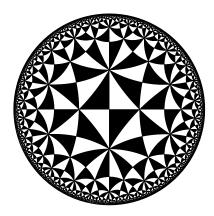
"Any discrete post-critical dynamics can occur for meromorphic maps."

Discrete = infinite set, accumulating only at boundary.

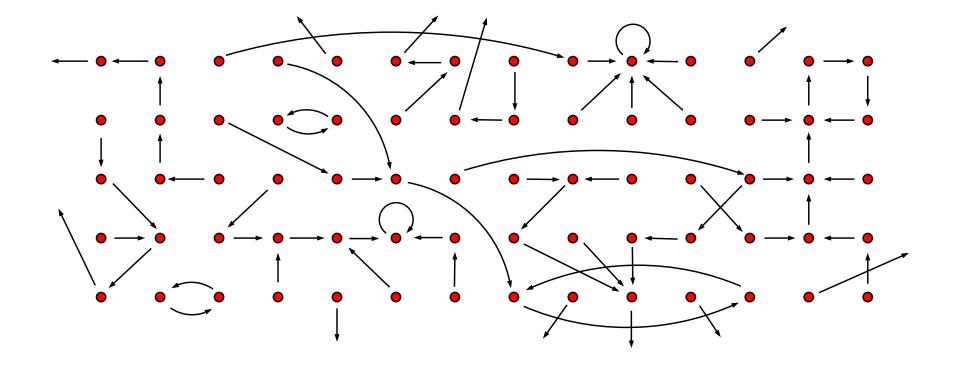
Proof uses a result on equilateral triangulations:

Thm (B-Lazebnik-Urbanski): For any decreasing $\eta : \mathbb{R} \to (0, 1)$, and any domain $\Omega \subset S^2$ there is an equilateral triangulation of Ω with $\operatorname{diam}(T) \leq \eta(\operatorname{dist}(T, \partial \Omega))$



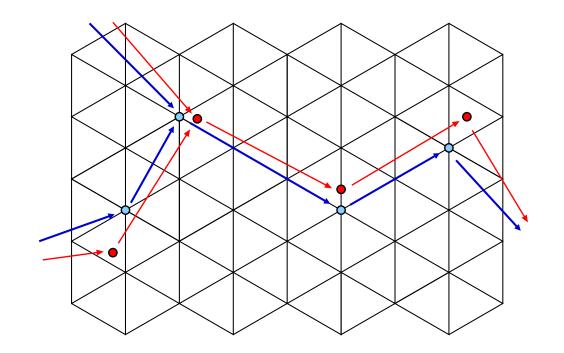


Thm: Given any discrete set $S \subset \mathbb{C}$ ($|S| \geq 3$) and any map $h: S \to S$, there is a meromorphic f and a bijection $\psi: P(f) \to S$ so that (1) $f = \psi^{-1} \circ h \circ \psi$ on P, (2) $|\psi(z) - z| \leq \eta(z)$ for all $z \in P$.



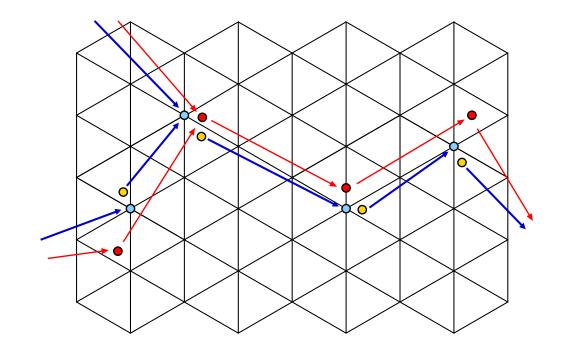
Idea of Meromorphic Proof:

- Take triangulation "finer" than S. Choose vertices approximating S.
- \bullet Perturb Belyi map to a QR map mimicking h on chosen vertices.



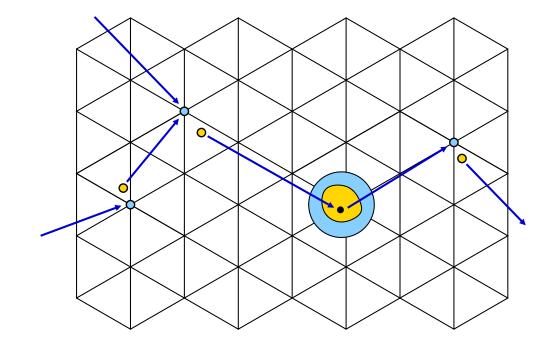
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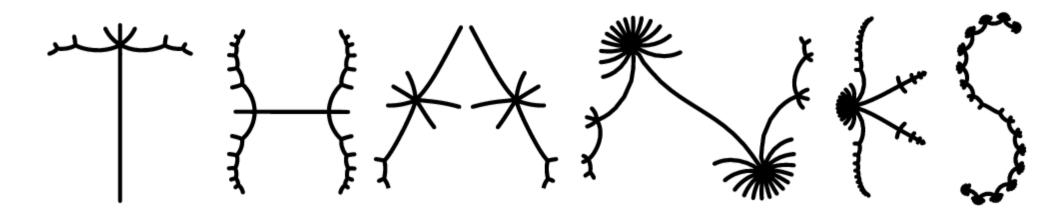
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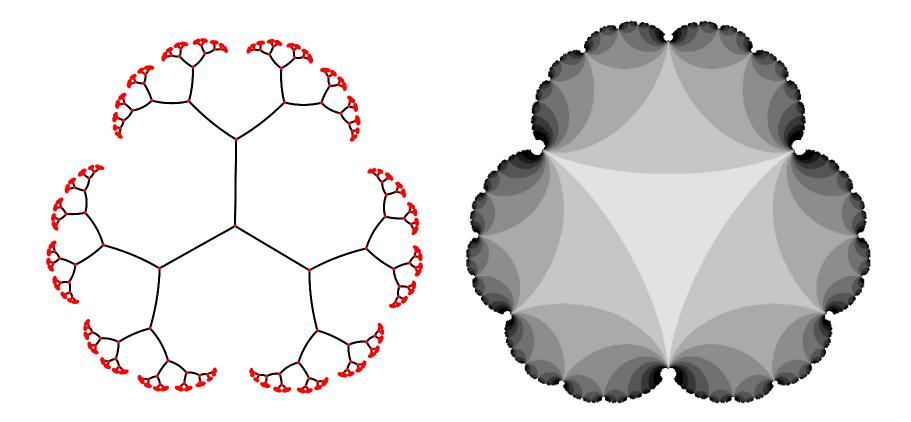


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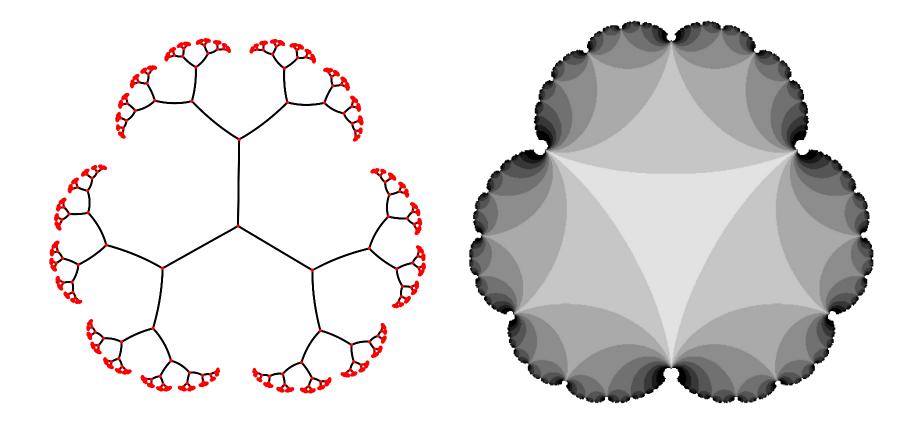
- Take triangulation "finer" than S. Choose vertices approximating S.
- \bullet Perturb Belyi map to a QR map mimicking h on chosen vertices.
- But holomorphic correction does not map vertices to vertices.
- If we move critical values, can show critical points move less.
- Infinite dimensional fixed point theorem gives desired meromorphic map.



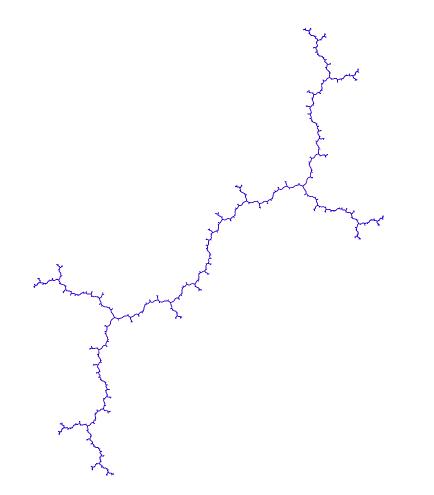




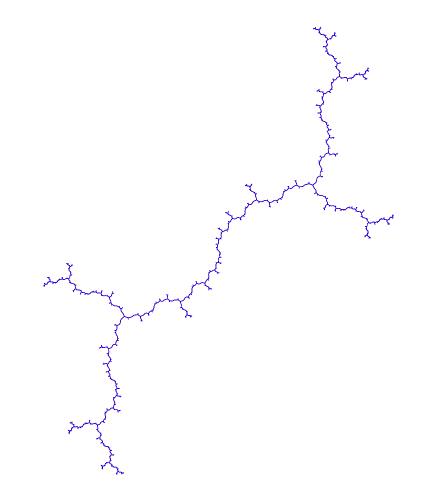
True form of truncated finite 3-regular tree, and limit set of reflected deltoid group.



Proved equal by Oleg Ivrii, Peter Lin, Steffen Rohde and Emanuel Sygal Related to matings of Julia sets with Kleinian groups, studied by Lee, Lyubich, Makarov and Mukherjee.



This is not the Julia set of $z^2 + c$, $c \approx 0.288 + 1.115i$. (But it looks exactly like it.)



Tree approximating combinatorics of the Julia set. Edges not all equal harmonic measure, but still identified by conformal map. Rigidity: combinatorial data determines geometry.