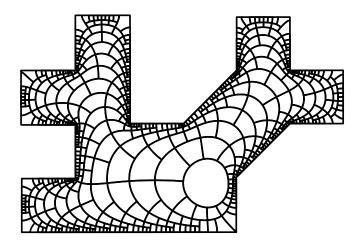
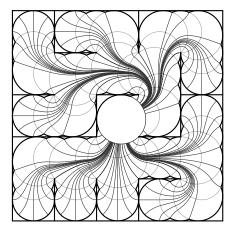
MAPPINGS AND MESHES I: CONNECTIONS BETWEEN CONTINUOUS AND DISCRETE GEOMETRY

Christopher Bishop, Stony Brook University

Geometric Methods for Analyzing Discrete Shapes Center of Mathematical Sciences and Applications Harvard University, May 7-9, 2021









THE IDEA

In analysis we often want to map a complicated region to a simple shape, such as a disk, preserving some property, e.g., angles.

In computational geometry we often want to decompose a complicated shape into simple regions, like triangles, bounding the number of pieces needed and the angles used.

These problems are closely connected.

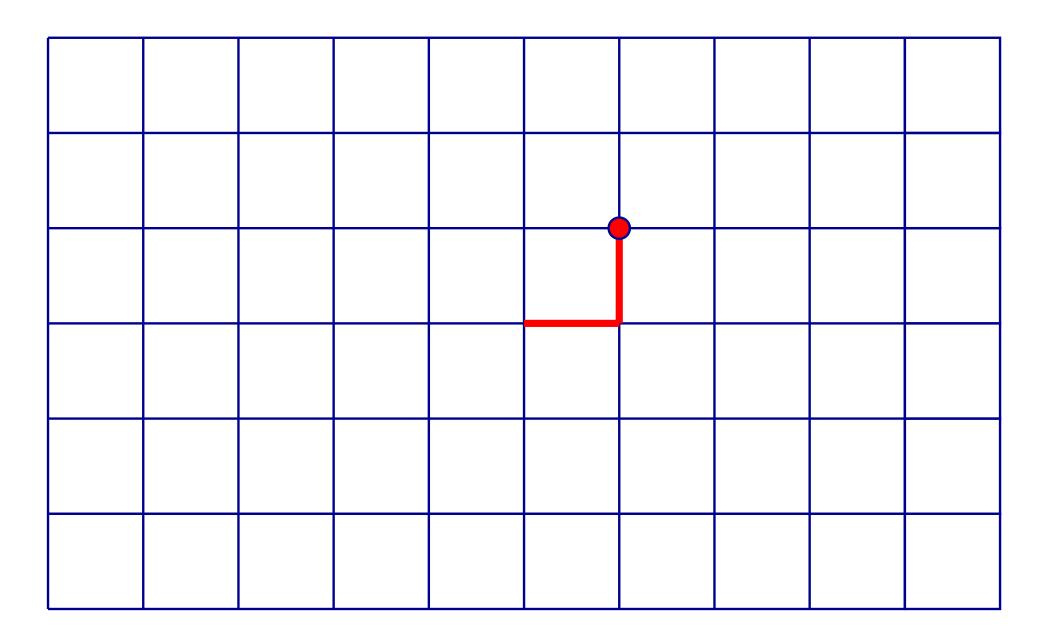
THE PLAN

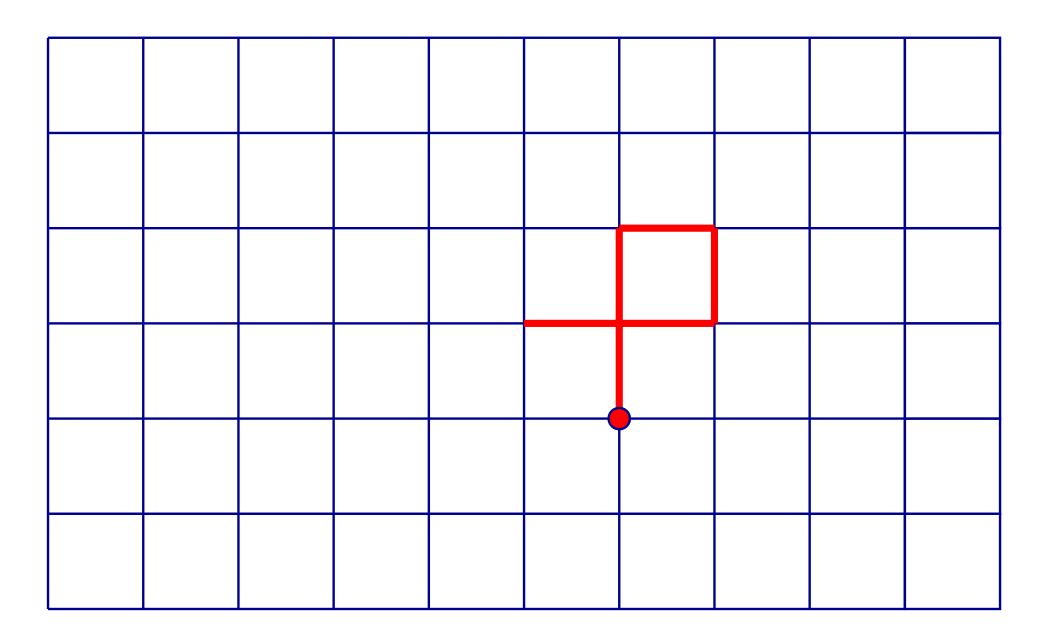
Lecture 1: Medial axis approximates conformal maps

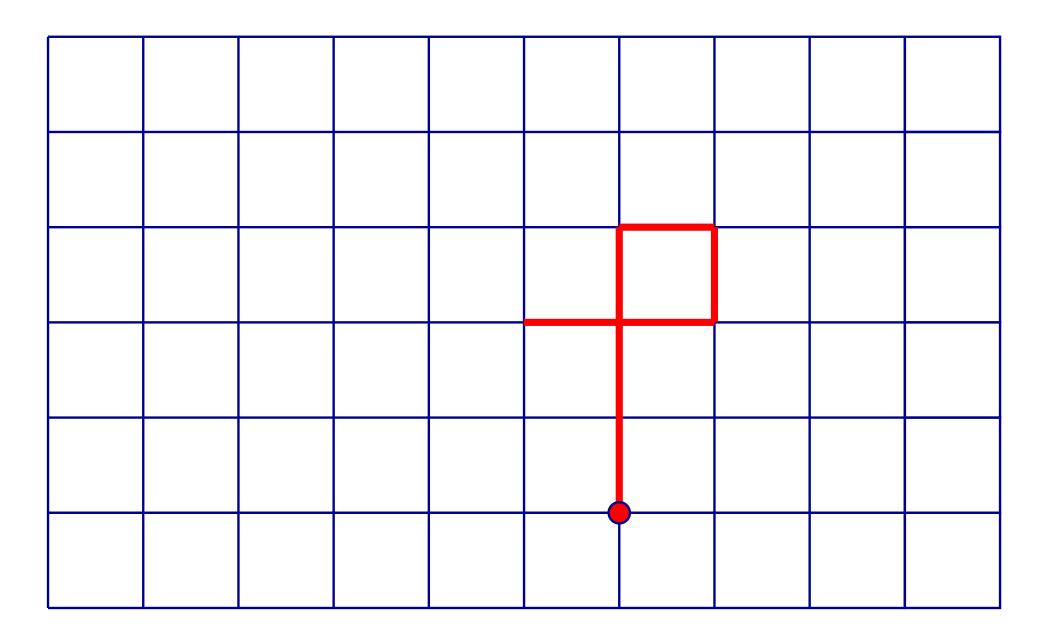
- Harmonic measure and conformal maps
- The Schwarz-Christoffel formula
- The medial axis
- Convex hulls in hyperbolic space

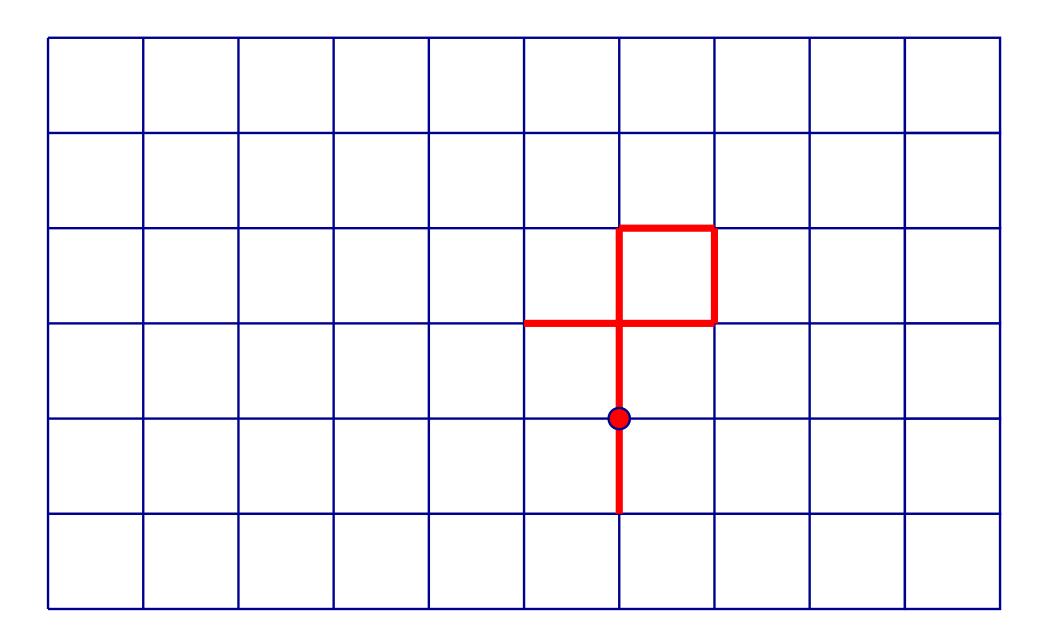
Lecture 2: Conformal maps give good meshes

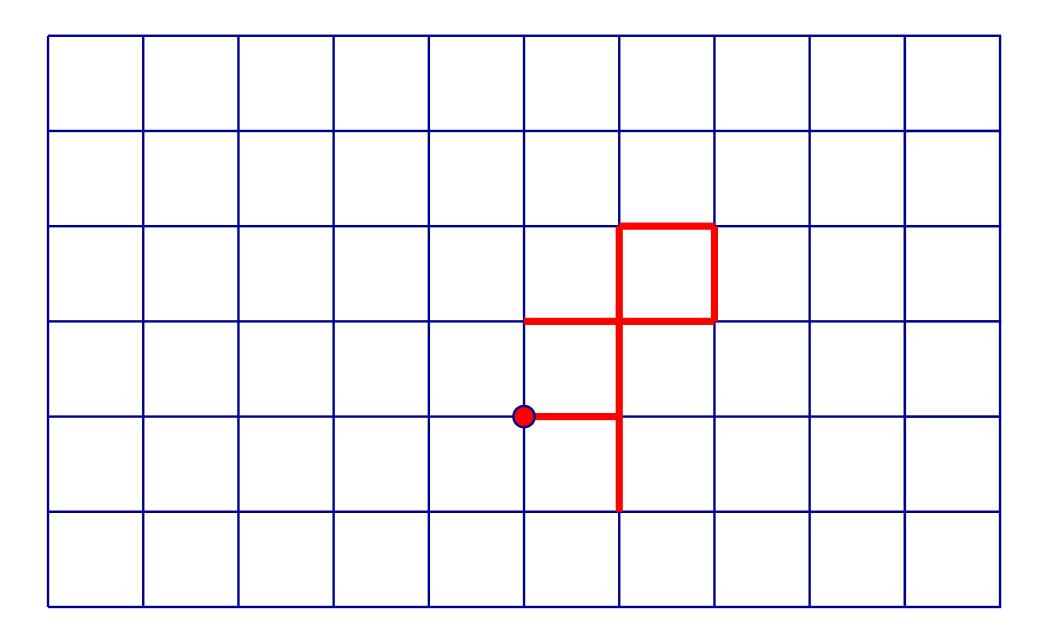
- Optimal quad-meshes
- Optimal triangulations of polygons
- Non-obtuse triangulations of PSLGs

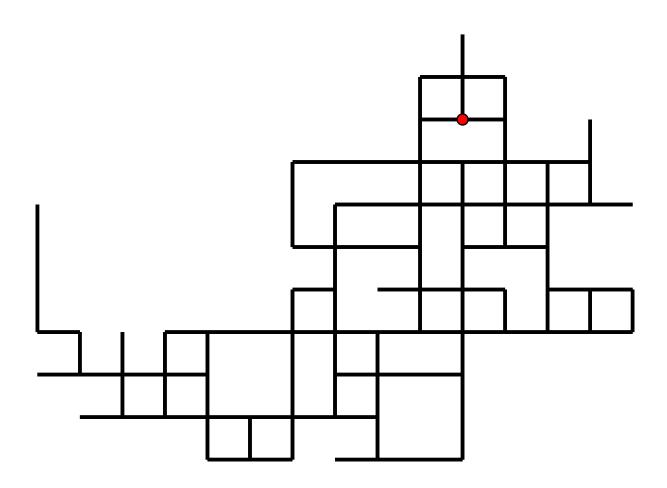




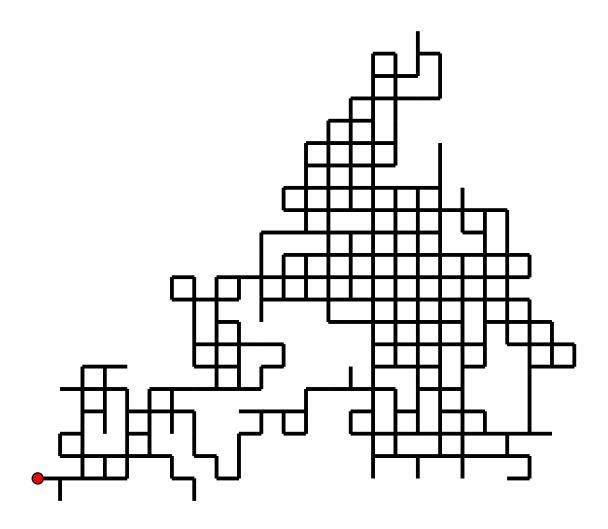




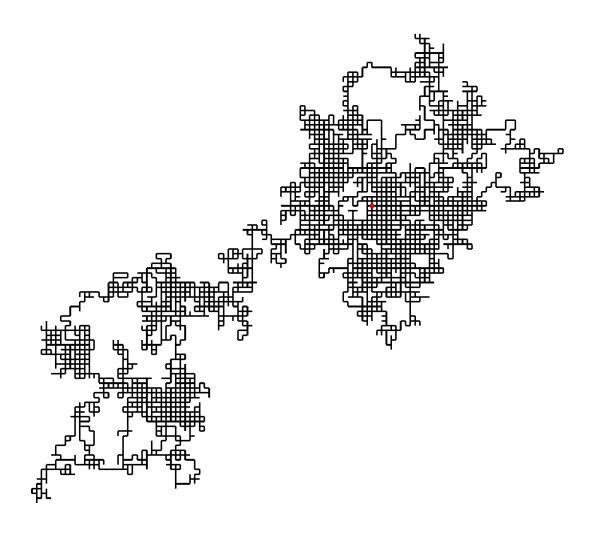




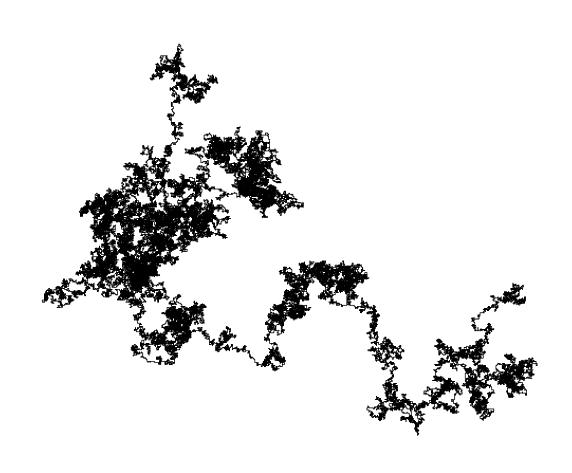
200 step random walk.



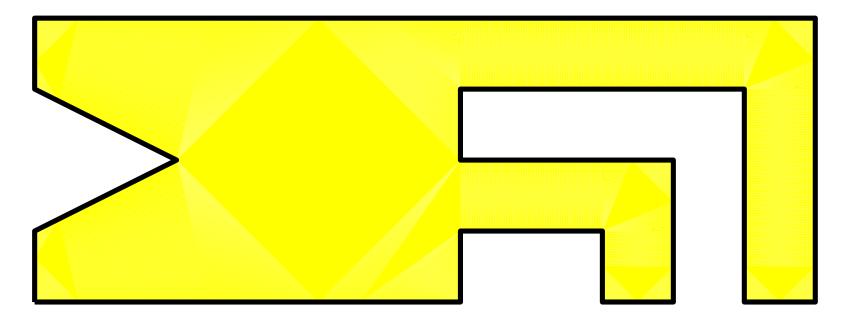
1000 step random walk.



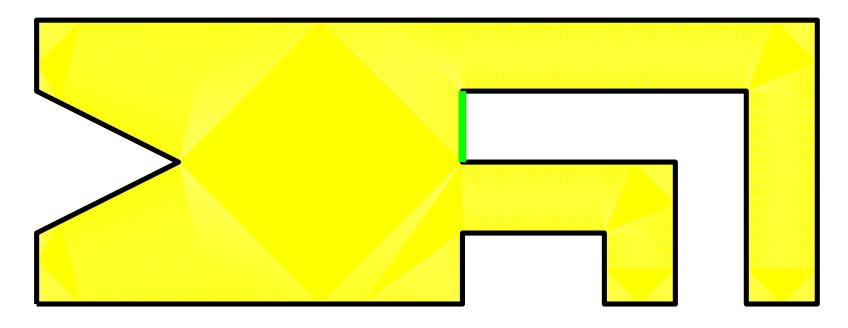
10,000 step random walk.



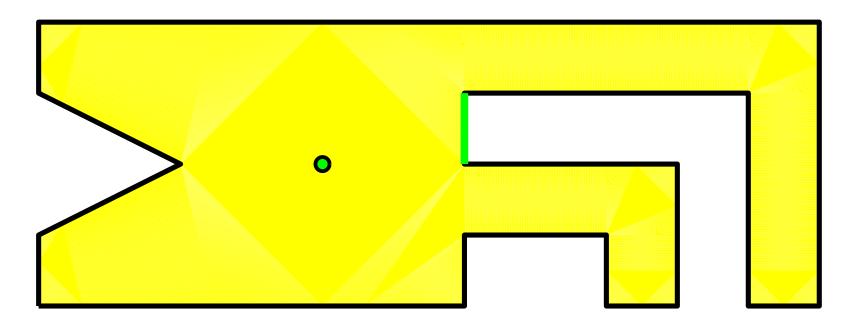
100,000 step random walk.



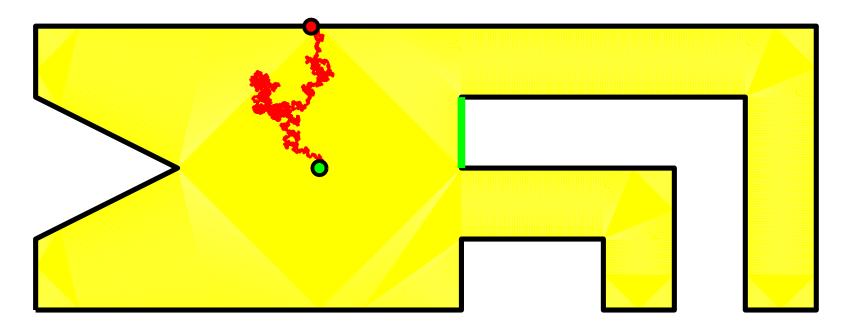
Suppose Ω is a planar Jordan domain.

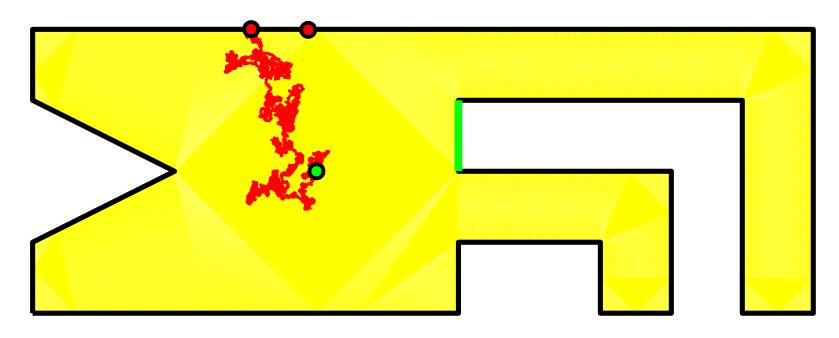


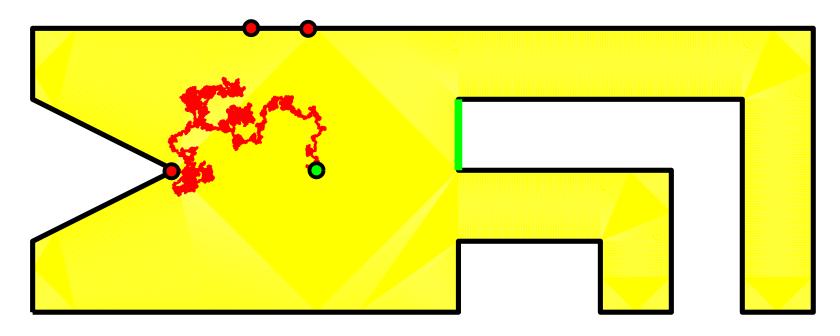
Let E be a subset of the boundary, $\partial\Omega$.

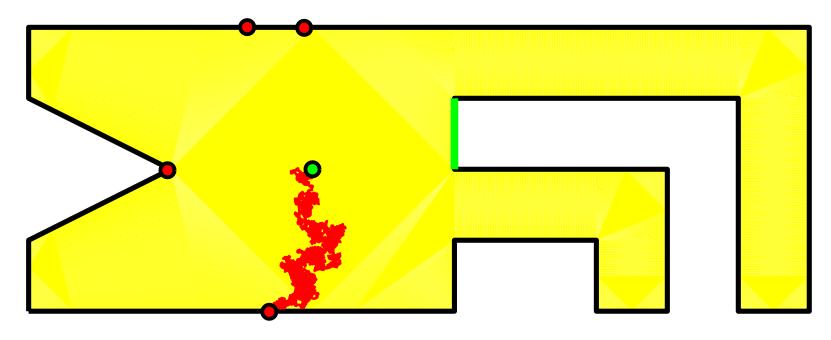


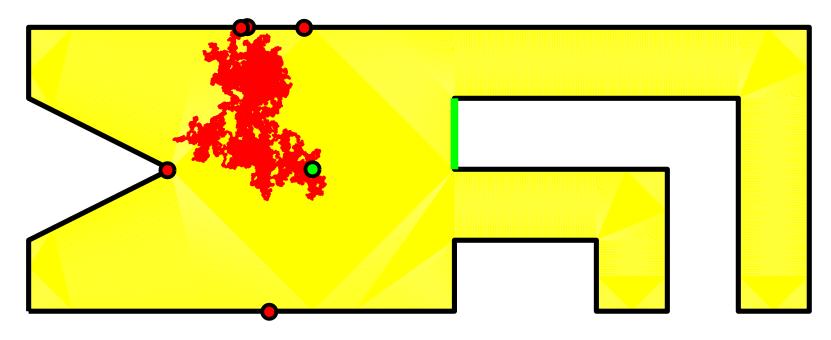
Choose an interior point $z \in \Omega$.

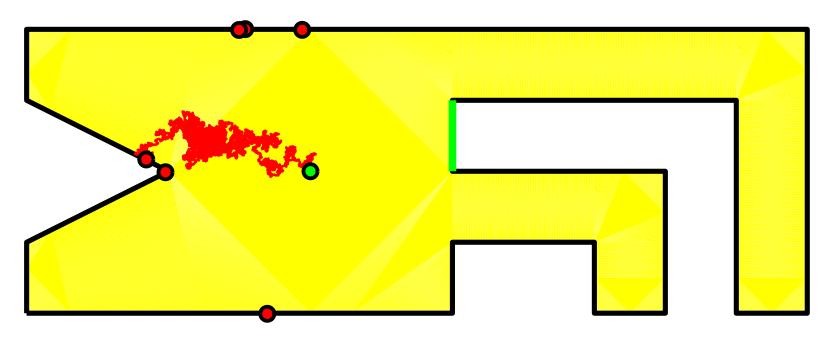


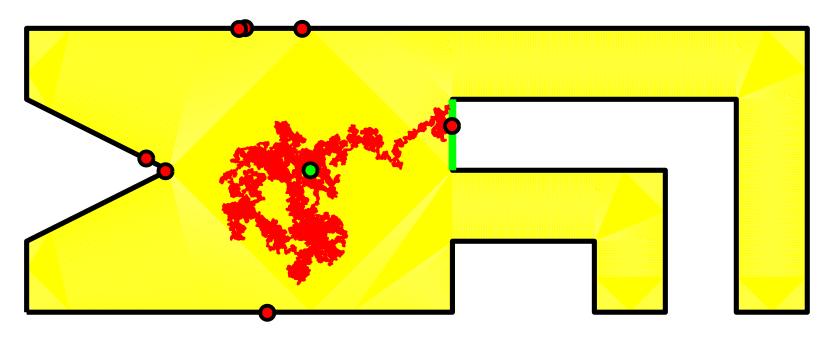


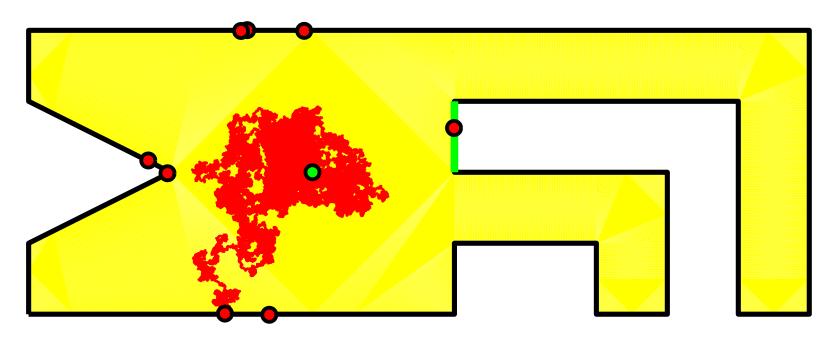


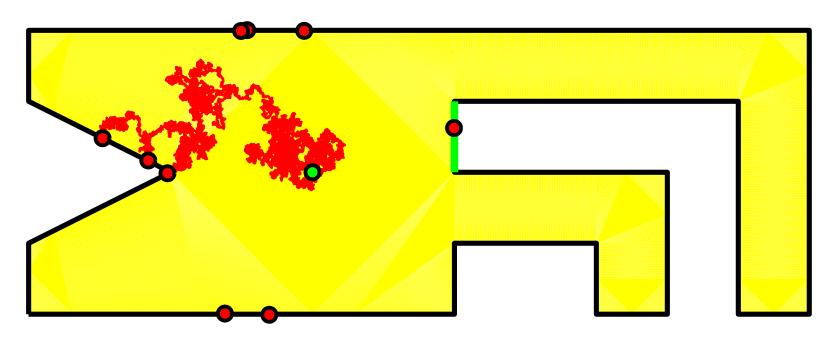


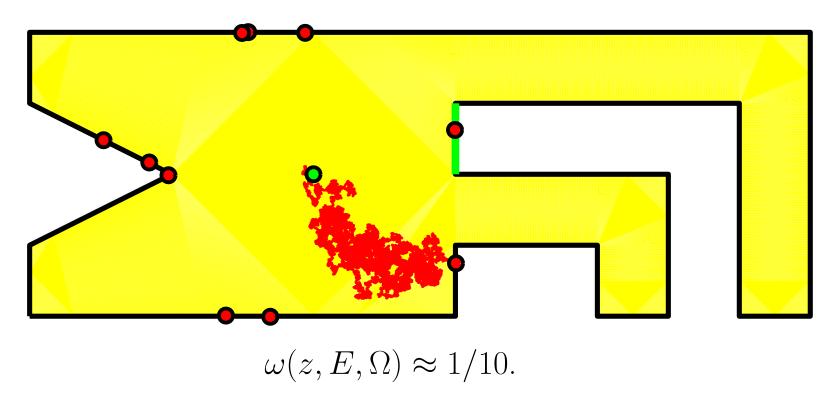


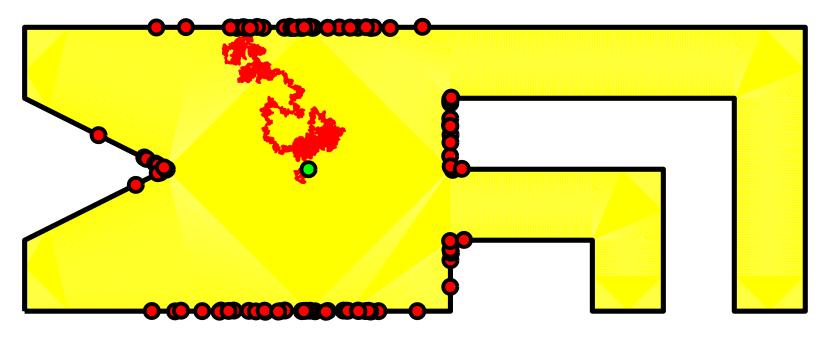




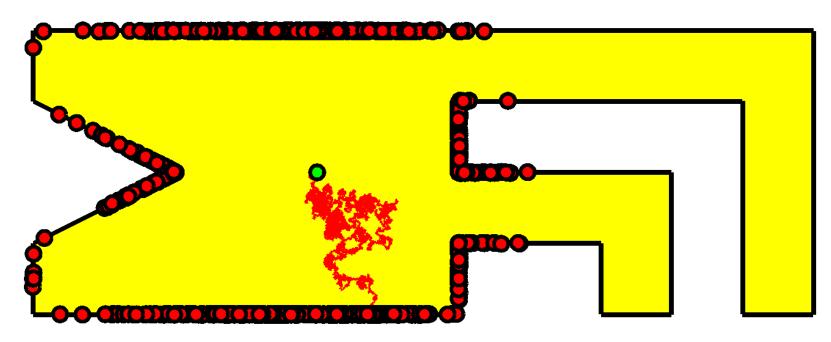






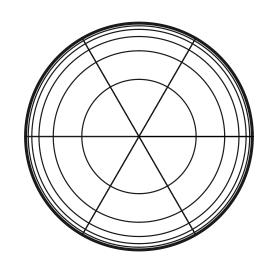


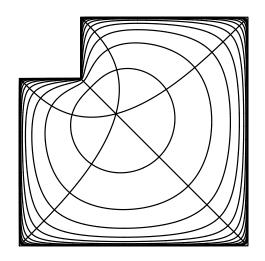
 $\omega(z, E, \Omega) \approx 13/100.$



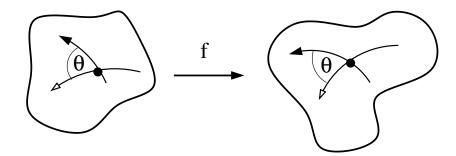
 $\omega(z, E, \Omega) \approx 126/1000.$

Riemann Mapping Theorem: If $\Omega \subsetneq \mathbb{R}^2$ is simply connected, then there is a conformal map $f: \mathbb{D} \to \Omega$.

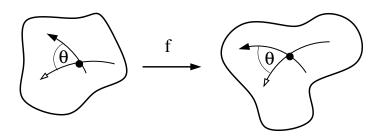




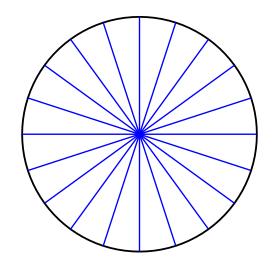
Conformal = angle preserving

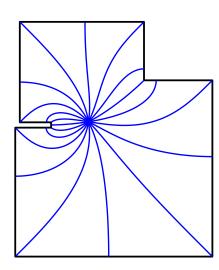


Riemann Mapping Theorem: If $\Omega \subsetneq \mathbb{R}^2$ is simply connected, then there is a conformal map $f: \mathbb{D} \to \Omega$. (conformal = angle preserving)

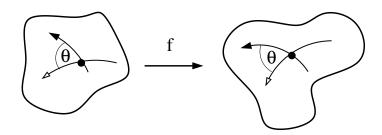


Brownian motion is conformally invariant, so normalized length measure maps to harmonic measure. Fastest way to compute harmonic measure.

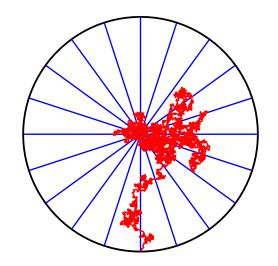


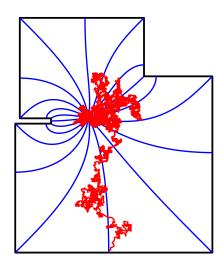


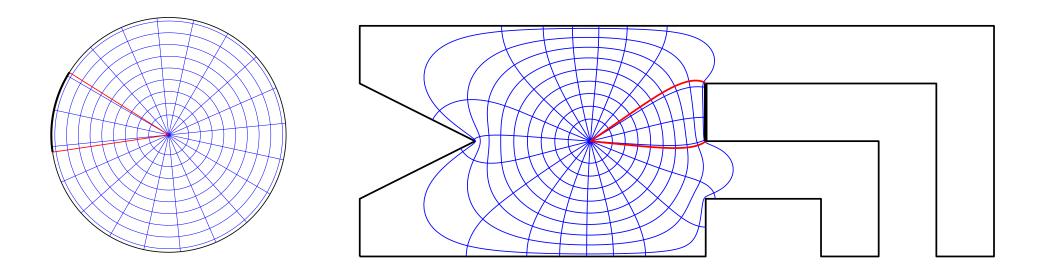
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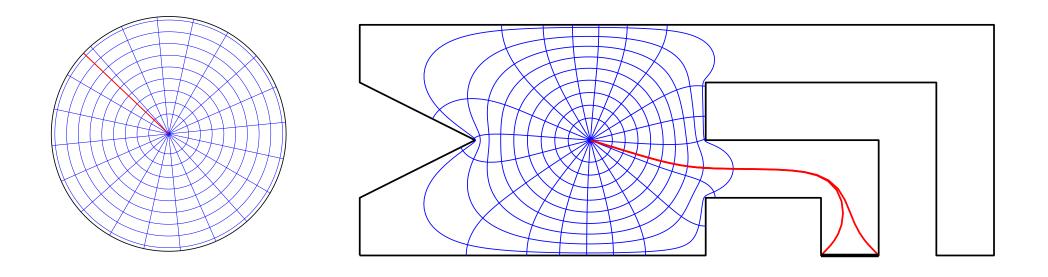
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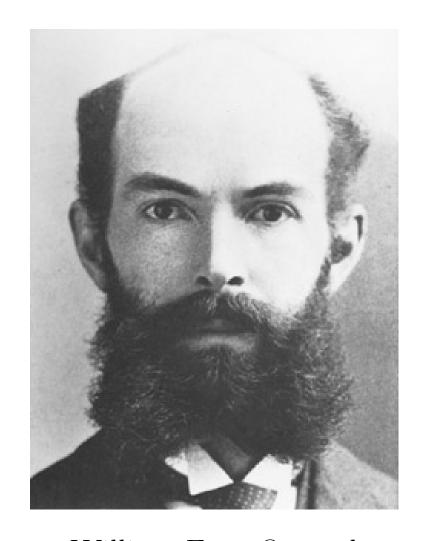
harmonic measure ≈ 0.1128027



harmonic measure $\approx 1.22155 \times 10^{-6}$



Georg Friedrich Bernhard Riemann Stated RMT in 1851



William Fogg Osgood First proof of RMT, Trans. AMS, vol. 1, 1900 Harvard 1866, Math Faculty 1890-1933, Chair 1918-22

The proof of Osgood represented, in my opinion, the "coming of age" of mathematics in America. Until then, ... the mathematical productivity in this country in quality lagged behind that of Europe, and no American before 1900 had reached the heights that Osgood then reached.

J.L. Walsh, "History of the Riemann mapping theorem", Amer. Math. Monthly, 1973.

Schwarz-Christoffel formula for maps to polygons (1867):

$$f(z) = A + C \int_{-\infty}^{z} \prod_{k=1}^{m} (1 - \frac{w}{z_k})^{\alpha_k - 1} dw,$$



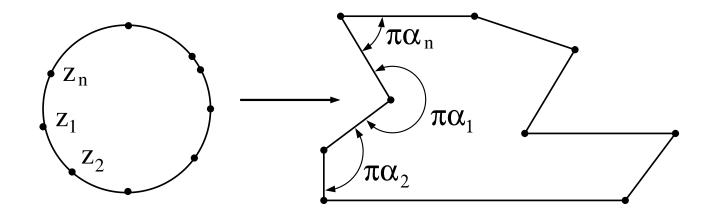
Christoffel



Schwarz

Schwarz-Christoffel formula for maps to polygons (1867):

$$f(z) = A + C \int_{-\infty}^{z} \prod_{k=1}^{m} (1 - \frac{w}{z_k})^{\alpha_k - 1} dw,$$



 α 's known. z's unknown (= SC-parameters = pre-vertices)

Finding SC-parameters = Finding harmonic measure of edges

Numerical conformal mapping:

- Koebe
- Theodorsen
- Fornberg
- Wegman
- Gaier
- Symm
- Kerzman-Stein
- Integral equations via fast multipole, Rokhlin
- Circle packing, Sullivan, Rodin, Stephenson
- CRDT, Driscoll and Vavasis
- SCToolbox, Trefethen, Driscoll
- ZIPPER, Marshall

Problem: given n-gon, how fast can we compute the SC-parameters?

Theorem: Can compute ϵ -conformal map onto n-gon in time $C_{\epsilon} \cdot n$.

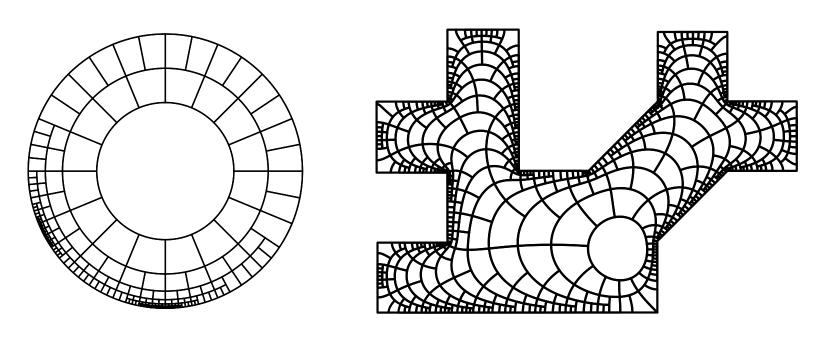
Theorem: Can compute ϵ -conformal map onto n-gon in time $C_{\epsilon} \cdot n$.

 ϵ -conformal = 1 + ϵ quasiconformal.

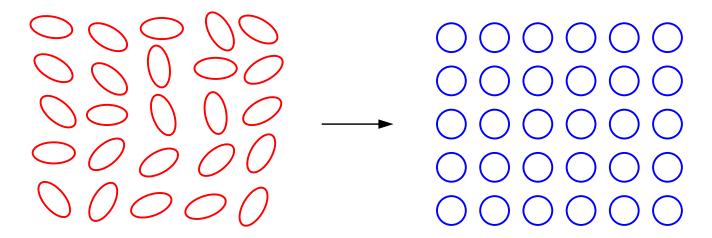
$$C_{\epsilon} = O(\log \frac{1}{\epsilon} \log \log \frac{1}{\epsilon}).$$

Data held as O(n) Laurent series of length $p = \log \frac{1}{\epsilon}$.

Bottleneck is doing O(1) FFTs per vertex of polygon.



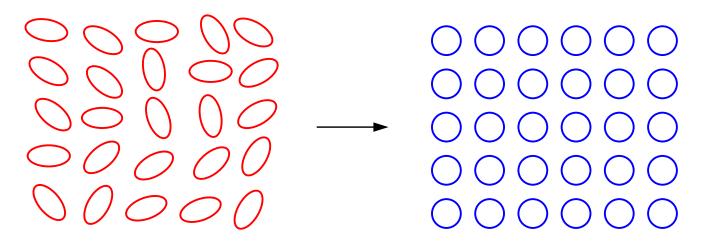
Quasiconformal (QC) maps are homeomorphisms that are differentiable a.e. and send infinitesimal ellipses to circles.



Eccentricity = ratio of major to minor axis of ellipse.

For K-QC maps, ellipses have eccentricity $\leq K$

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Eccentricity = ratio of major to minor axis of ellipse.

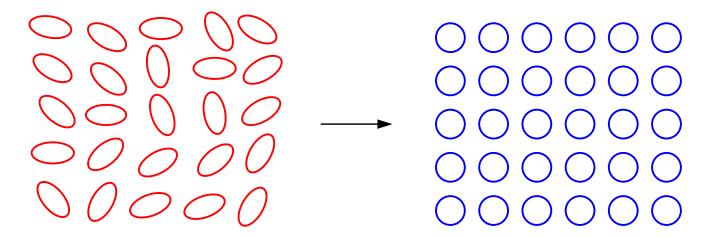
For K-QC maps, ellipses have eccentricity $\leq K$

Ellipses determined a.e. by measurable dilatation

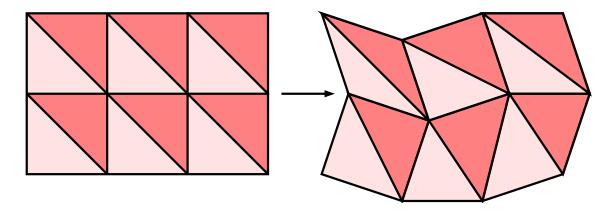
$$\mu = f_{\overline{z}}/f_z, \qquad f_{\overline{z}} = \mu \cdot f_z, \quad \text{with } |\mu| \le \frac{K-1}{K+1} < 1.$$

Here $f_z = f_x - if_y$ and $f_{\overline{z}} = f_x + if_y$.

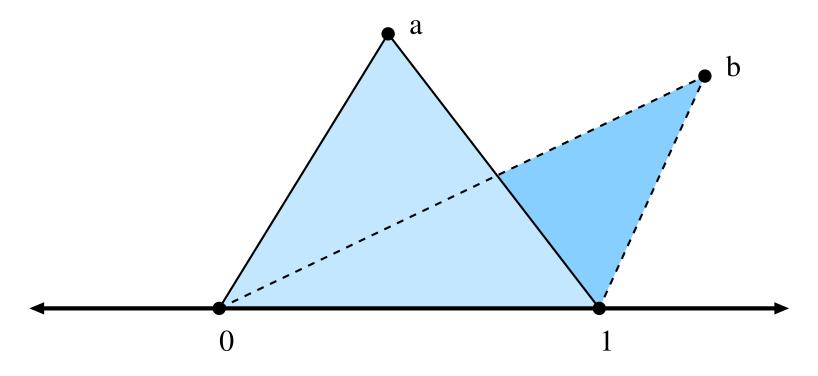
Quasiconformal (QC) maps are homeomorphisms that are differentiable a.e. and send infinitesimal ellipses to circles.



Example: piecewise affine maps between triangulations.



Map is QC if all angles bounded above and below.

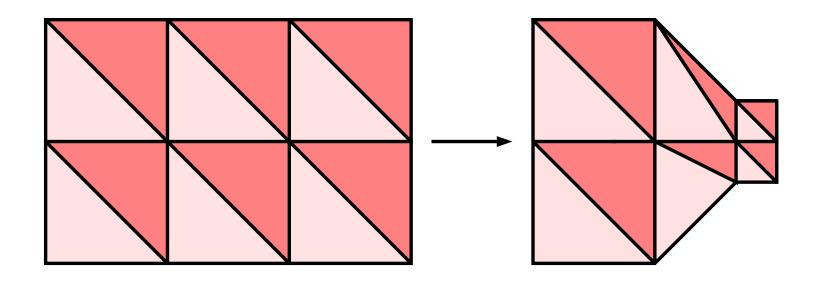


Affine map between triangles $\{0, 1, a\}$ and $\{0, 1, b\}$ has constant dilatation

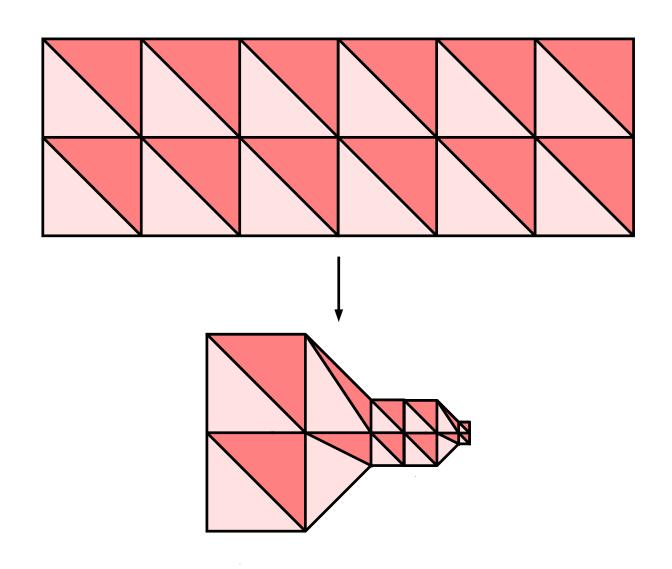
$$\mu = \frac{b - a}{b - \overline{a}}$$

(For experts, this is pseudo-hyperbolic distance in upper half-plane.)

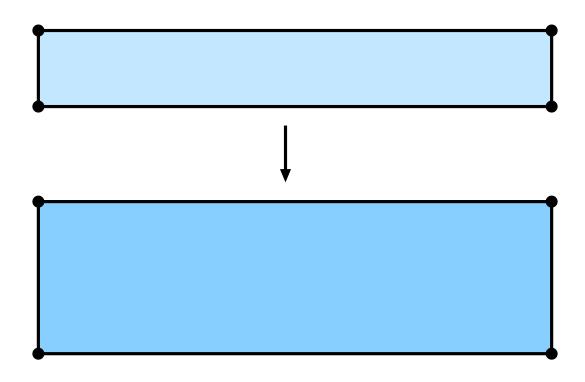
Example: piecewise affine maps between triangulations.



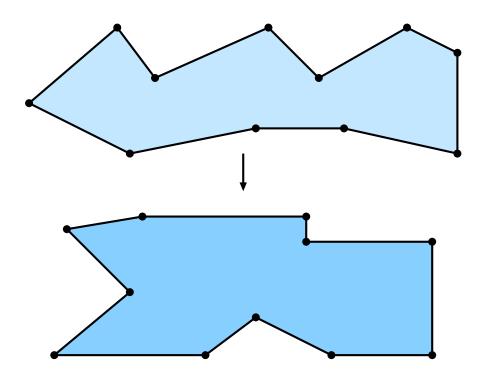
Example: piecewise affine maps between triangulations.



QC-distance for n-gons defined by optimal QC map preserving vertices.

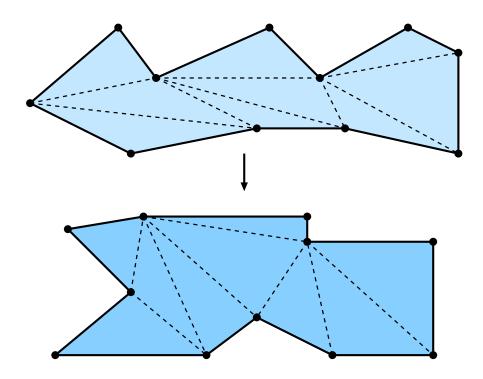


For rectangles, optimal map is linear stretch $(x, y) \to (xa \cdot y)$.



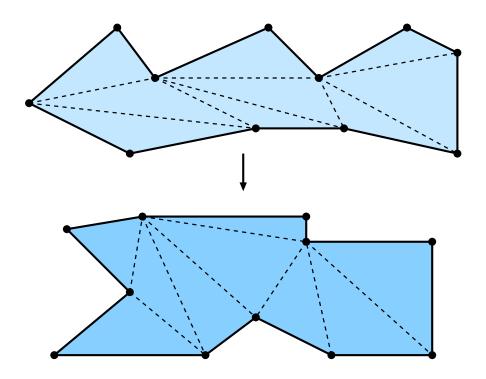
In general, optimal QC map is hard to compute.

See "Computing Teichmüller Maps between Polygons", Goswami, Gu, Pingali, Telang (2014)



Any map gives an upper bound.

Find compatible triangulation and compute dilatation of affine maps.



Any map gives an upper bound.

Find compatible triangulation and compute dilatation of affine maps.

If P_1, P_2 are ϵ -close in QC sense, SC-parameters are $O(\epsilon)$ -close on circle.

Estimates distance to SC-parameters without knowing SC-parameters.

Recall that the dilatation of a QC map is

$$\mu = \frac{f_{\overline{z}}}{f_z}$$
 or $f_{\overline{z}} = \mu \cdot f_z$

Measurable Riemann Mapping Theorem:

Given a measurable dilatation μ on the unit disk with $\|\mu\|_{\infty} < 1$, there is a quasiconformal $f: \mathbb{D} \to \mathbb{D}$ with this dilatation.

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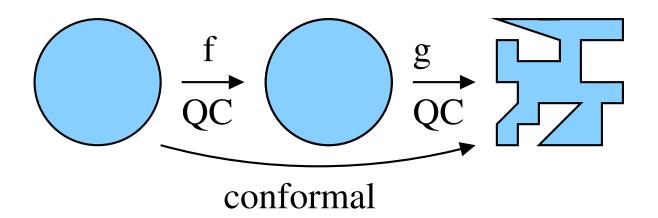
Measurable Riemann Mapping Theorem:

Given a measurable dilatation μ on the unit disk with $\|\mu\|_{\infty} < 1$, there is a quasiconformal $f: \mathbb{D} \to \mathbb{D}$ with this dilatation.

- Exact solution by power series of singular integral operators.
- Linearization can be solved by convolution with 1/z.
- Newton's method: solve linear approximation, compute new μ , repeat.
- Converges if $\|\mu\|_{\infty} \leq \epsilon_0$.

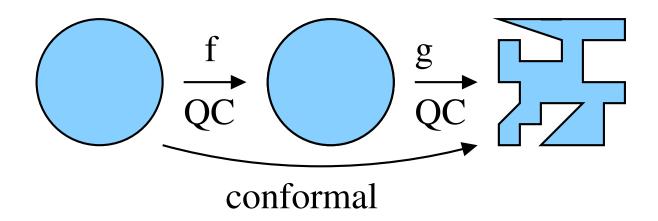
Corollary:

Given QC $g: \mathbb{D} \to \Omega$, there is $f: \mathbb{D} \to \mathbb{D}$ so that $g \circ f$ is conformal.



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Given QC $g: \mathbb{D} \to \Omega$, there is $f: \mathbb{D} \to \mathbb{D}$ so that $g \circ f$ is conformal.



Fast mapping theorem reduces to two steps:

- \bullet Find initial QC map g to polygon.
- Solve Beltrami for $\mu = \mu_g$ to get $f : \mathbb{D} \to \mathbb{D}$.

We ignore 2nd part; just find good g.

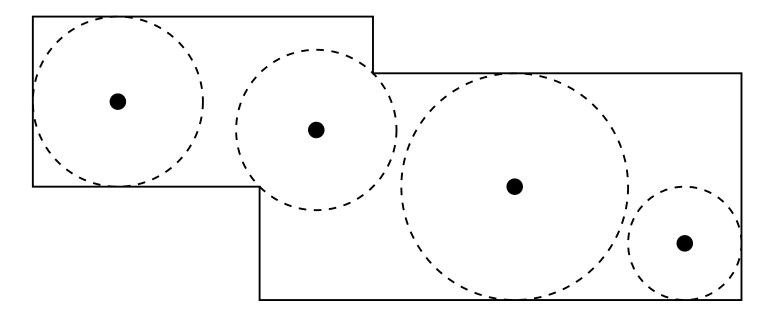
A "good" g is **fast** to compute and guaranteed **close** to correct answer.

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- fast comes from computational geometry.
- close comes from hyperbolic geometry.

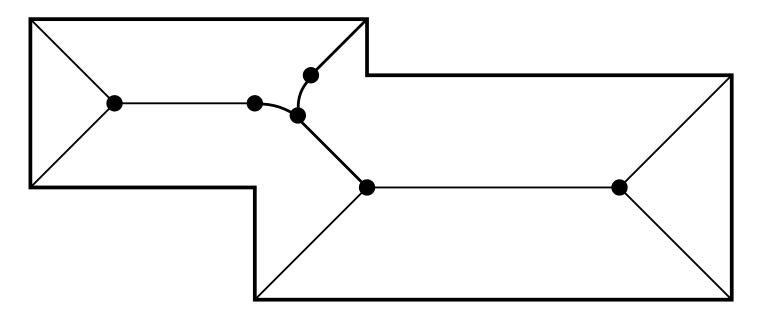
Medial axis:

centers of disks that hit boundary in at least two points.



Medial axis:

centers of disks that hit boundary in at least two points.



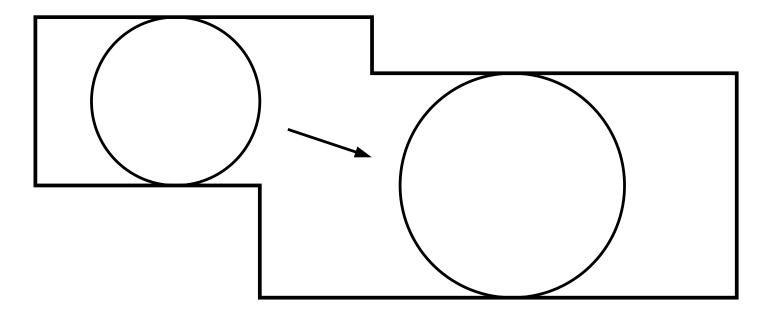
Medial axis of a polygon is a finite tree.

Computable in O(n), Chin-Snoeyink-Wang (1999).

Related to Voronoi diagrams: divides polygon according to nearest edge.

Medial axis:

centers of disks that hit boundary in at least two points.



Claim: there is a "natural" choice of conformal map between any two medial axis disks.

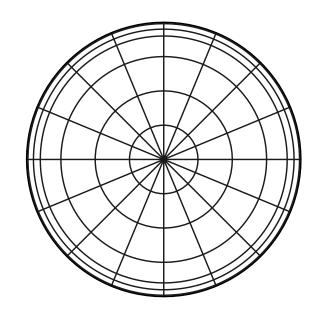
A Möbius transformation is a map of the form

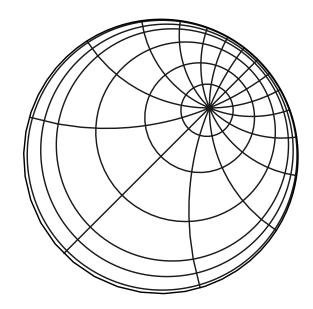
$$z \to \frac{az+b}{cz+d}$$
.

Conformally maps disks to disks (or half-planes).

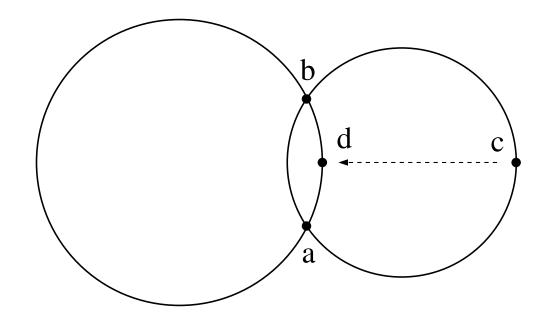
Form a group under composition.

Uniquely determined by images of 3 distinct points.





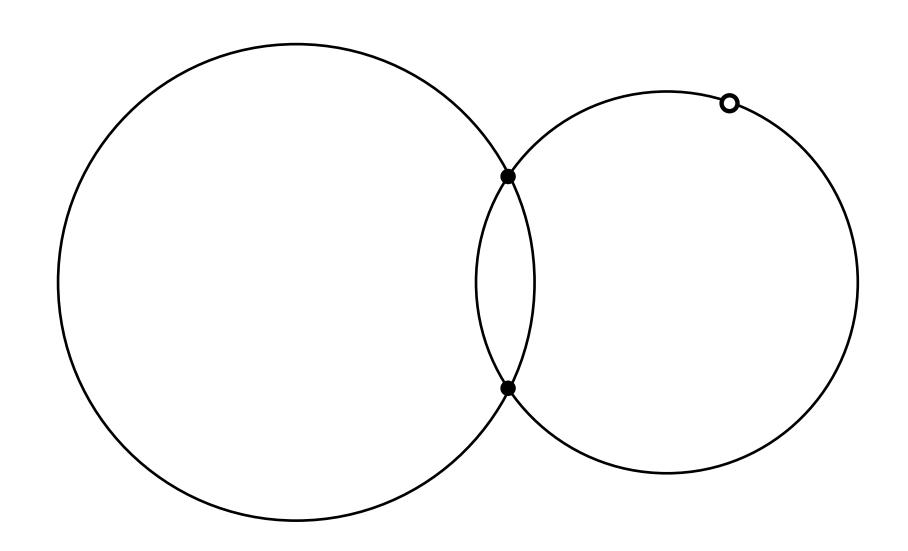
Intersecting circles:

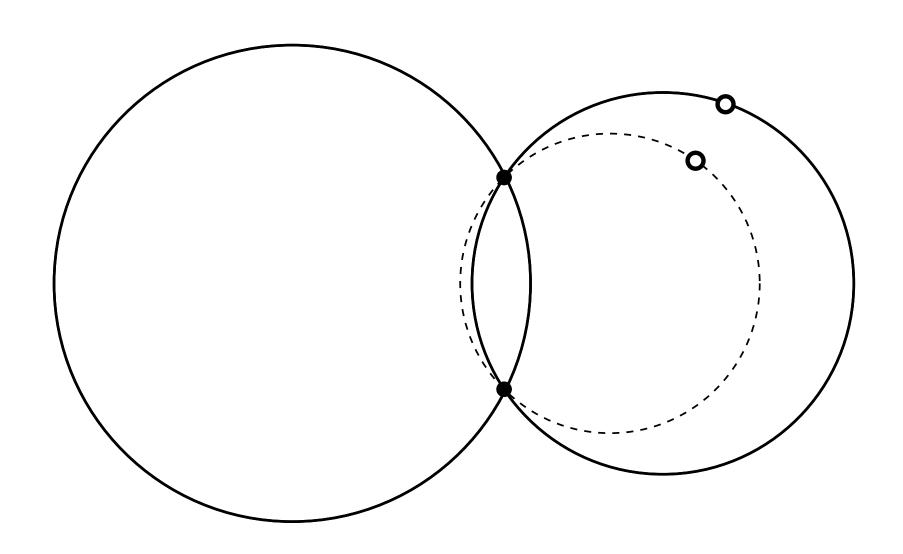


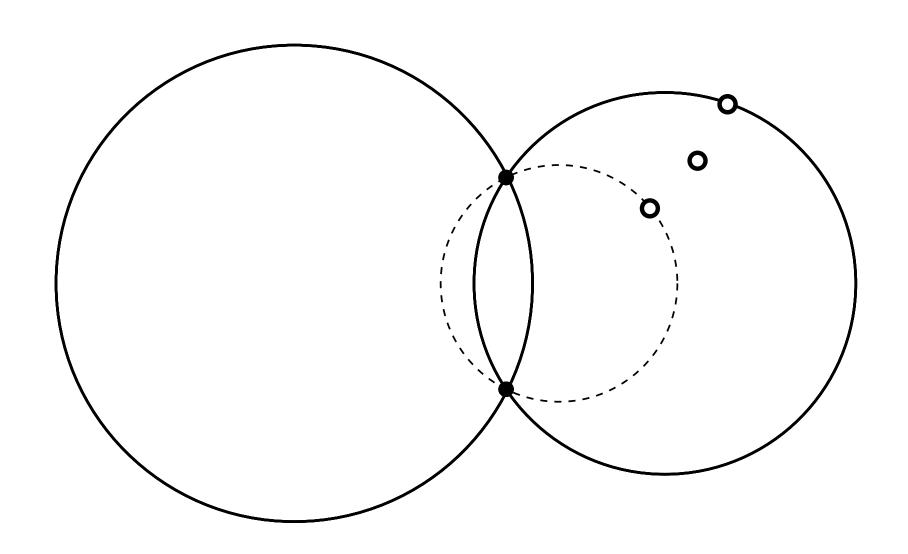
Fix intersection points a, b and map $c \to d$ as shown.

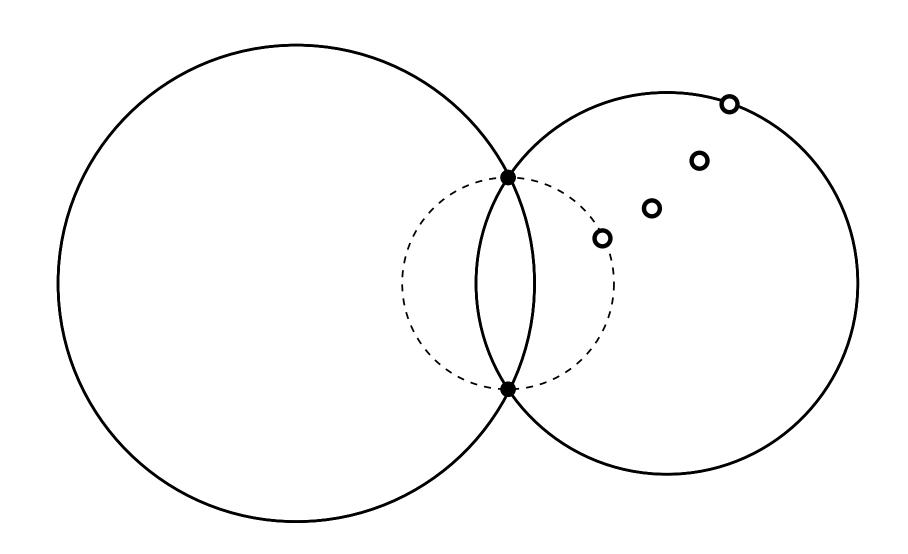
Determines unique Möbius map between disks.

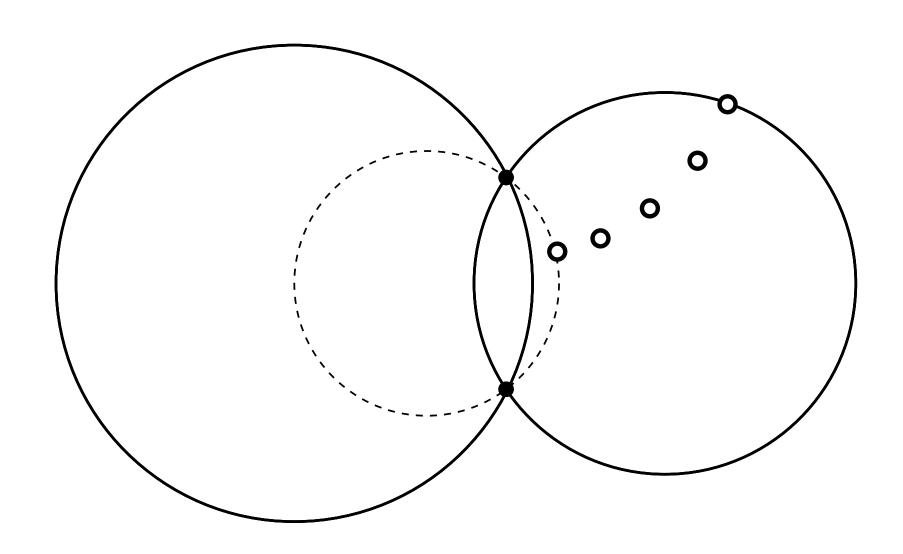
Part of 1-parameter symmetric family fixing a, b.

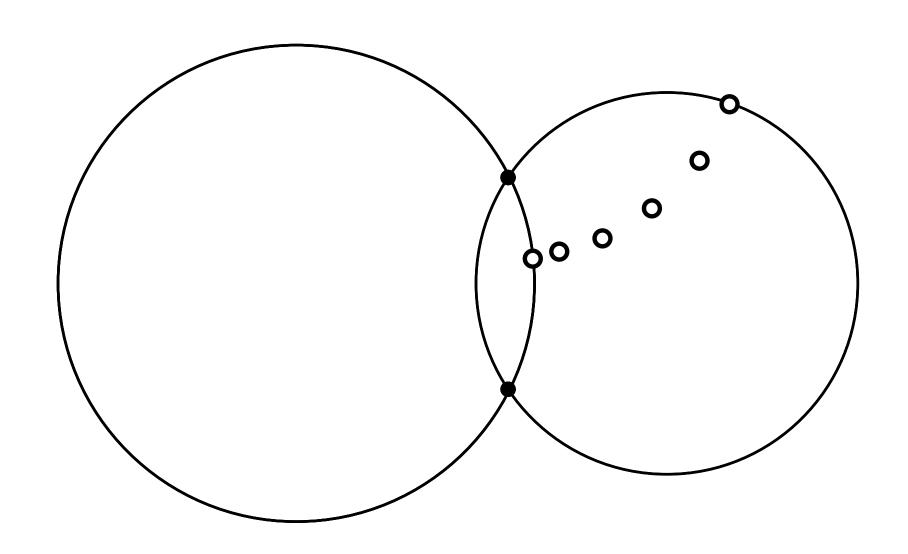


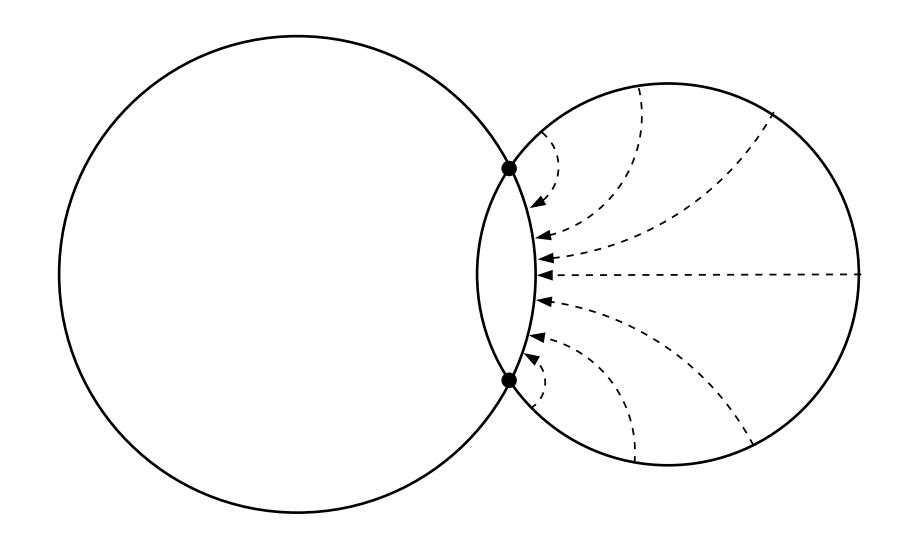






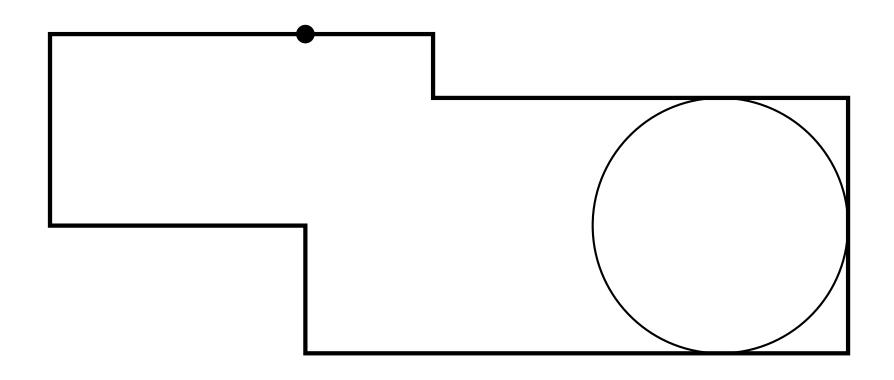




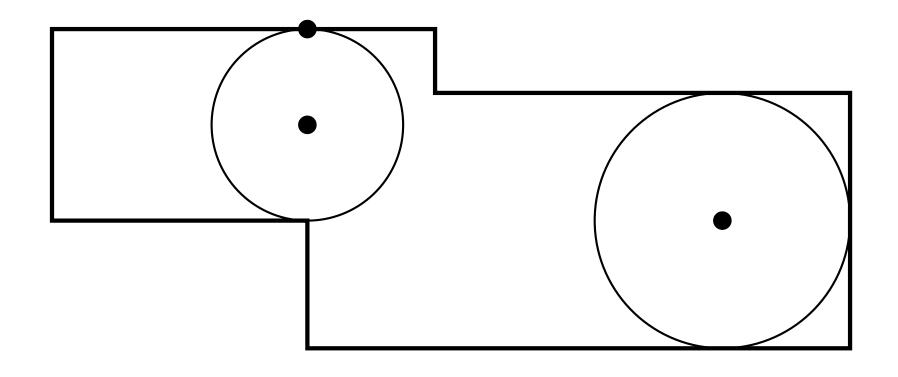


Points follow circular paths, perpendicular to boundary.

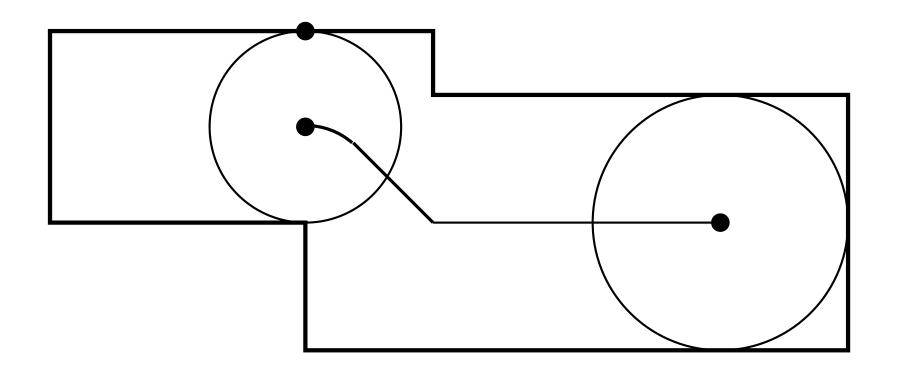
How does this give a map from polygon P to a circle?



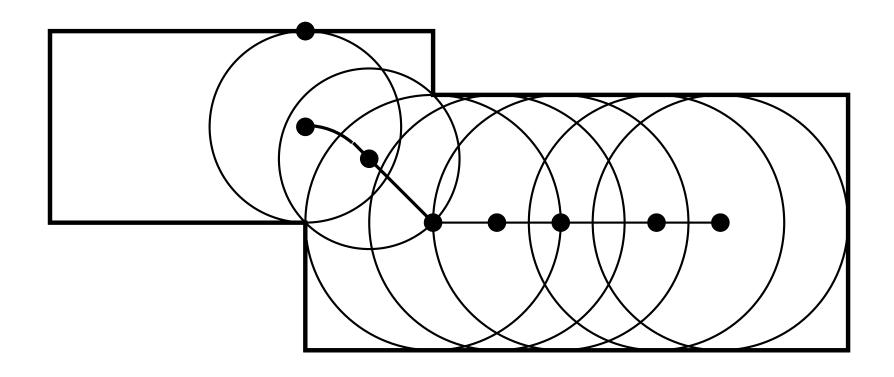
 \bullet Fix a "root" MA disk D.

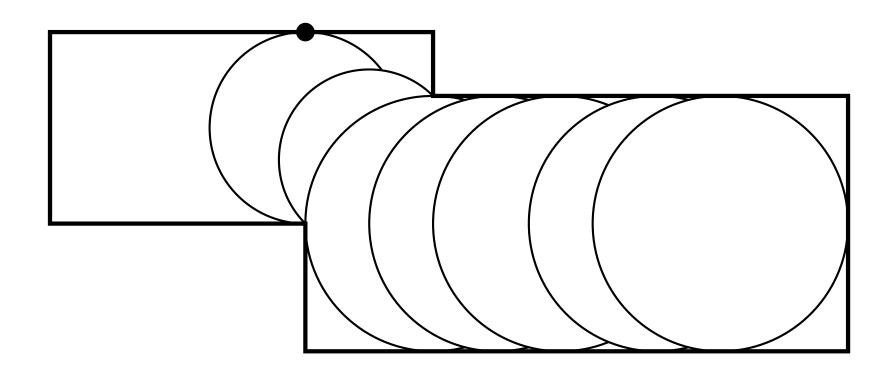


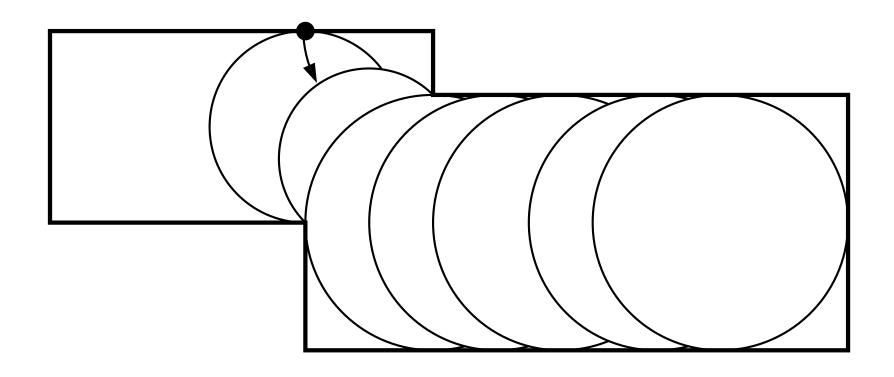
• For any $z \in P$, take MA disk D_z touching z.

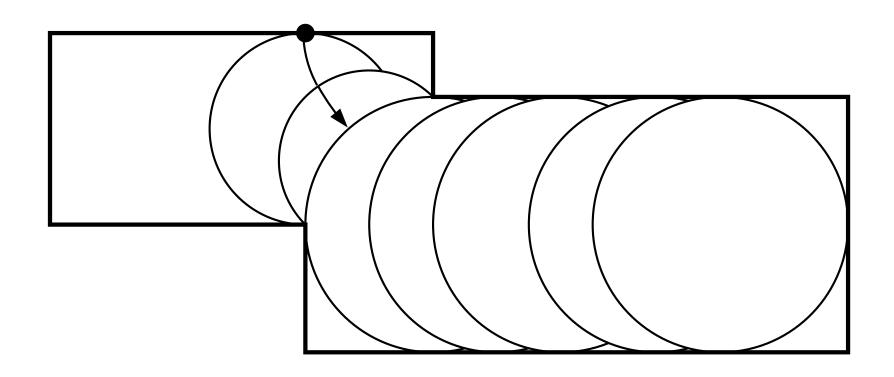


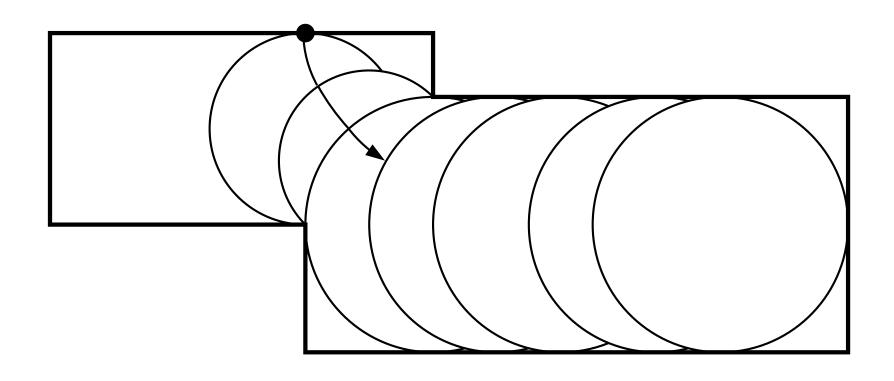
• Connect D_z to D on MA.

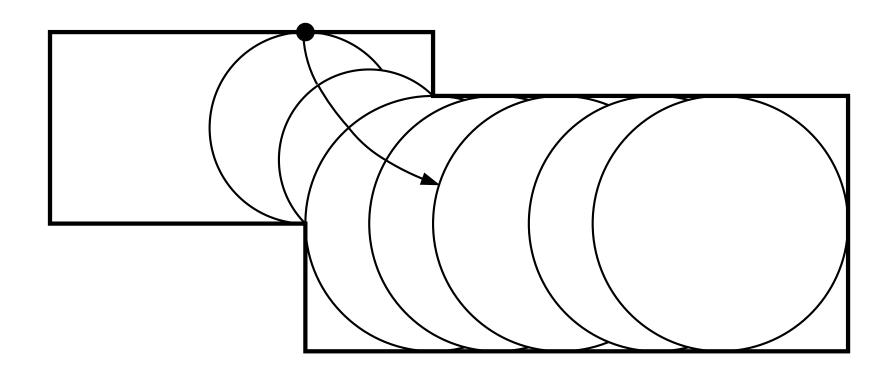


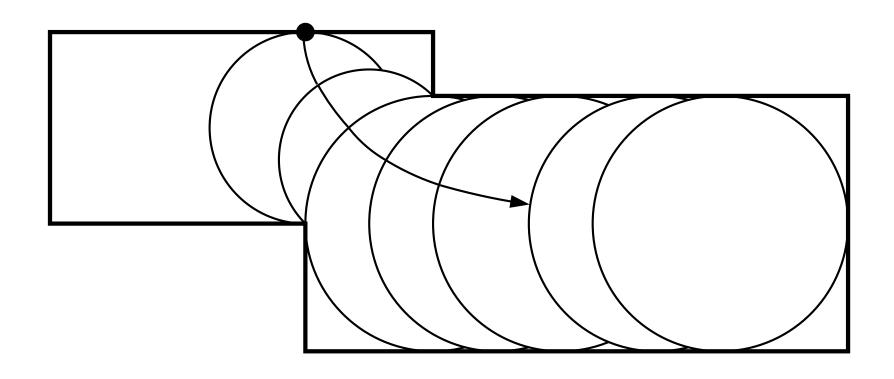


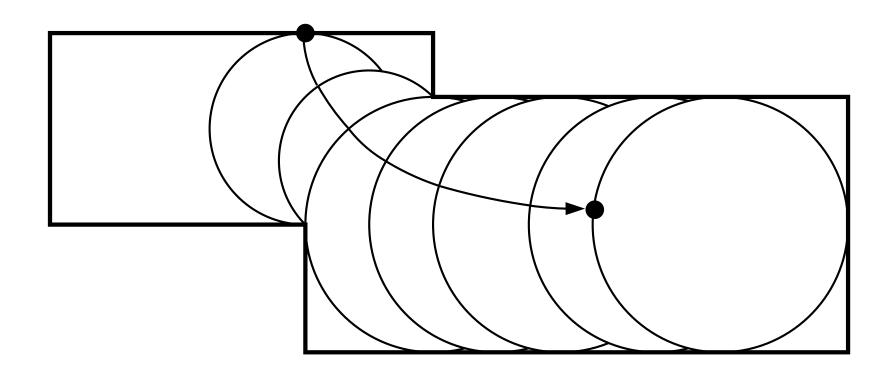


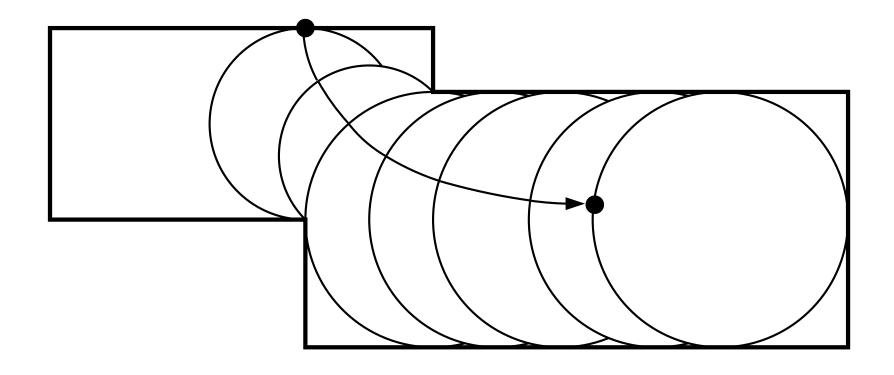








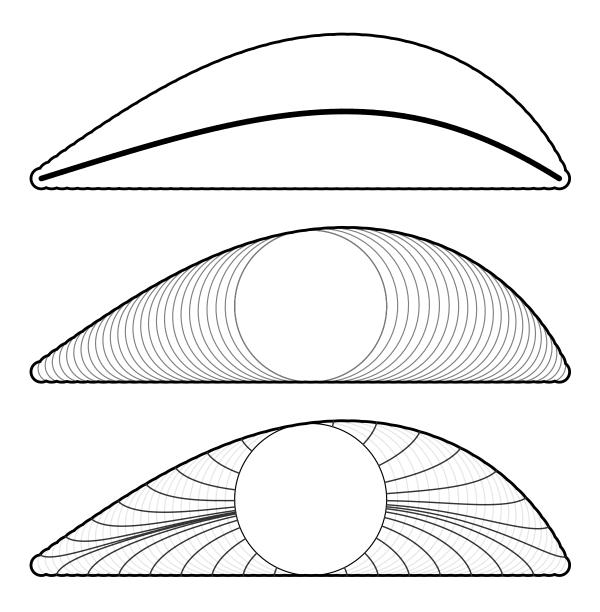


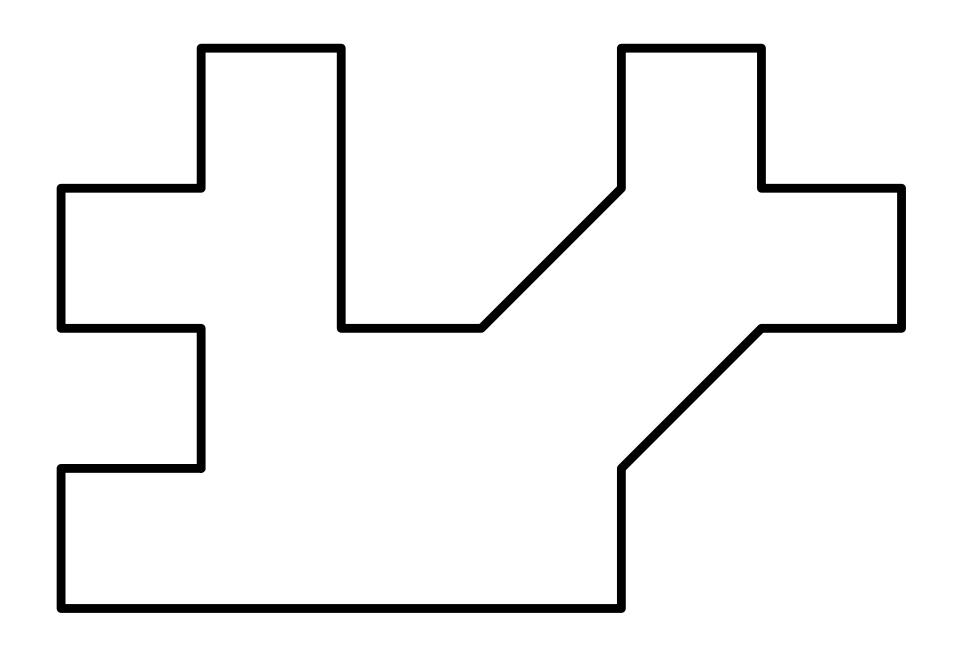


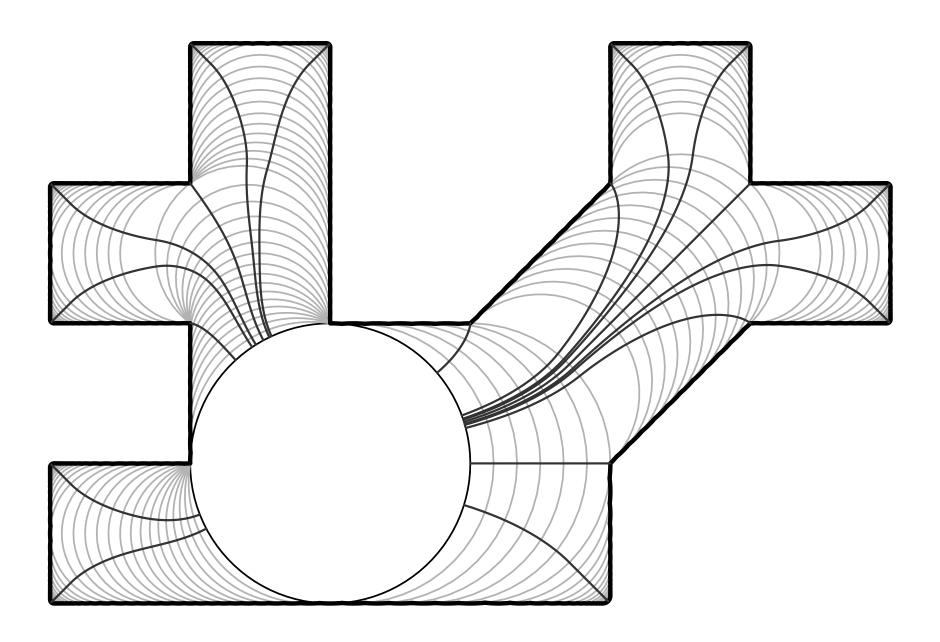
We discretize only to draw picture.

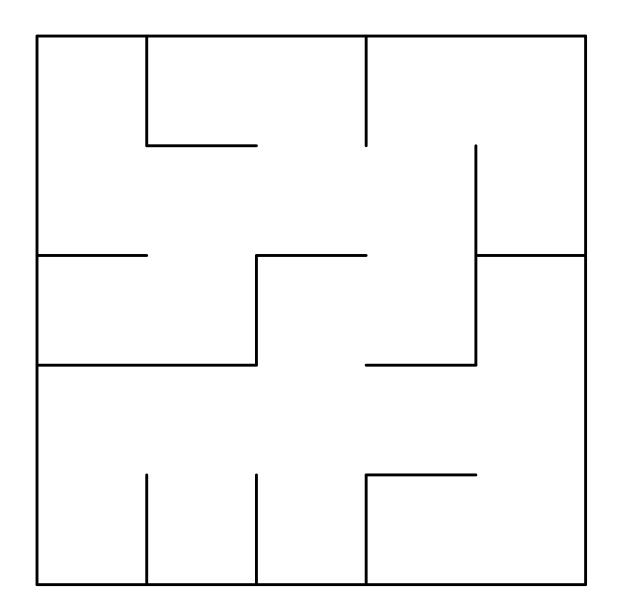
Limiting map has **formula** in terms of medial axis.

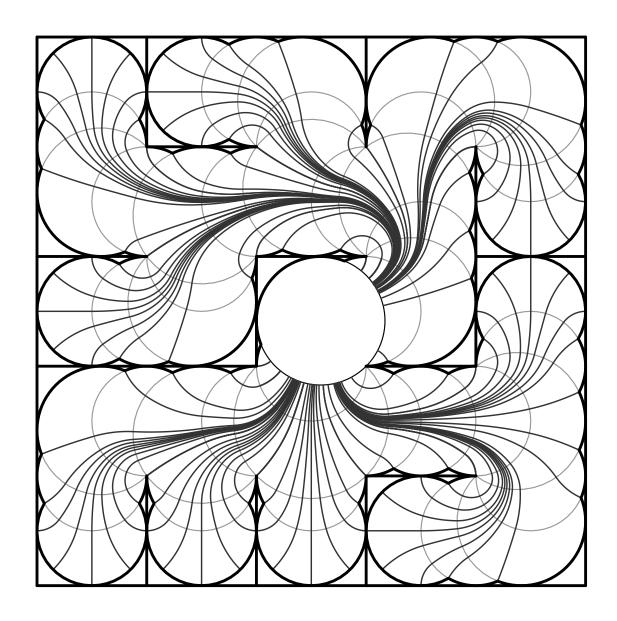
Similar flow for any simply connected domain.





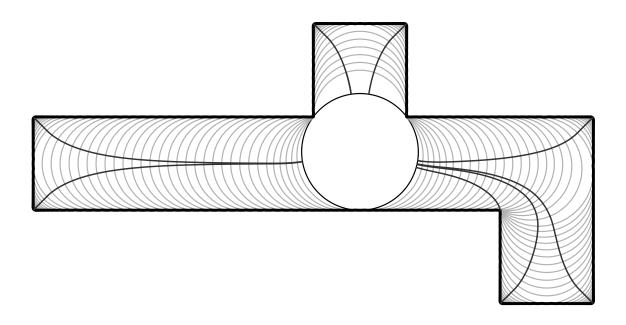






Theorem: Mapping all n vertices takes O(n) time.

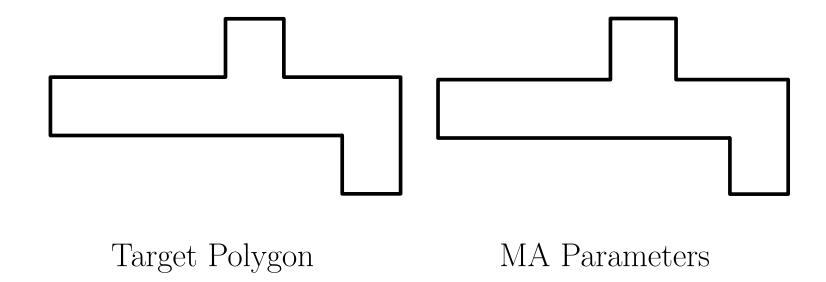
Uses linear time computation of MA (Chin-Snoeyink-Wang) and book-keeping with cross ratios.



How close is medial axis map to conformal map?

How close is medial axis map to conformal map?

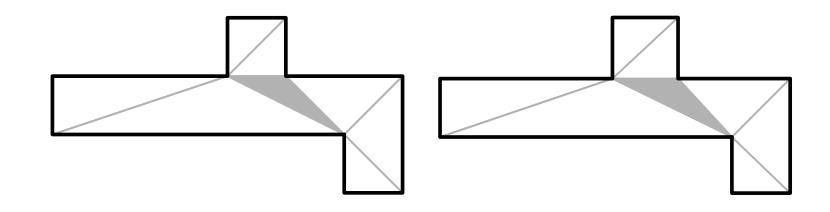
Use "MA-parameters" in Schwarz-Christoffel formula.



Looks pretty close. What is QC distance?

How close is medial axis map to conformal map?

Use "MA-parameters" in Schwarz-Christoffel formula.



Target Polygon

MA Parameters

Looks pretty close. What is QC distance?

The most distorted triangle is shaded. Here K = 1.24.

Theorem: Medial axis map always gives QC-error K < 8.

Theorem: Medial axis map always gives QC-error K < 8.

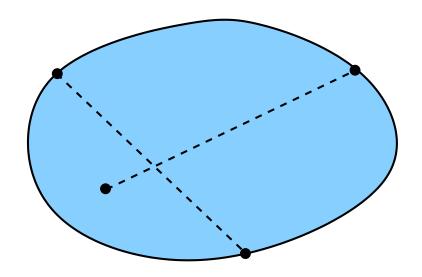
Why is this theorem true?

Theorem: Medial axis map always gives QC-error K < 8.

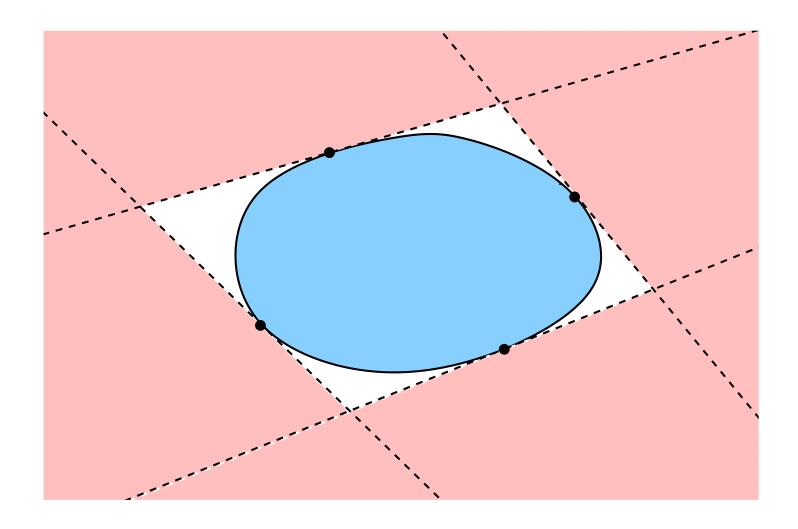
Why is this theorem true?

Short answer: convex sets in hyperbolic 3-space

Usual definition of convex: contains geodesic between any two points.



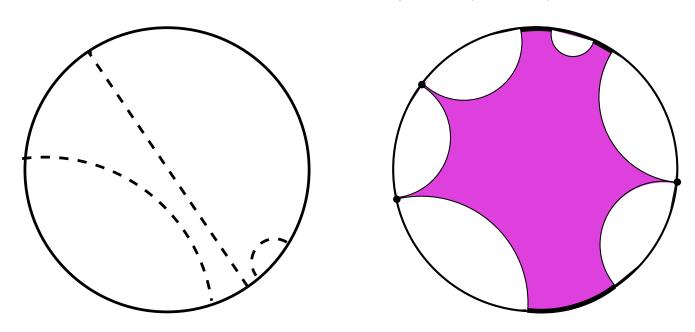
More useful for us: complement is a union of half-spaces.



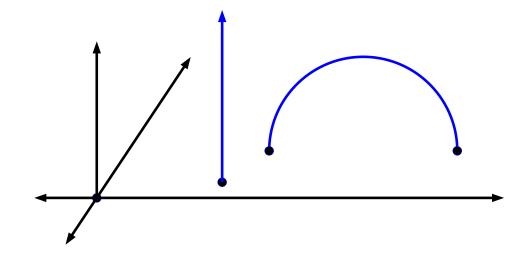
Hyperbolic metric on disk given by

$$d\rho = \frac{ds}{1 - |z|^2} \simeq \frac{ds}{\operatorname{dist}(z, \partial D)}.$$

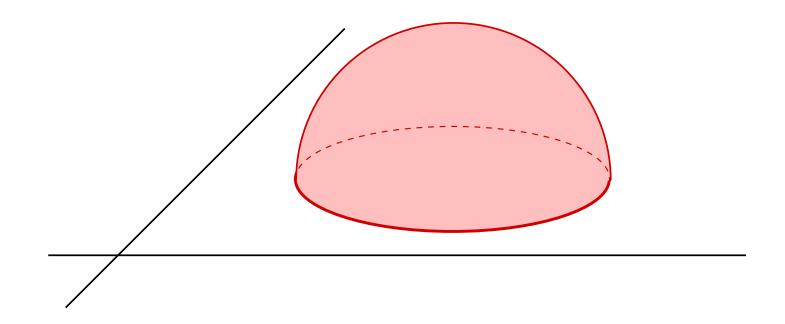
- Geodesics are circles perpendicular to boundary.
- Shaded region is hyperbolically convex.
- Metric transfers via conformal maps to other domains.
- For simply connected regions $d\rho \simeq ds/\mathrm{dist}(z,\partial\Omega)$.



In the upper half-space $\mathbb{R}^3_+ = \{(x, y, t) : t > 0\}$, metric is $d\rho = ds/2t$.



Geodesics in \mathbb{R}^3_+ are vertical rays or semi-circles perpendicular to \mathbb{R}^2 .

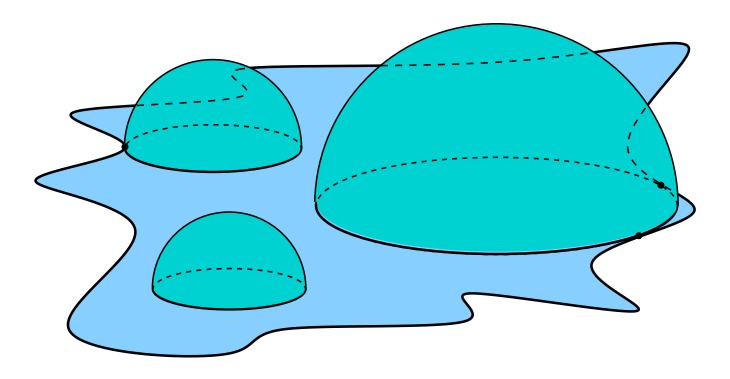


In \mathbb{R}^3_+ , a hyperbolic half-space = hemisphere.

Given $\Omega \subset \mathbb{R}^2$, compute hyperbolic convex hull its complement.

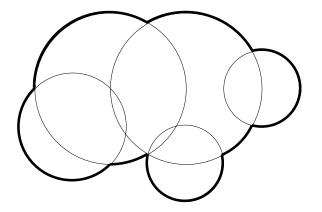
Easier to visualize the complement of convex hull = union of hemispheres.

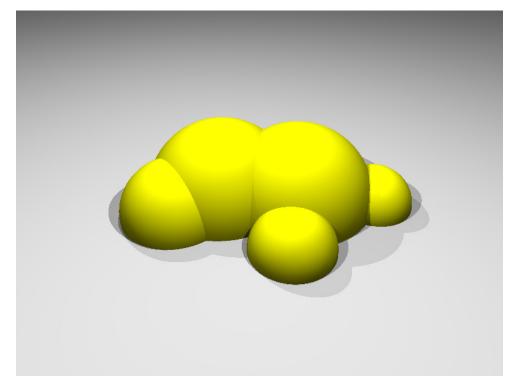
 $Dome(\Omega)$ is union of hemi-spheres with base disks in Ω .



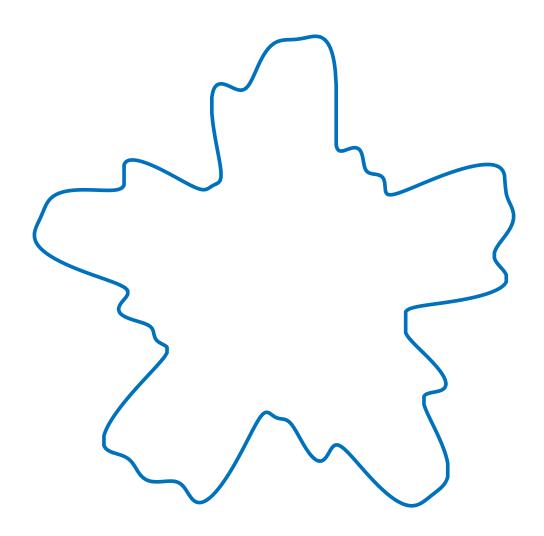
Region above dome is intersection of half-spaces, hence convex.

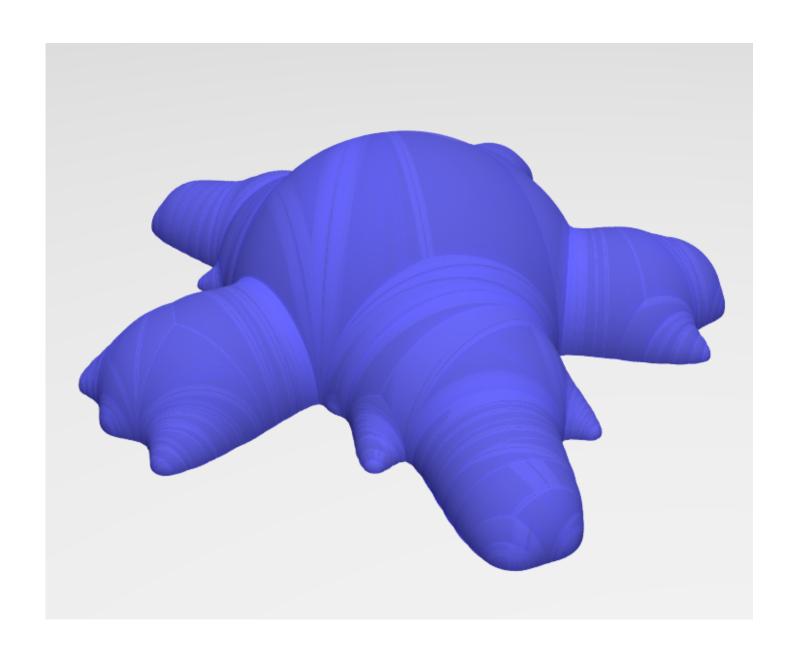
Upper boundary S of dome is a surface in \mathbb{R}^3_+ with $\partial S = \partial \Omega$.

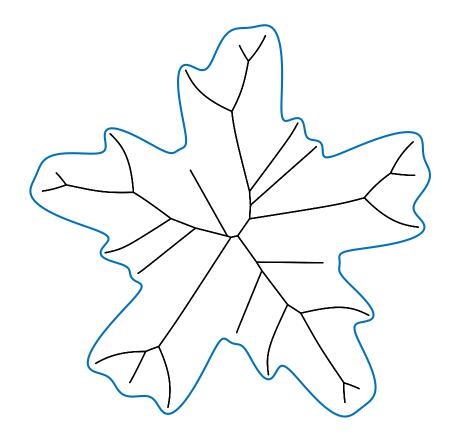




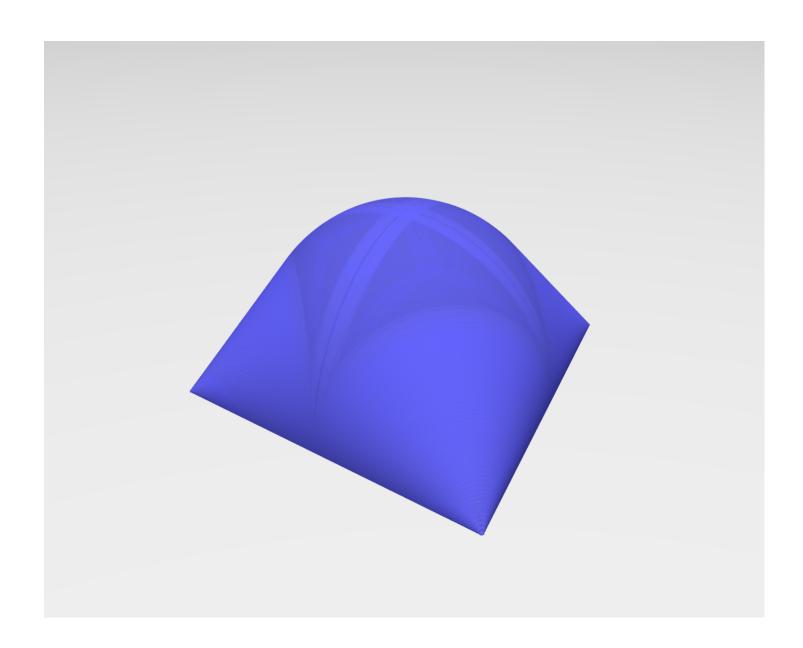
Finite dome

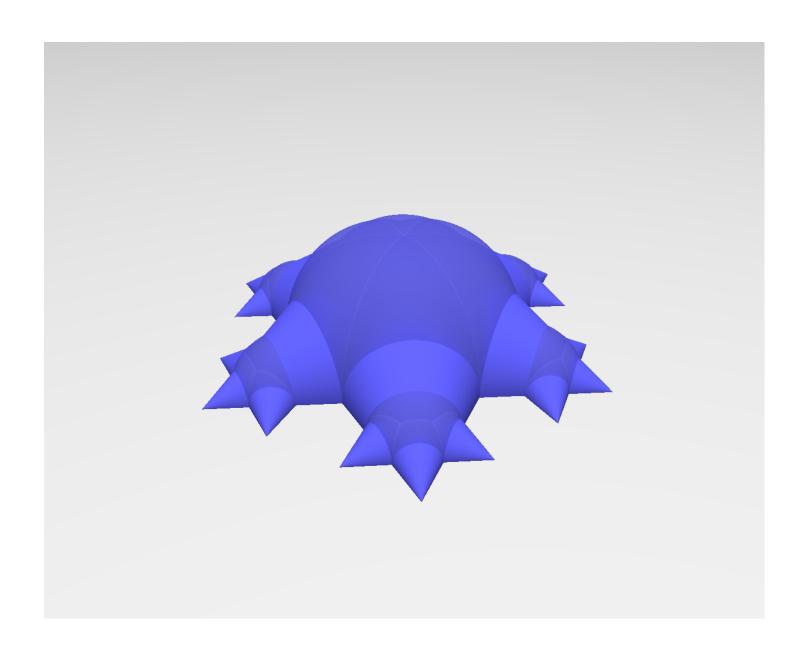






The medial axis. Equidistant from at least two boundary points. Corresponding hemispheres give the dome.





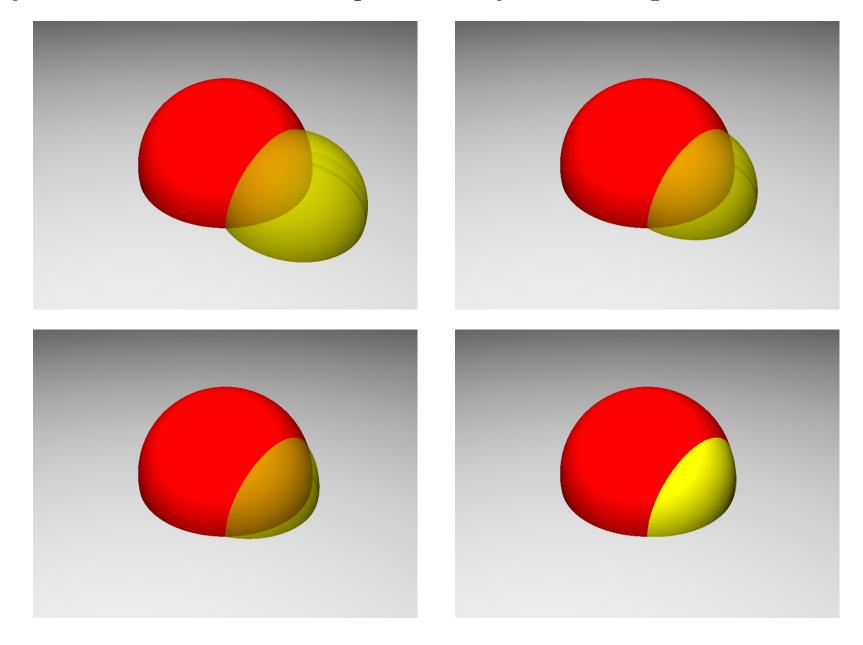
Thm: Simply connected domes are isometric to hyperbolic disk.

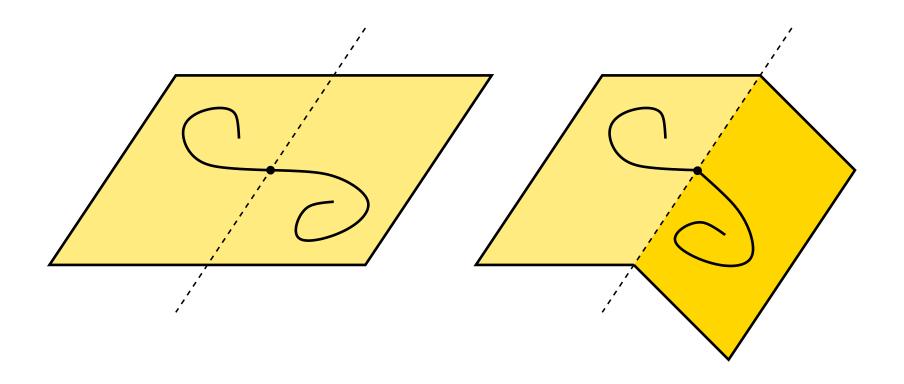
We assume dome has hyperbolic path metric.

- Prove for finite unions of disks.
- Every dome is a limit of finite domes.
- Limit of isometries is an isometry.

Isometry on boundary Γ defines a map Γ to circle.

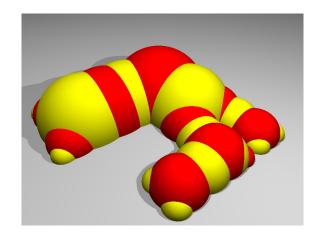
Every dome has conformal map to disk by "flattening".

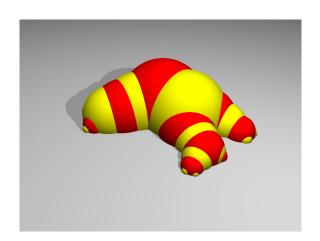


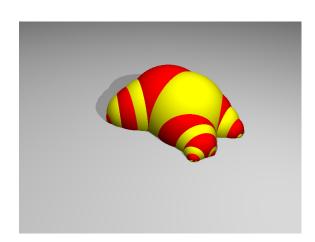


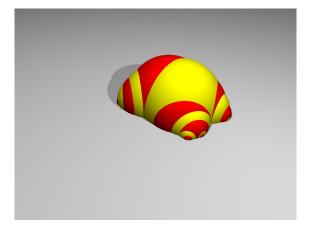
Folding plane along geodesic does not change length.

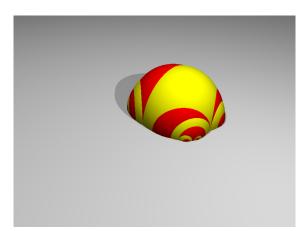
Pleated surface (folded along disjoint geodesics) = Flat plane

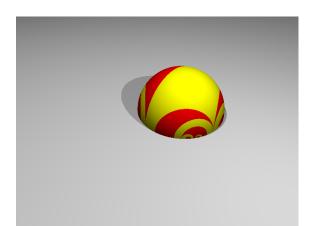


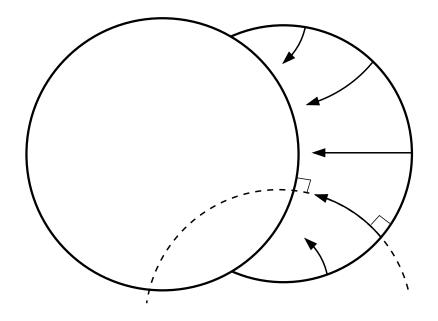






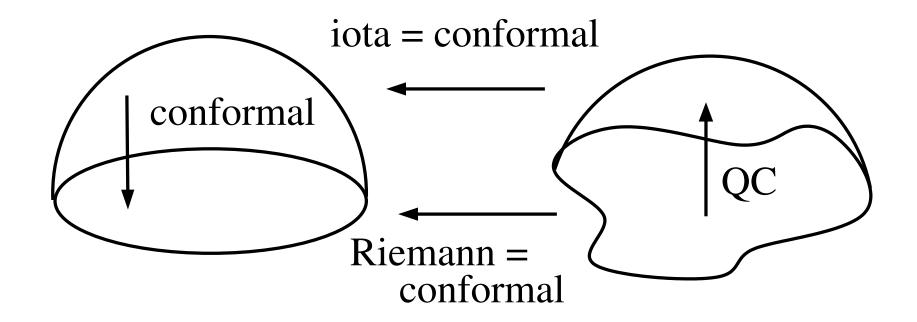






Medial axis map = boundary of flattening map (iota)

= boundary of conformal map of dome to hemisphere



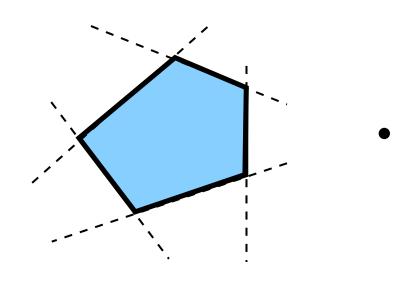
Iota = conformal from dome to disk.

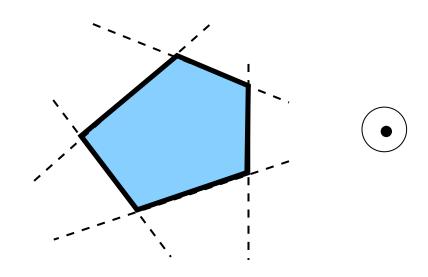
Medial axis flow = boundary values of iota

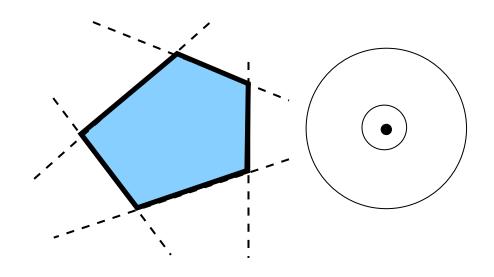
Claim: There is QC map base \rightarrow dome fixing boundary pointwise.

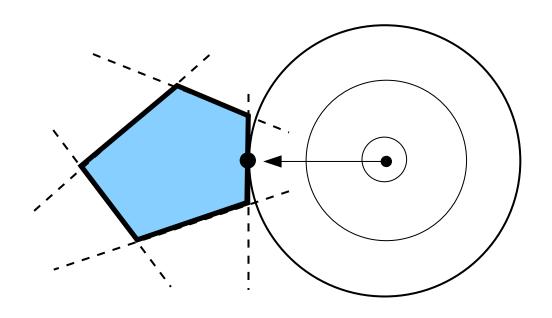
Implies that medial axis map has QC extension $\Omega \to \mathbb{D}$.

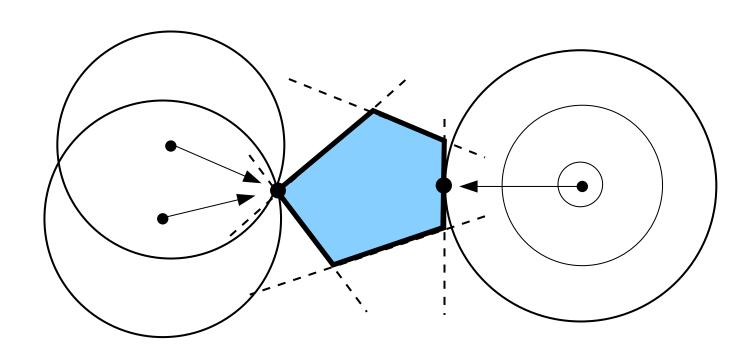
Claim is proven using nearest point projection onto convex sets.

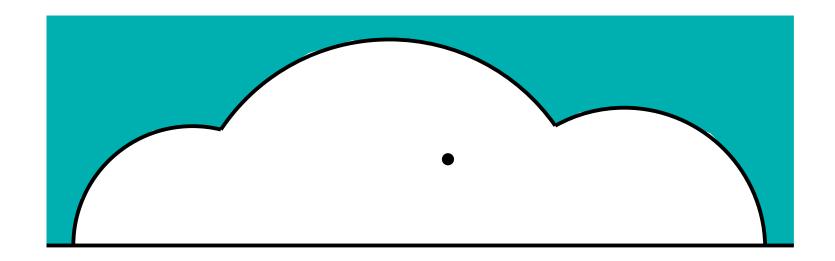










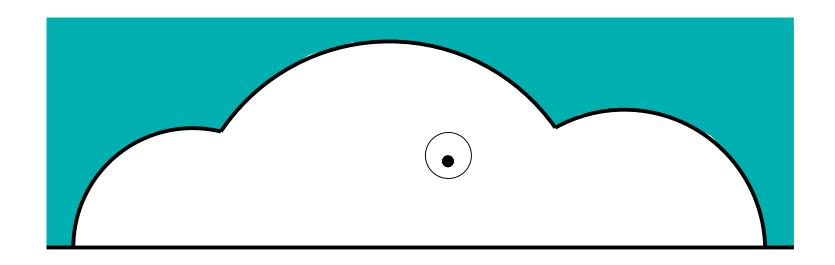


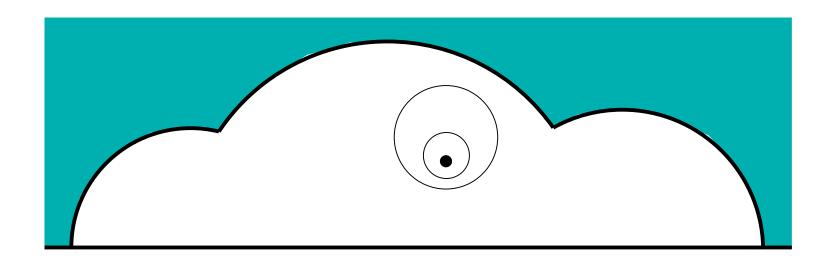
Region below dome is union of hemispheres

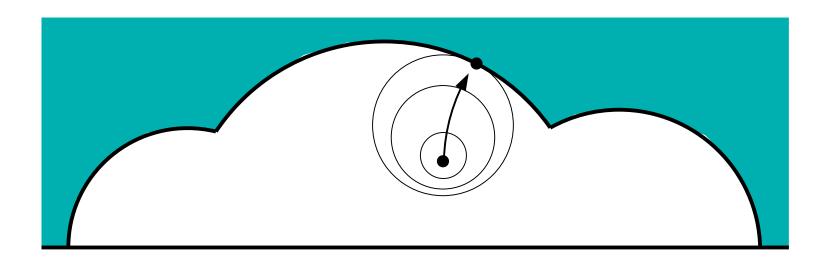
Hemispheres = hyperbolic half-spaces.

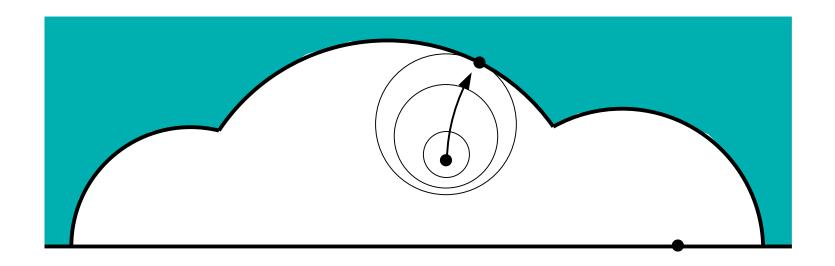
Region above dome is hyperbolically convex.

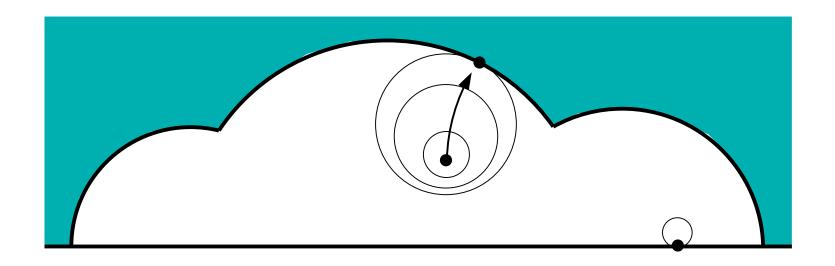
Consider nearest point retraction onto this convex set.

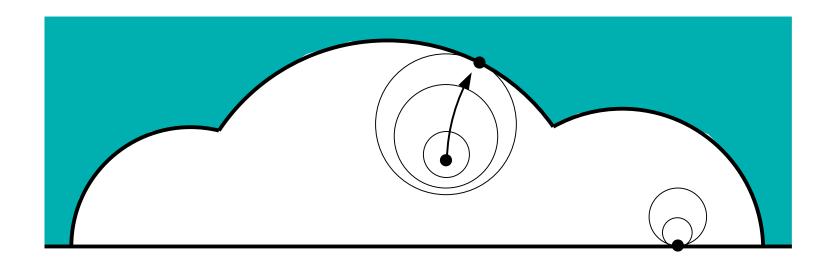


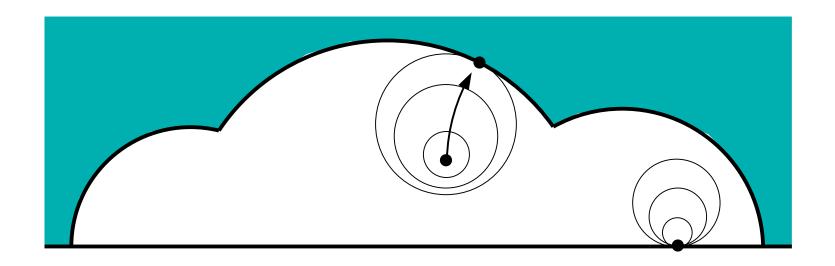


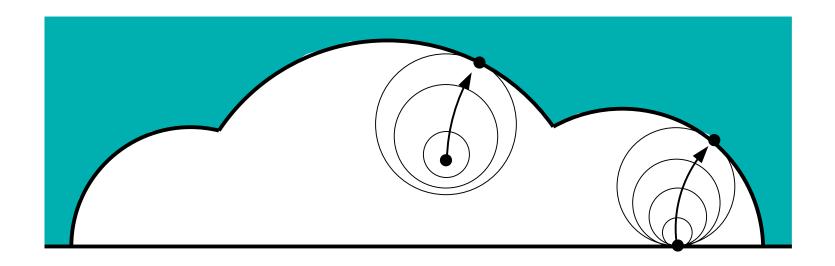


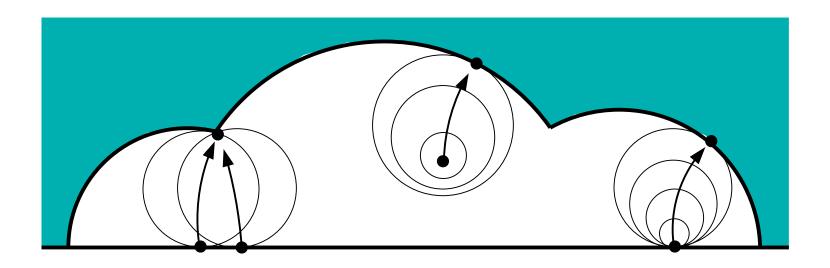




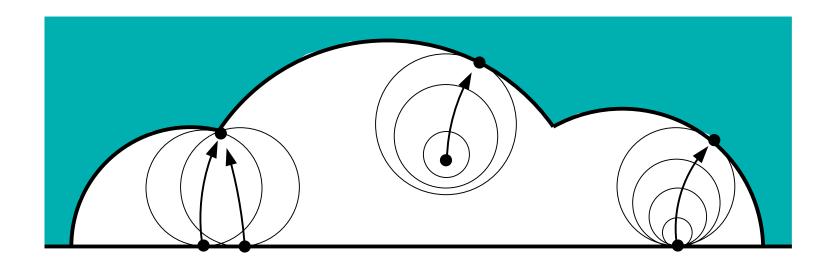








Need not be a homeomorphism, but ...

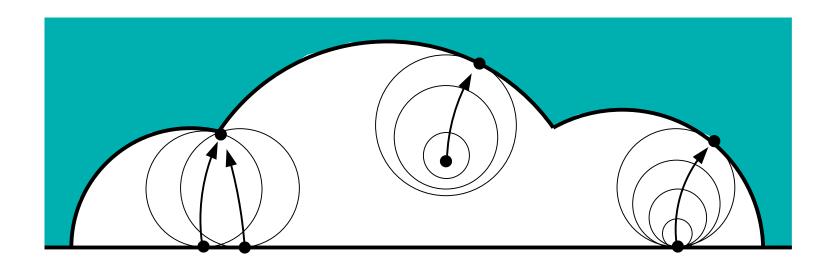


Need not be a homeomorphism, but it is a quasi-isometry

$$\frac{1}{A} \le \frac{\rho(R(x), R(y))}{\rho(x, y)} \le A, \quad \text{if } \rho(x, y) \ge B.$$

i.e., R is bi-Lipschitz on large scales.

Metrics are hyperbolic metrics on Ω and S.



"Smoothing" gives K-QC map fixing boundary points.

Sullivan's convex hull theorem: K is independent of domain.

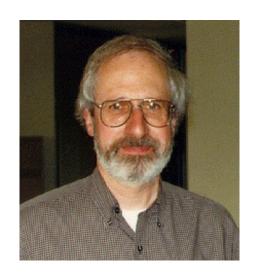
Dennis Sullivan, David Epstein and Al Marden, C.B.







David Epstein



Al Marden

Dennis Sullivan proved this assuming invariance under a group of Möbius transformations. This was used by William Thurston to prove certain 3-manifolds have a hyperbolic metric.

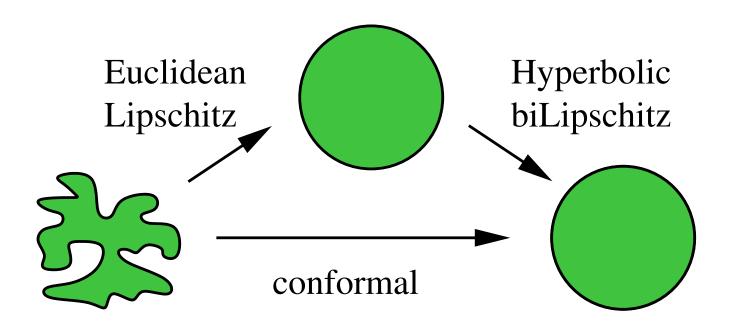
Epstein and Marden extended to general simply connected Ω . $K \approx 85$.

Best value unknown, but 2.1 < K < 7.82.

Application: factorization: Riemann map $f = h \circ g$ where

• $g: \Omega \to \mathbb{D}$ is Lipschitz in Euclidean path metrics,

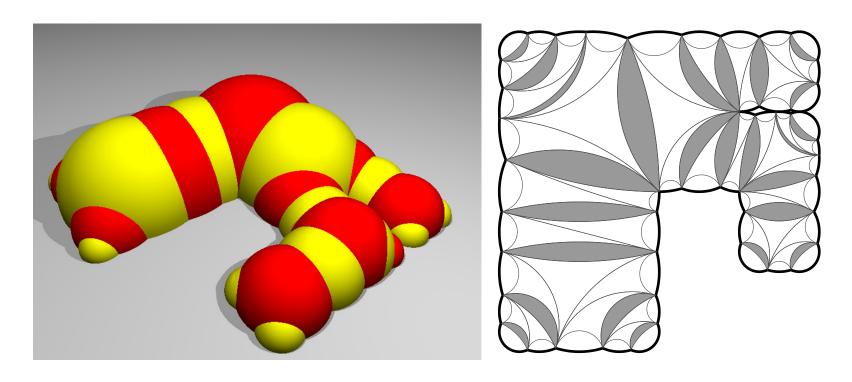
• $h: \mathbb{D} \to \mathbb{D}$ is biLipschitz in hyperbolic metric



Cor: Any simply connected domain can be mapped 1-1, onto a disk D by a contraction for the internal path metric.

Application: Angle scaling:

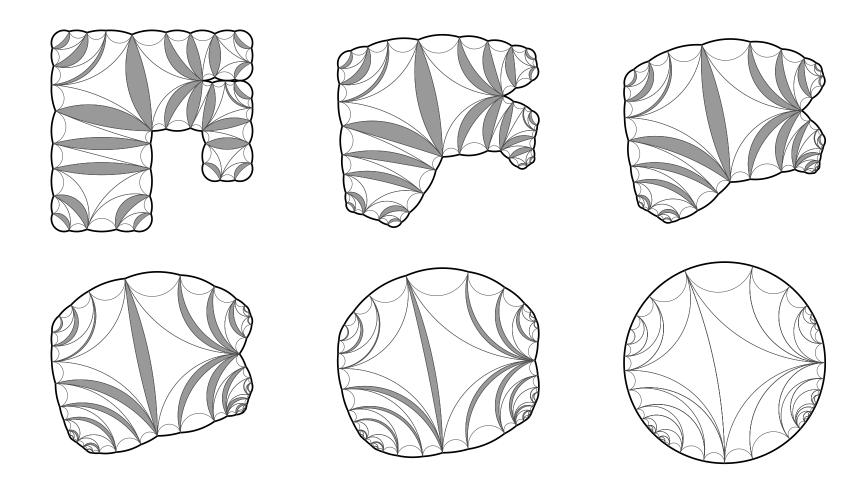
Crescents in base can map to folding geodesics on surface.



Gray collapses to bending lines, "width = angle".

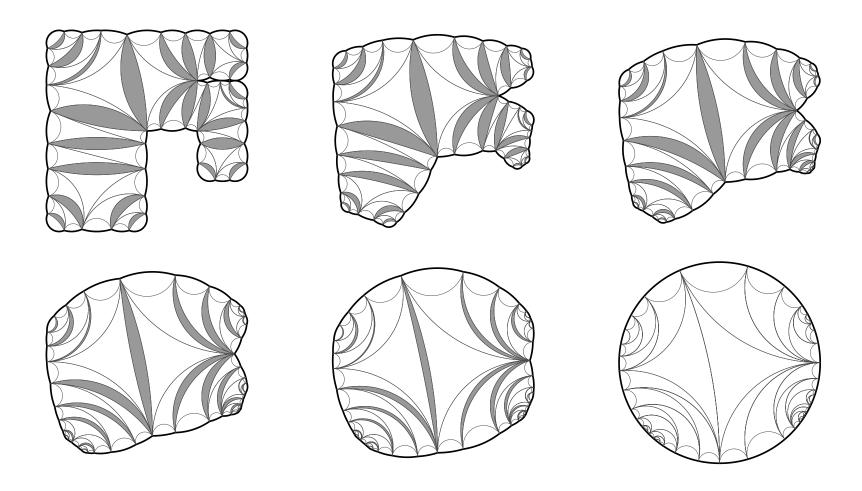
White maps isometrically to dome.

Discrete Riemann map: collapses crescents (gray), Möbius elsewhere (white).



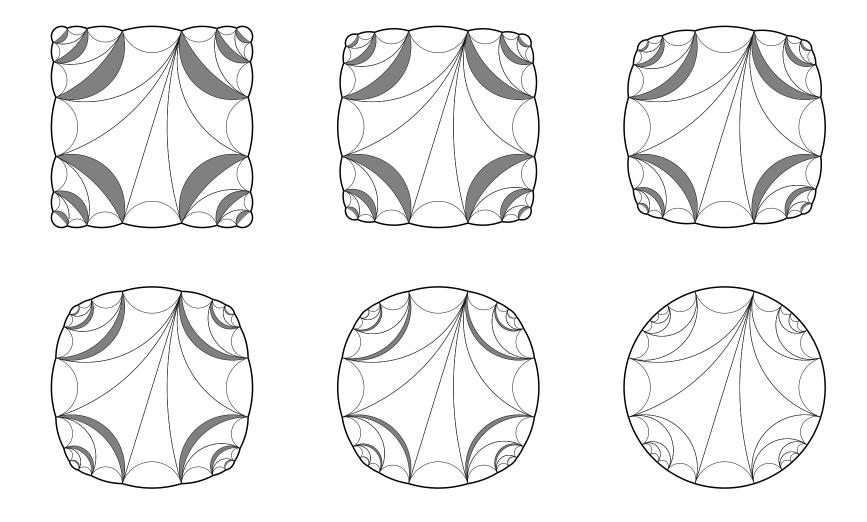
Angle scaling family - crescent angles decrease "Morphs" region to disk.

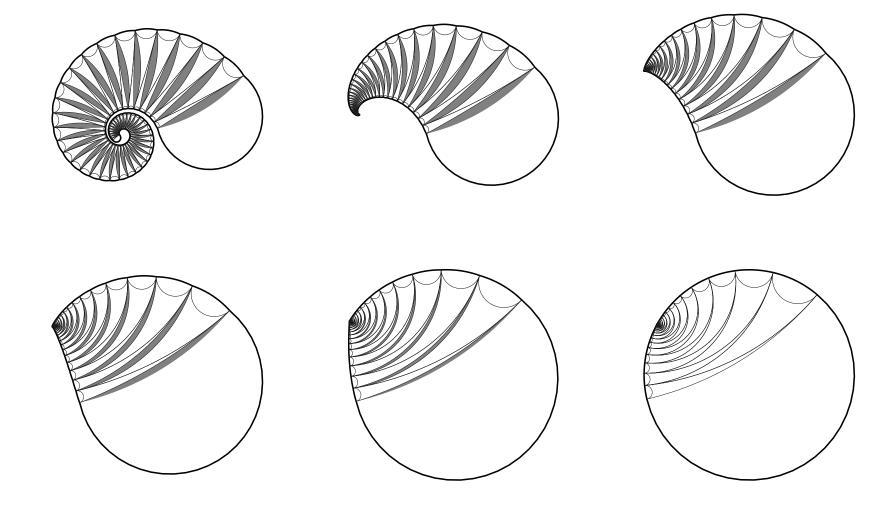
Reduces solving Beltrami equation to case of small dilatations.

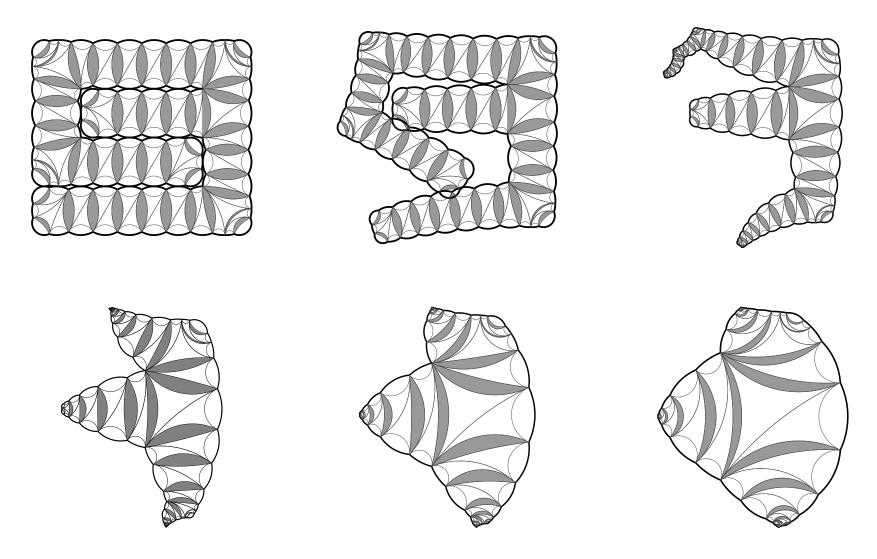


Riemann map approximated by cutting into simple pieces and rearranging. Gray pices collapse orthogonally.

White pieces map by Möbius transformations.

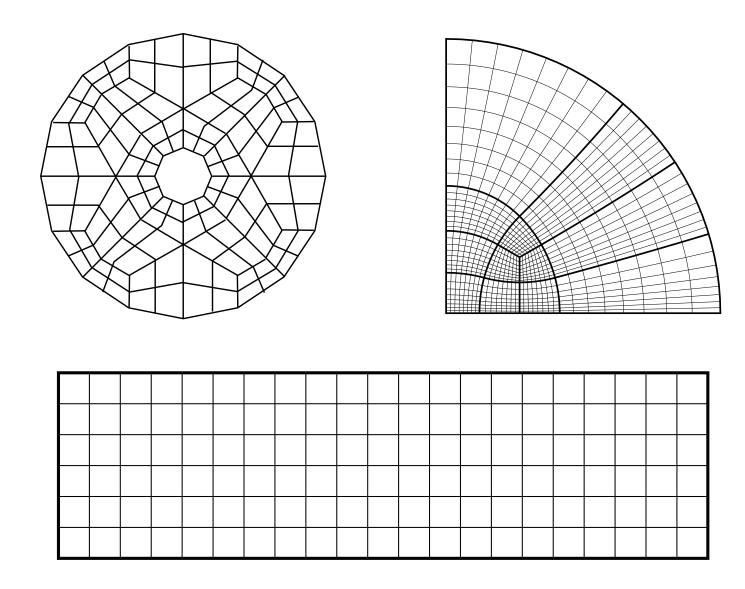






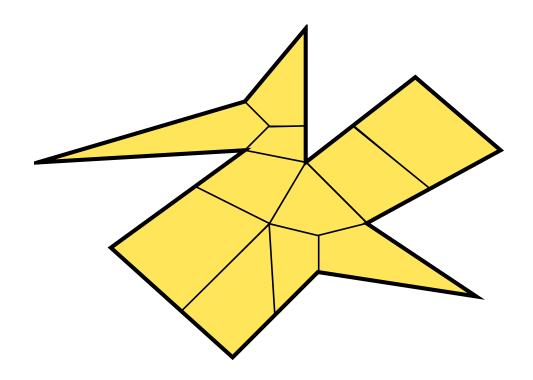
Intermediate regions need not be planar.

Application: Quadrilateral meshes



Marshall Bern and David Eppstein (2000) proved:

- n-gons have O(n) quad-mesh with angles $\leq 120^{\circ}$.
- $O(n \log n)$ work.
- Regular hexagon (and Euler's formula) shows 120° is sharp.



Bern asked: can we bound angles from below?

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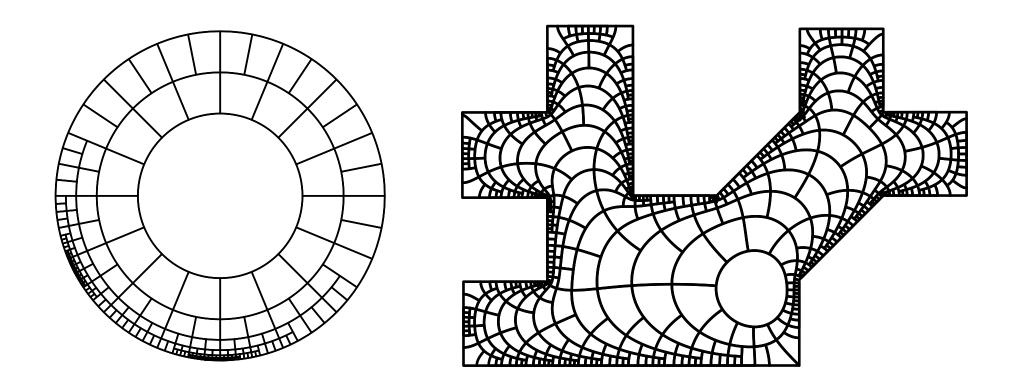
Theorem: Every n-gon has O(n) quad-mesh with all angles $\leq 120^{\circ}$ and new angles $\geq 60^{\circ}$. O(n) work.

Bern asked: can we bound angles from below?

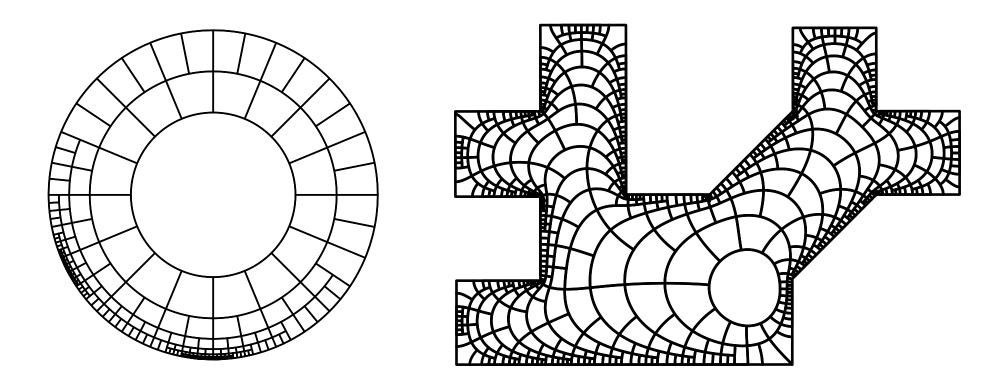
Theorem: Every n-gon has O(n) quad-mesh with all angles $\leq 120^{\circ}$ and new angles $\geq 60^{\circ}$. O(n) work.

Original angles $< 60^{\circ}$ remain unchanged. 60° is sharp.

Idea of proof: transfer mesh from disk



Idea of proof: transfer mesh from disk



More about this next time.

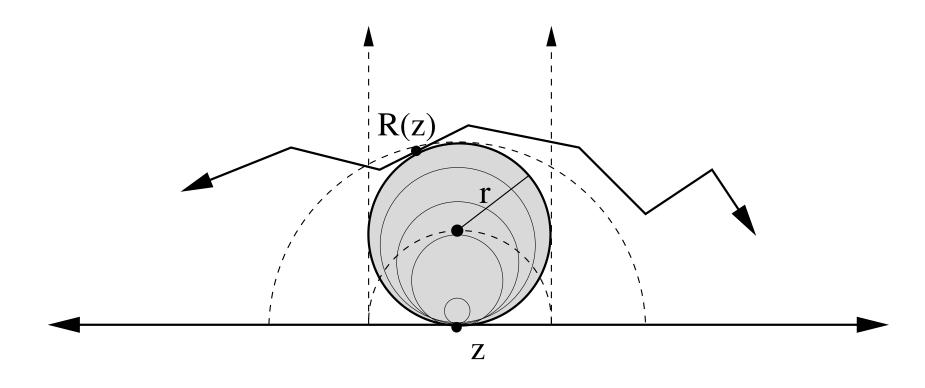


Sketch of proof that R is quasi-isometry

One direction: R is Lipschitz.

Other direction: R^{-1} is Lipschitz at distances ≥ 1 .

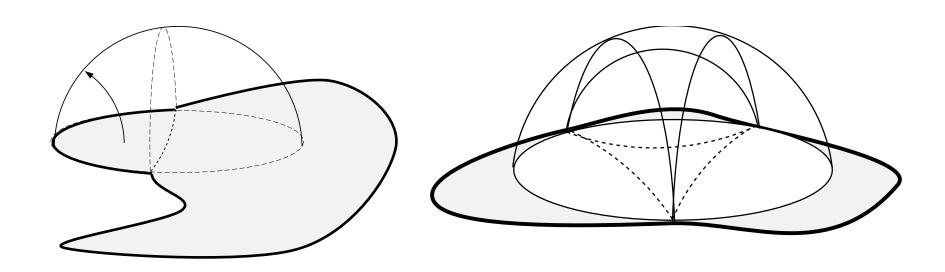
Fact 1: If $z \in \Omega$, $\infty \notin \Omega$, $r \simeq \operatorname{dist}(z, \partial \Omega) \simeq \operatorname{dist}(R(z), \mathbb{R}^2) \simeq |z - R(z)|.$



Fact 2: R is Lipschitz.

- Ω simply connected $\Rightarrow d\rho \simeq |dz|/\mathrm{dist}(z,\partial\Omega)$.
- $z \in D \subset \Omega$ and $R(z) \in Dome(D) \Rightarrow z$ in hyperbolic convex hull of $\partial \Omega \cap \partial D$ in D.

$$\Rightarrow \operatorname{dist}(z, \partial \Omega) / \sqrt{2} \leq \operatorname{dist}(z, \partial D) \leq \operatorname{dist}(z, \partial \Omega)$$
$$\Rightarrow \rho_{\Omega}(z) \simeq \rho_{D}(z) = \rho_{\operatorname{Dome}}(R(z)).$$



Fact 3: $\rho_S(R(z), R(w)) \leq 1 \Rightarrow \rho_{\Omega}(z, w) \leq C$.

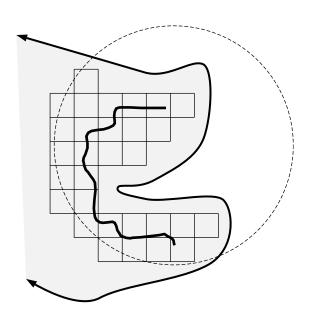
Suppose $dist(R(z), \mathbb{R}^2) = r$.

Suppose γ is geodesic on dome from R(z) to R(w).

$$\Rightarrow \operatorname{dist}(\gamma, \mathbb{R}^2) \simeq r$$

$$\Rightarrow \operatorname{dist}(R^{-1}(\gamma), \partial \Omega) \simeq r, \qquad R^{-1}(\gamma) \subset D(z, Cr)$$

$$\Rightarrow \rho_{\Omega}(z, w) \leq C$$



Moreover, $g = \iota \circ \sigma : \Omega \to \mathbb{D}$ is locally Euclidean Lipschitz.

$$|g'(z)| \simeq \frac{\operatorname{dist}(g(z), \partial \mathbb{D})}{\operatorname{dist}(z, \partial \Omega)}.$$

Use Fact 1

$$\operatorname{dist}(z, \partial\Omega) \simeq \operatorname{dist}(R(z), \mathbb{R}^2)$$

$$\simeq \exp(-\rho_{\mathbb{R}^3_+}(R(z), z_0))$$

$$\gtrsim \exp(-\rho_S(R(z), z_0))$$

$$= \exp(-\rho_D(g(z), 0))$$

$$\simeq \operatorname{dist}(g(z), \partial D)$$