

BRAID FORCING AND STAR-SHAPED TRAIN TRACKS

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ABSTRACT. Global results are proved about the way in which Boyland's forcing partial order organizes a set of braid types: those of periodic orbits of Smale's horseshoe map for which the associated train track is a *star*. This is a special case of a conjecture introduced in [9], which claims that forcing organizes all horseshoe braid types into linearly ordered families which are, in turn, parameterized by homoclinic orbits to the fixed point of code 0.

1. INTRODUCTION AND BACKGROUND

Given a discrete dynamical system and information about one of its periodic orbits, can one derive further information about the system: for example, the existence of other periodic orbits, or that it has positive topological entropy? The problem of periodic orbit forcing in particular — that is, to determine whether the presence of a certain periodic orbit implies the existence of other periodic orbits — has interested dynamicists for many years. Sharkovskii's theorem for self-maps of the real line is unrivalled in elegance: it defines a total order \preceq on the positive integers with the property that if $m \preceq n$ then any continuous self-map of \mathbb{R} which has a periodic orbit of period n must also have a periodic orbit of period m . Moreover, the theorem is sharp: for any initial segment of the order, there exists a continuous self-map of \mathbb{R} having periodic orbits of the corresponding periods and no others. In this context, periodic orbits are specified by their period alone. One could also, for example, use the order on \mathbb{R} to specify a periodic orbit by its permutation, and then consider forcing among permutations. Many authors have studied this problem, which is now quite well understood.

The corresponding problem in dimension 2 (i.e. for self-homeomorphisms of the disk D^2) is much harder. In this case, the period alone is an inadequate specification of a periodic orbit: given any set of positive integers which includes 1, a self-homeomorphism of D^2 (or indeed of any other surface) can be constructed which has periodic orbits of the given periods and no others. Boyland and others observed that, by analogy with the permutation in dimension 1, adding topological information about the way in which the points of a periodic orbit of a disk homeomorphism 'braid around' one another produces a non-trivial theory. More precisely, Boyland defined the *braid type* $\text{bt}(P, f)$ of a periodic orbit P of a disk homeomorphism f to be the isotopy class of f relative to P , up to topological change of coordinates. A braid type γ is then said to *force* another braid type β if every disk homeomorphism having a periodic orbit of type γ also has one of type β . Boyland showed that forcing is a partial order on the set of braid types. This leads to the question of how forcing organizes this set. Of course, if the question is to be answered, one must be able to decide how two given braid types are related by the partial order. This can be done using Bestvina and Handel's algorithmic proof of Thurston's classification theorem for surface homeomorphisms up to isotopy. However, not only is this process very time consuming in practice, but also the information it gives is

only local, in the sense that the ability to compare two given braid types gives no information about the global structure of the partially ordered set of braid types.

This paper gives global information about the restriction of the forcing order to a subset of braid types, namely those of periodic orbits of Smale's horseshoe map for which the associated train track is a *star* (that is, a tree with exactly one vertex of valence larger than 1). In [9], a conjecture is stated which describes how forcing organizes all braid types of the horseshoe. It claims that the symbolic code of each periodic orbit can be parsed into two segments, the *prefix* and the *decoration*. All orbits with the same decoration have the same *topological train track type*, and form a family which is totally ordered by the forcing relation; the position of an orbit within this totally ordered set is determined by its prefix. Within families, this trivializes the problem of comparing braid types: simply compare their symbolic codes using the unimodal order. The conjecture also describes the forcing relation between families in terms of the forcing between homoclinic orbits associated to each family.

In this paper the conjecture is proved for the (infinite) family of decorations for which the corresponding train track type is a star. There is one such decoration for each rational $m/n \in (0, 1/2]$: the train track associated to m/n is a star with n edges, and the train track map rotates the central vertex of the star by m/n .

The approach taken is to start with train track maps of the appropriate topological type, and to identify those of their periodic orbits for which the star is itself a train track (after it has been truncated outside the span of the orbit). The combinatorics of such orbits are intricately related to the position of the rational m/n within the Farey graph.

Only after these *train track orbits* have been determined is their relationship with periodic orbits of Smale's horseshoe investigated. Conceptually, the most important relationship between the horseshoe and the star maps (and indeed the relationship which motivates the conjecture under discussion) is that the star maps (or, more accurately, the corresponding thick tree maps) can be obtained from the horseshoe by *pruning*: that is, by performing an isotopy which destroys all of the dynamics within a given open subset of the disk, while leaving the dynamics unchanged elsewhere. Since this construction is not used elsewhere in the paper, however, it is only described on an intuitive level, and a less general approach to showing that star periodic orbits have horseshoe braid types is adopted.

Section 1.1 describes the parsing of the code of a horseshoe periodic orbit into prefix and decoration, and provides a statement of the conjecture of [9] and the special case of it which will be proved here. Section 1.2 contains a brief summary of the properties of the Farey graph which will be used in the remainder of the paper. The main results of the paper can be found in Section 2, in which star maps and the concept of a train track orbit are defined, and the set of such orbits is determined. Section 3 details the connection between the star maps and the horseshoe, and Section 4 is devoted to the proof of a technical lemma.

The paper is often technical: it is the nature of the subject. Its structure, however, is simple. Lemma 1.17 identifies the periodic orbits of the horseshoe which have the property that their period is equal to the sum of the denominators of the endpoints of their rotation interval. The symbolic codes of these orbits have the property that the prefix is determined by one endpoint of the rotation interval, and the decoration by the other. Having described the combinatorics of train track orbits of star maps (Section 2), it is possible to compute their rotation intervals, which have the property considered in Lemma 1.17. Since star orbits have horseshoe braid type (Section 3), it follows that the symbolic codes of the train track

orbits are as given by the lemma. The total order within families follows directly from the nature of the maps being considered.

1.1. Prefix and decoration for horseshoe periodic orbits. In this paper the standard model $F: D^2 \rightarrow D^2$ of Smale's horseshoe map [15] depicted in Fig. 1 will be used; symbolic dynamics in the set $\Sigma_2 = \{0, 1\}^{\mathbb{Z}}$ is applied in the usual way to describe points $x \in D^2$ whose (past and future) orbits lie entirely in the square S . The definitions and results summarized in this section can be found in [11, 9].

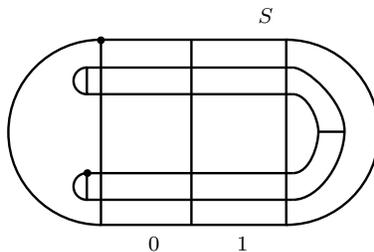


FIGURE 1. Symbolic dynamics for the horseshoe

A periodic orbit P of F of (least) period n is described by its *code* $c_P \in \{0, 1\}^n$, which is given by the first n symbols of the itinerary $k(p)$ of its rightmost point p : thus, for example, the period 5 orbit which contains the point with itinerary $\overline{01001}$ has code 10010. A word $w \in \{0, 1\}^n$ is therefore the code of a period n horseshoe orbit if and only if it is *maximal*: that is, the infinite repetition $\overline{w} \in \{0, 1\}^{\mathbb{N}}$ of w is strictly greater than its shifts $\sigma^i(\overline{w})$ in the unimodal order for $1 \leq i < n$.

The following definitions are due to Boyland [5, 6]. Let $f: D^2 \rightarrow D^2$ and $g: D^2 \rightarrow D^2$ be orientation-preserving homeomorphisms having periodic orbits P and Q respectively. Then (P, f) and (Q, g) *have the same braid type* if there is a homeomorphism $h: (D^2, P) \rightarrow (D^2, Q)$ such that $f: (D^2, P) \rightarrow (D^2, P)$ is isotopic (rel. P) to $h^{-1} \circ g \circ h: (D^2, P) \rightarrow (D^2, P)$ (if either P or Q lies on ∂D^2 , then the corresponding homeomorphism should first be extended arbitrarily over an exterior collar). The *braid type* $\text{bt}(P, f)$ of (P, f) is its equivalence class under this relation. Since the braid types of period n orbits correspond to conjugacy classes in the mapping class group of the n -punctured disk, they can be classified as finite order, reducible, or pseudo-Anosov by means of Thurston's classification [16]. Boyland's *forcing relation* \leq on the set BT of braid types is defined as follows: $\beta \leq \gamma$ if and only if every orientation-preserving homeomorphism $f: D^2 \rightarrow D^2$ which has a periodic orbit of braid type γ also has one of braid type β . Boyland proved [5, 6] that \leq is a partial order on BT.

An alternative characterisation of braid type indicates the dynamical significance of the definition: (P, f) and (Q, g) have the same braid type if and only if there exists an isotopy $\{f_t\}$ from f to g and a path P_t in $(D^2)^n$ from P to Q such that P_t is a (least) period n orbit of f_t for all t .

The braid type $\text{bt}(P, F)$ of a horseshoe periodic orbit P will here be denoted simply by $\text{bt}(P)$. It is well known that two periodic orbits P and Q of the horseshoe whose codes c_P and c_Q differ only in their final symbol have the same braid type. Thus, for example, the two orbits with codes 10010 and 10011 have the same braid type; the code of either one of

these orbits is often written $c_P = 1001_1^0$ to reflect the fact that the distinction between the two is unimportant in so far as braid type is concerned.

The conjecture presented in [9] is based upon the parsing of the code of any horseshoe periodic orbit P which is not of finite order braid type into two parts: the *prefix* and the *decoration*. In order to define these, it is necessary first to introduce a word $c_q \in \{0, 1\}^{n+1}$ for each rational $q = m/n \in (0, 1/2]$, and then to use these to define the *height* $q(P) \in \mathbb{Q} \cap (0, 1/2]$ of P : this is an invariant of braid type which plays a central role in the conjecture. Motivation for the definition is given in [11], and a program for computing heights of horseshoe periodic orbits can be found at [12].

The definition of the words c_q is first given in an easily accessible form: a more practical way of computing them is described afterwards.

Definition 1.1. Given $q = m/n \in \mathbb{Q} \cap (0, 1/2]$, define a word $c_q \in \{0, 1\}^{n+1}$ as follows. Let L_q be the straight line in \mathbb{R}^2 from $(0, 0)$ to (n, m) . For $0 \leq i \leq n$, let $s_i = 1$ if L_q crosses some line $y = \text{integer}$ for $x \in (i - 1, i + 1)$, and $s_i = 0$ otherwise. Then $c_q = s_0 s_1 \dots s_n$.

Example 1.2. Figure 2 shows that $c_{3/10} = 10011011001$.

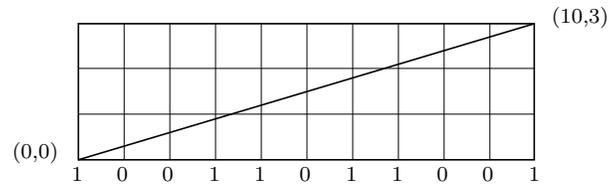


FIGURE 2. $c_{3/10} = 10011011001$

The words c_q are manifestly palindromic. Their general form is indicated by the examples in Figure 3, in which the column headings and row headings denote the numerator and denominator of q respectively. The $n - 2m + 1$ zeros are partitioned ‘as even-handedly as possible’ into m subwords (possibly empty), separated by 11.

	1	2	3	4
3	1001			
4	10001			
5	100001	101101		
6	1000001			
7	10000001	10011001	10111101	
8	100000001		101101101	
9	1000000001	1000110001		1011111101
10	10000000001		10011011001	
11	100000000001	100001100001	100110011001	101101101101

FIGURE 3. Examples of the words c_q

The following more practical method for determining c_q is easily shown to be equivalent: given $q = m/n$, define integers $\kappa_i(q)$ for $1 \leq i \leq m$ by

$$(1) \quad \kappa_i(q) = \begin{cases} \lfloor n/m \rfloor - 1 & \text{if } i = 1 \\ \lfloor in/m \rfloor - \lfloor (i-1)n/m \rfloor - 2 & \text{if } 2 \leq i \leq m \end{cases}$$

(here $\lfloor x \rfloor$ denotes the greatest integer which does not exceed x). Then $c_q = 10^{\kappa_1(q)} 110^{\kappa_2(q)} 11 \dots 110^{\kappa_m(q)} 1$.

Example 1.3. Let $q = 3/7$, so $\kappa_1(q) = \lfloor 7/3 \rfloor - 1 = 1$, $\kappa_2(q) = \lfloor 14/3 \rfloor - \lfloor 7/3 \rfloor - 2 = 4 - 2 - 2 = 0$, and $\kappa_3(q) = \lfloor 21/3 \rfloor - \lfloor 14/3 \rfloor - 2 = 7 - 4 - 2 = 1$. Thus $c_{3/7} = 10111101$.

Height is an invariant of horseshoe braid type taking rational values in $(0, 1/2]$. The next lemma [11] is necessary for the definition, or at least for the definition to be sensible: it implies that the height function $q: \{0, 1\}^{\mathbb{N}} \rightarrow [0, 1/2]$ is decreasing with respect to the unimodal order \prec on $\{0, 1\}^{\mathbb{N}}$.

Lemma 1.4. For each $q \in \mathbb{Q} \cap (0, 1/2]$, the two words c_{q1}^0 are both maximal. Moreover, if $q, r \in \mathbb{Q} \cap (0, 1/2]$ with $q < r$ then $c_{r1}^0 \prec c_{q1}^0$.

Definitions 1.5. Let $c \in \{0, 1\}^{\mathbb{N}}$. Then the *height* $q(c) \in [0, 1/2]$ of c is given by

$$q(c) = \inf\{q \in \mathbb{Q} \cap (0, 1/2] : q = 1/2 \text{ or } \overline{c_q 0} \prec c\}.$$

The *height* $q(P)$ of a horseshoe periodic orbit P of code c_P is given by $q(P) = q(\overline{c_P})$.

The algorithm described in the following lemma [11] can be used to compute $q(c)$ for any $c \in \{0, 1\}^{\mathbb{N}}$ which contains the word 010. In particular (see [11]), it can be used to compute $q(P)$ for any periodic orbit P : if $\overline{c_P}$ doesn't contain the word 010, then change the final symbol of c_P from 1 to 0 before applying the algorithm. In particular, the height of a horseshoe periodic orbit is rational and strictly positive.

Lemma 1.6. Let $c \in \{0, 1\}^{\mathbb{N}}$, and suppose that c contains the word 010. Then $q(c)$ can be calculated as follows. If c does not begin 10, then $q(c) = 1/2$. Otherwise, write

$$c = 10^{\kappa_1} 1^{\mu_1} 0^{\kappa_2} 1^{\mu_2} \dots,$$

where each $\kappa_i \geq 0$, each μ_i is either 1 or 2, and $\mu_i = 1$ only if $\kappa_{i+1} > 0$ (thus κ_i and μ_i are uniquely determined by c). For each $r \geq 1$, define

$$I_r(c) = \left(\frac{r}{2r + \sum_{i=1}^r \kappa_i}, \frac{r}{(2r-1) + \sum_{i=1}^r \kappa_i} \right],$$

and let $s \geq 1$ be the least integer such that either $\mu_s = 1$ or $\bigcap_{i=1}^{s+1} I_i(c) = \emptyset$. Write $\bigcap_{i=1}^s I_i(c) = (x, y]$. Then

$$q(c) = \begin{cases} x & \text{if } \mu_s = 2 \text{ and } w < z \text{ for all } w \in I_{s+1}(c) \text{ and } z \in \bigcap_{i=1}^s I_i(c) \\ y & \text{if } \mu_s = 1, \text{ or } \mu_s = 2 \text{ and } w > z \text{ for all } w \in I_{s+1}(c) \text{ and } z \in \bigcap_{i=1}^s I_i(c). \end{cases}$$

Notice that some μ_i is equal to 1 (since c contains the word 010), and hence the algorithm terminates.

Example 1.7. Let P be the horseshoe periodic orbit of code 10111101111. Then $q(P) = q(\overline{10111101110})$. Using the notation of the definition, $\kappa_1 = 1$, $\mu_1 = 2$, $\kappa_2 = 0$, $\mu_2 = 2$, $\kappa_3 = 1$, $\mu_3 = 2$, $\kappa_4 = 0$, and $\mu_4 = 1$. Hence $I_1 = (1/3, 1/2]$, $I_2 = (2/5, 2/4]$ (so $\bigcap_{i=1}^2 I_i = (2/5, 1/2]$),

$I_3 = (3/8, 3/7]$ (so $\bigcap_{i=1}^3 I_i = (2/5, 3/7]$), and $I_4 = (4/10, 5/9]$ (so $\bigcap_{i=1}^4 I_i = (2/5, 3/7]$). Since $\mu_4 = 1$, it follows that $q(P) = 3/7$. Although tedious to carry out by hand, the height algorithm is easily automated [12].

The next theorem [11] summarizes the properties of height which will be used here.

Theorem 1.8. *a) Let P be a horseshoe periodic orbit with height $q(P) = m/n$ in lowest terms. Then P has period n if and only if it has finite order braid type: in this case, F is isotopic rel. P to a rigid rotation through $2\pi m/n$. Otherwise, the period of P is at least $n + 2$, and c_P starts with the word $c_{m/n}$.*

b) Let P and R be horseshoe periodic orbits. If $\text{bt}(P) \geq \text{bt}(R)$ then $q(P) \leq q(R)$. In particular, height is an invariant of braid type.

Thus the code c_P of a horseshoe periodic orbit which is not of finite order braid type can be written $c_P = c_{q(P)}v$ for some word v of length at least 1. This makes possible the following definitions, which are taken from [9]:

Definitions 1.9. Let P be a period N orbit of the horseshoe which is not of finite order braid type, with height $q = q(P) = m/n$. The *prefix* of P is the word c_q . The *decoration* of P is defined to be $*$ if $N = n + 2$, and to be the element w of $\{0, 1\}^{N-n-3}$ such that

$$c_P = c_{q_1^0} w_1^0$$

otherwise.

Example 1.10. Let P be the period 17 orbit with code $c_P = 10011011001011010$. Then $q(P) = 3/10$ (using the algorithm of Lemma 1.6). Hence P has prefix $10011011001 = c_{3/10}$, and decoration 1101.

Only certain heights q are compatible with a given decoration w (namely those for which $c_{q_1^0} w_1^0$ is a maximal word, and hence describes the rightmost point of a periodic orbit). The following lemma (from [9]) gives the compatibility conditions:

Lemma 1.11. *Let w be a decoration, and define $q_w \in \mathbb{Q} \cap (0, 1/2]$ by $q_* = 1/2$ and*

$$q_w = \min_{0 \leq i \leq k+2} q(\sigma^i(\overline{10w0}))$$

*if $w \in \{0, 1\}^k$. Then each of the four words $c_{q_1^0} w_1^0$ (or each of the two words $c_{q_1^0}$ when $w = *$) is maximal of height q when $0 < q < q_w$, and none is maximal of height q when $q_w < q \leq 1/2$.*

The reason that the four orbits with codes $c_{q_1^0} w_1^0$ are considered together is that they all have the same braid type: the following general result will appear in [10]. That the braid type is unchanged on changing the final symbol of the code is a triviality: the content of the theorem is that changing the symbol between prefix and decoration also leaves the braid type unchanged.

Theorem 1.12. *Let $w \neq *$ be a decoration, and $q \in \mathbb{Q} \cap (0, q_w)$. Then the four horseshoe periodic orbits with codes $c_{q_1^0} w_1^0$ have the same braid type.*

The following conjecture is motivated and stated in [9]. It involves the notion of two horseshoe periodic orbits of pseudo-Anosov braid type having the *same topological train track type*. A formal definition of this term can be found in [9]: however its intuitive meaning should be clear in the context of this paper (in which ‘star’ train track types are considered).

- Conjecture 1.13.** a) *Let P be a horseshoe periodic orbit with height q and decoration w . If $q \neq q_w$ then P has pseudo-Anosov braid type.*
- b) *Let P' be another periodic orbit with height q' and decoration w . Then $\text{bt}(P') \leq \text{bt}(P)$ if and only if $q \leq q'$ (i.e. the set of braid types of periodic orbits with decoration w is totally ordered by the forcing relation). Moreover, if $q' < q_w$ then P and P' have the same topological train track type.*
- c) *There is an equivalence relation \sim on the set \mathcal{D} of decorations with the property that two horseshoe periodic orbits have the same braid type if and only if they have equal heights and equivalent decorations.*
- d) *There is a partial order \preceq on \mathcal{D}/\sim with the property that if P and P' are periodic orbits with heights q, q' and decorations w, w' such that $q < q'$ and $[w'] \preceq [w]$, then $\text{bt}(P') \leq \text{bt}(P)$.*

Notice that part b) trivializes the problem of comparing braid types within a family of fixed decoration: the forcing order is given by the unimodal order on the codes of the periodic orbits. As explained in [9], parts a) and b) of this conjecture (i.e., the statements which only concern a single decoration) can be proved for certain particular choices of decoration w . In this paper an infinite family of decorations is considered, making it possible to address parts c) and d) of the conjecture meaningfully. The decorations concerned are those which give rise to ‘star’ topological train track types, and are given by words w_q with $q \in \mathbb{Q} \cap (0, 1/2]$ defined as follows:

Definition 1.14. Let $q = m/n \in \mathbb{Q} \cap (0, 1/2]$. If $q \neq 1/2$ then w_q is the element of $\{0, 1\}^{n-3}$ obtained by deleting the initial symbols 10 and the final symbols 01 from c_q . If $q = 1/2$ then $w_q = *$.

Thus, for example, $w_{1/3}$ is the empty decoration, $w_{1/4} = 0$, $w_{1/5} = 00$, and $w_{2/5} = 11$. These decorations satisfy $q_{w_q} = q$, are mutually non-equivalent under \sim , and are totally ordered by \preceq , with $w_q \preceq w_{q'}$ if and only if $q \leq q'$. Thus the proof of Conjecture 1.13 above for these decorations gives the following theorem, which summarizes some of the main results of this paper.

Theorem 1.15. *Let $r, r' \in \mathbb{Q} \cap (0, 1/2]$ and $q, q' \in \mathbb{Q} \cap (0, 1/2)$ with $q \leq r$ and $q' \leq r'$. Let P be a horseshoe periodic orbit of height q and decoration w_r , and P' be a horseshoe periodic orbit of height q' and decoration $w_{r'}$. Then*

- a) *If $q \neq r$ then P has pseudo-Anosov braid type.*
- b) *If $r = r'$ and $q, q' < r$ then P and P' have the same (star) topological train track type. Moreover $\text{bt}(P') \leq \text{bt}(P)$ if and only if $q \leq q'$.*
- c) *P and P' have the same braid type if and only if $q = q'$ and $r = r'$.*
- d) *If $q < q'$ and $r \geq r'$ then $\text{bt}(P') \leq \text{bt}(P)$.*

Part a) of this theorem is given by Corollary 1.18 below, Part b) by Theorem 2.34, and Part c) by Theorem 1.12 and Lemma 1.17, as described following the statement of Lemma 1.17. Part d) is proved using quite different techniques from those of this paper: a proof will appear in [10].

Lemma 1.17 below will be used to identify periodic orbits with decoration $w_{m/n}$ for some $m/n \in (0, 1/2]$. Its proof is relatively long and technical, and uses methods from [11] which are not required elsewhere in this paper — it has therefore been relegated to Section 4.

Definitions 1.16. Let P be a horseshoe periodic orbit of period $N > 1$. The *rotation number* $\rho(P) \in \mathbb{Q} \cap (0, 1/2]$ of P is its F -rotation number about the fixed point with code 1. The *rotation interval* of P is the set

$$\rho i(P) = \{\rho(P') : \text{bt}(P') \leq \text{bt}(P)\}.$$

The rotation number of a horseshoe periodic orbit is an invariant of braid type [11] (this statement is not quite as obvious as it may at first appear), and hence so also is the rotation interval. By a theorem of Handel [13], $\rho i(P)$ is equal either to $\{\rho(P)\}$ or to a set of the form $\mathbb{Q} \cap [a, b]$ (which will here be denoted simply $[a, b]$) for some $a < b \in \mathbb{Q}$.

Lemma 1.17. *Let P be a period $N > 1$ orbit of the horseshoe with non-trivial rotation interval $\rho i(P) = [u/v, m/n]$. Then $N \geq v + n$. Moreover, $N = v + n$ if and only if c_P is one of the four words $c_{u/v_1^0} w_{m/n_1^0}$ (or one of the two words c_{u/v_1^0} in the case $m/n = 1/2$).*

In other words, if the period of P is equal to the sum of the denominators of the endpoints of its rotation interval, then P has height u/v and decoration $w_{m/n}$. Since the rotation interval is a braid type invariant, it follows that if P has height q and decoration w_r , and P' is a horseshoe periodic orbit with the same braid type as P , then P' also has height q and decoration w_r . In particular, this establishes (and extends) the ‘only if’ part of Theorem 1.15 c). The ‘if’ part follows from Theorem 1.12. Moreover, the following corollary provides a proof of Theorem 1.15 a).

Corollary 1.18. *Let $u/v < m/n \leq 1/2$. Then the horseshoe periodic orbits of height u/v and decoration $w_{m/n}$ have pseudo-Anosov braid type.*

Proof. Let P be a periodic orbit of height u/v and decoration $w_{m/n}$. Then $\rho i(P) = [u/v, m/n]$. Periodic orbits of finite order type have trivial rotation intervals. If P had reducible braid type, then a horseshoe periodic orbit of period less than $v+n$ (namely one given by the isotopy class corresponding to the outermost reducible component of $F: (D^2, P) \rightarrow (D^2, P)$) would also have rotation interval $[u/v, m/n]$, contradicting Lemma 1.17. \square

Lemma 1.17 may be of independent interest. The authors know of no pseudo-Anosov homeomorphism f of the annulus, relative to a single periodic orbit P , for which the sum of the denominators of the endpoints of the rotation interval of f is less than the period n of P . If this more general result could be proved, then (together with the fact [3] that the rotation number m/n of P is contained in the interior of the rotation interval of f) it would provide a natural generalization of Boyland’s theorem [4], which gives bounds on the size of the rotation interval in the case where m and n are coprime.

1.2. The Farey Graph. In this section some well known results about Farey sequences and the Farey graph are presented, some notation is introduced, and some simple number-theoretic lemmas which will be used later are stated. In the usual treatment (see for example [14], where the assertions made in this section are proved), the Farey sequences contain rationals between 0 and 1: since only rationals between 0 and $1/2$ are of interest in this paper, the definitions have been modified accordingly. Throughout the paper, all rationals are assumed to be written in lowest terms.

Definitions 1.19. Let $n \geq 1$ be an integer. The *Farey sequence* \mathcal{F}_n of order n is the (finite) sequence of rational numbers in $[0, 1/2]$ whose denominators do not exceed n , arranged in ascending order. Two rationals h/k and h'/k' in $[0, 1/2]$ are *Farey neighbours* if there is some

n such that they are consecutive elements of \mathcal{F}_n . Given a rational $h/k \in (0, 1/2)$, the *left and right Farey parents* $\text{LFP}(h/k)$ and $\text{RFP}(h/k)$ are the elements of \mathcal{F}_k which precede and follow h/k ; and h/k is said to be the *Farey child* of its parents.

In particular, both $\text{LFP}(h/k)$ and $\text{RFP}(h/k)$ are Farey neighbours of h/k . It is well known that if $h/k < h'/k'$ are Farey neighbours, then $h'k - hk' = 1$; and that if $\text{LFP}(h/k) = u/v$ and $\text{RFP}(h/k) = p/q$, then $h = u + p$ and $k = v + q$.

The rationals in $[0, 1/2]$ can be organized as the vertices of the *Farey Graph*, in which an edge joins two vertices if and only if one of the associated rationals is the left or right Farey parent of the other. Part of the Farey graph is depicted in Fig. 4. Notice that every vertex has infinite valence (that is, every rational is the parent of infinitely many other rationals). The *immediate left (respectively right) Farey child* of a rational m/n is the rational h/k of smallest denominator for which $\text{LFP}(h/k) = m/n$ (respectively $\text{RFP}(h/k) = m/n$): it is the child of m/n and $\text{LFP}(m/n)$ (respectively $\text{RFP}(m/n)$). Thus, for example, the immediate children of $2/5$ are $3/8$ and $3/7$. The Farey sequences appear as finite connected subtrees of the Farey graph: for example, \mathcal{F}_5 is shown in the figure with bolder lines.

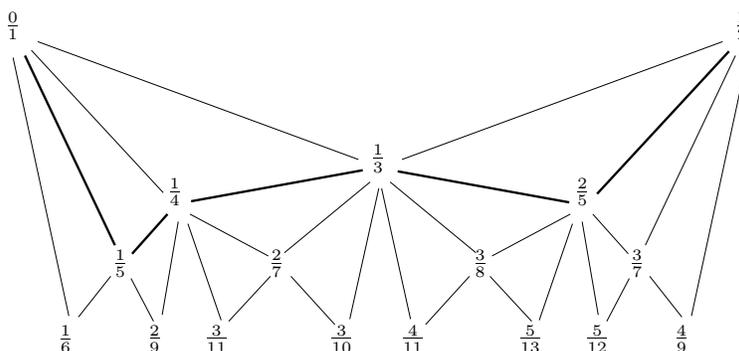


FIGURE 4. Part of the Farey Graph

Definition 1.20. Let m/n be a rational in $(0, 1/2)$. Then the *left Farey sequence* $\text{LFS}(m/n)$ of m/n is the (finite) sequence $(0 = u_1/v_1, u_2/v_2, \dots, u_\alpha/v_\alpha)$, where $u_\alpha/v_\alpha = \text{LFP}(m/n)$, and $u_i/v_i = \text{LFP}(u_{i+1}/v_{i+1})$ for $1 \leq i < \alpha$.

Thus, for example, $\text{LFS}(3/10) = (0, 1/4, 2/7)$.

Definition 1.21. Let m/n be a rational in $(0, 1/2)$, and denote addition modulo n by $+_n$. Given $r, s \in \mathbb{Z}_n$, let

$$\mathcal{O}_{m/n}[r, s] = \{r +_n jm : 0 \leq j \leq k, \text{ where } k \geq 0 \text{ is least such that } r +_n km = s\}.$$

That is, $\mathcal{O}_{m/n}[r, s]$ is the shortest segment of the orbit of r in \mathbb{Z}_n under addition of m which ends with s . When $r \leq s$ are integers, the notation $[r, s]$ will also be used as a shorthand for $\{r, r + 1, \dots, s\}$.

The following simple Lemma contains the results about the Farey graph which will be needed in the remainder of the paper.

Lemma 1.22. Let $m/n \in (0, 1/2)$, and suppose that $\text{LFP}(m/n) = u/v$ and $\text{RFP}(m/n) = p/q$. Then

- a) $\#\mathcal{O}_{m/n}[m, n-1] = q = n - v$, and $\#\mathcal{O}_{m/n}[m-1, 0] = n - q = v$.
b) $\#(\mathcal{O}_{m/n}[m, n-1] \cap [1, m]) = \#(\mathcal{O}_{m/n}[m, n-1] \cap [n-m, n-1]) = p = m - u$.
c) $\#(\mathcal{O}_{m/n}[m-1, 0] \cap [1, m]) = \#(\mathcal{O}_{m/n}[m-1, 0] \cap [n-m, n-1]) = m - p = u$.

Proof.

- a) By definition of $R = \mathcal{O}_{m/n}[m, n-1]$ its cardinality is $k+1$, where $k \in \mathbb{Z}_n$ is such that $m + km \equiv n-1 \pmod{n}$: that is, $(k+1)(m-n) \equiv -1 \pmod{n}$, so $k+1$ is the multiplicative inverse $(n-m)_n^{-1}$ of $n-m$ in \mathbb{Z}_n .

Since $p/q = \text{RFP}(m/n)$, it follows that $q \in [1, n-1]$, and that $qm = np - 1$. Thus $q(n-m) \equiv -qm \equiv 1 \pmod{n}$, so $(n-m)_n^{-1} = q$ as required. However $v + q = n$, so $\#R = n - v$ also. Since R and $S = \mathcal{O}_{m/n}[m-1, 0]$ partition \mathbb{Z}_n , the cardinality of S follows.

- b) Let $\#(R \cap [1, m]) = s$. Since R has cardinality q , it follows that $m + (q-1)m = n-1 + (s-1)n$, or $qm + 1 = sn$. Thus $s = p$. However $u + p = m$, so $s = m - u$ also. A similar argument applies to $R \cap [n-m, n-1]$.

- c) Immediate from part b). □

The function defined in the following lemma will play a central role in this paper.

Lemma 1.23. *Let $u/v < p/q$ be Farey neighbours. Then the function*

$$\xi_{u/v, p/q}: (0, 1) \cap \mathbb{Q} \rightarrow (u/v, p/q) \cap \mathbb{Q}$$

defined by

$$\xi_{u/v, p/q} \left(\frac{r}{s} \right) = \frac{rp + (s-r)u}{rq + (s-r)v}$$

is an increasing bijection.

Proof. Since $u/v < p/q$ are Farey neighbours, $pv - qu = 1$. Hence if $m/n \in (u/v, p/q)$, then $m/n = \xi_{u/v, p/q}((vm - un)/((v - q)m + (p - u)n))$ (note that $0 < vm - un < (v - q)m + (p - u)n = (vm - un) + (pn - qm)$ (since $u/v < m/n < p/q$), and that $vm - un$ and $(v - q)m + (p - u)n$ are coprime (since $m = (vm - un)(p - u) + ((v - q)m + (p - u)n)u$ and $n = (vm - un)(q - v) + ((v - q)m + (p - u)n)v$ are coprime)). Thus $\xi_{u/v, p/q}$ is surjective. It is also strictly increasing, since if $0 < r/s < a/b < 1$ then $\xi_{u/v, p/q}(a/b) - \xi_{u/v, p/q}(r/s) = \frac{as - br}{(aq + (b-a)v)(rq + (s-r)v)} > 0$. □

Notice in particular that $\xi_{u/v, p/q}(1/2) = (u + p)/(v + q)$ is the Farey child of u/v and p/q .

2. PERIODIC ORBITS OF STAR MAPS

2.1. Star maps and *-orbits. For each $n \geq 2$ let $\Gamma_n \subseteq D^2$ be an n -star: that is, a tree with n edges e_0, \dots, e_{n-1} of equal length, each of which has a valence 1 vertex at its initial point, and which meet at a valence n vertex v at their final points, cyclically ordered according to their indices. For each rational $m/n \in (0, 1/2]$ (written in its lowest terms), let $f_{m/n}: \Gamma_n \rightarrow \Gamma_n$ be the tree map (see Fig. 5) with image edge paths

$$\begin{aligned} f_{m/n}(e_0) &= e_0 \bar{e}_1 e_1 \bar{e}_2 e_2 \dots \bar{e}_m e_m \\ f_{m/n}(e_r) &= e_{r+n} \quad (r \neq 0), \end{aligned}$$

which expands each edge uniformly away from the preimages of vertices (each connected component of which is either a nontrivial interval or, in the case of the preimage component

containing v , is homeomorphic to Γ_n). Figures depicting $f_{m/n} : \Gamma_n \rightarrow \Gamma_n$ (with the exception of Fig. 11) are always drawn with e_0 and e_{n-m} horizontal.

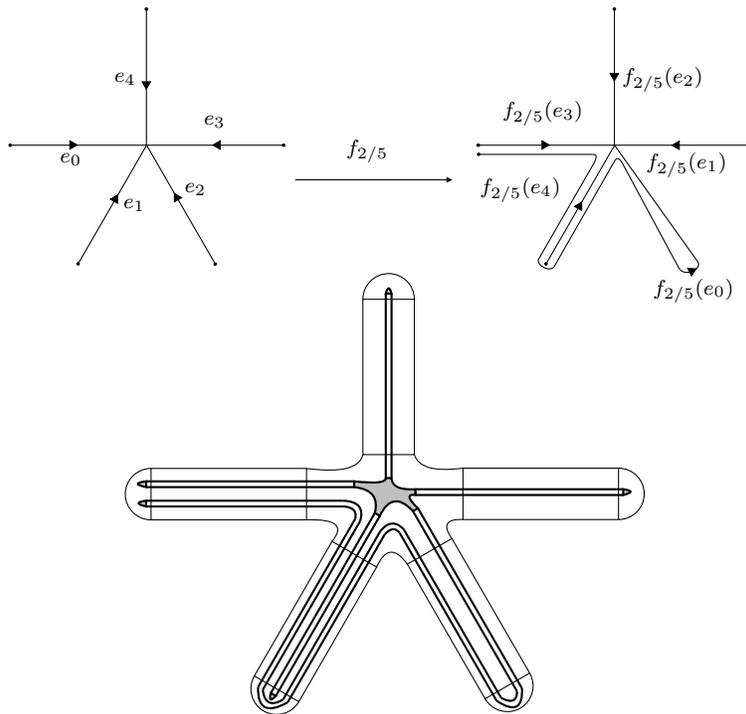


FIGURE 5. The tree map $f_{2/5} : \Gamma_5 \rightarrow \Gamma_5$ and the thick tree map $F_{2/5} : T_5 \rightarrow T_5$

Let $F_{m/n} : T_n \rightarrow T_n$ be a *thick tree map* (see for example [8] for a formal definition) corresponding to $f_{m/n} : \Gamma_n \rightarrow \Gamma_n$ (see Fig. 5). Thus $T_n \subseteq D^2$ is a topological disk, and there is a continuous map $p : T_n \rightarrow \Gamma_n$ such that $p^{-1}(x)$ is a disk if x is a vertex of Γ_n , and an interval otherwise. The map $F_{m/n}$ is an embedding which contracts each such *decomposition element* $p^{-1}(x)$ into a decomposition element, in such a way that it induces $f_{m/n}$ on Γ_n (so $F_{m/n}(p^{-1}(x)) \subset p^{-1}(f_{m/n}(x))$ for all $x \in \Gamma_n$). Where necessary, $F_{m/n}$ is considered as a homeomorphism $D^2 \rightarrow D^2$, by extending from T_n without introducing any new periodic orbits.

There is a natural one-to-one correspondence (which will be invoked without comment in the remainder of the paper) between the periodic orbits of $f_{m/n}$ and those of $F_{m/n}$ (with a periodic point x of $f_{m/n}$ corresponding to the unique periodic point of $F_{m/n}$ in $p^{-1}(x)$). Notice that $F_{1/2}$ is the horseshoe map F after blowing up the leaf containing the fixed point of code 1 into a disk: in particular, there is a braid type preserving bijection between the set of periodic orbits of $F_{1/2}$ and the set of periodic orbits of F .

In this paper only periodic orbits P which satisfy certain non-triviality conditions are considered:

Definition 2.1. A periodic orbit P of $f_{m/n} : \Gamma_n \rightarrow \Gamma_n$ is called a **-orbit* if

- * a) $P \neq \{v\}$.
- * b) $P \cap e_r \neq \emptyset$ for all r .
- * c) $f_{m/n}(P \cap e_0) \not\subseteq e_m$.

*d) If p_r denotes the point of $P \cap e_r$ closest to the initial point of e_r , then $f_{m/n}(p_r) = p_{r+nm}$ for all $r \neq 0$.

The reasons for imposing three of these conditions are intuitively clear: a) states that P is not a fixed point; b) that it explores each of the edges of Γ_n ; and c) that the points of P in e_0 have images in more than one edge: if this were not true, the braid type of P would either be finite order (if each edge contained just one point of P) or reducible (with n reducing curves, each bounding a disk containing the points of P on one of the edges of Γ_n). Condition d) is less clear: the motivation is that when $f_{m/n}$ is ‘truncated’ with respect to P (see Definition 2.5), there is only one point of P (namely the preimage of p_m) at which $f_{m/n}$ is not locally injective. Lemma 2.3 below is one important consequence of this condition.

Definition 2.2. Let P be a periodic orbit of $f_{m/n}$. Then the *span* Γ_n^P of P is the smallest connected subset of Γ_n containing P .

Lemma 2.3. Let $m/n \in (0, 1/2]$. Then the set of spans of $*$ -orbits of $f_{m/n}$ is totally ordered by inclusion.

Proof. Let P and P' be $*$ -orbits of $f_{m/n}$, and for each $0 \leq r < n$ let $p_r, p'_r \in e_r$ be the points of the orbits closest to the initial point of e_r (which exist by *b)). If p_m is closer to the initial point of e_m than p'_m , then applying *d) inductively gives that p_r is closer than p'_r to the initial point of e_r for all r , and hence $\Gamma_n^{P'} \subseteq \Gamma_n^P$. \square

The following result, which is contained in Theorem 3.3 below, makes it possible to identify the braid types of $*$ -orbits using Lemma 1.17. It reflects the fact that the thick tree maps $F_{m/n}: T_n \rightarrow T_n$ can be obtained from the horseshoe by pruning (i.e. by destroying some dynamics), as described in Section 3.1.

Lemma 2.4. Every $*$ -orbit of $f_{m/n}$ has the braid type of some periodic orbit of the horseshoe.

2.2. Describing $*$ -orbits. In this section a combinatorial method for describing $*$ -orbits is developed. Since $f_{m/n}(e_r) = e_{r+nm}$ for $r \neq 0$, the main work required is in describing the images of the points of the orbit in e_0 . The first step is to make precise the notion of truncating $f_{m/n}$ with respect to a $*$ -orbit P .

Definition 2.5. Let P be a $*$ -orbit of $f_{m/n}$. Let $r_P: \Gamma_n \rightarrow \Gamma_n^P$ be the map defined by $r_P(x) = x$ if $x \in \Gamma_n^P$, and $r_P(x)$ is the endpoint of Γ_n^P contained on the same edge of Γ_n as x if $x \notin \Gamma_n^P$. The *truncation* $f_{m/n}^P$ of $f_{m/n}$ with respect to P is the map $f_{m/n}^P = r_P \circ f_{m/n}: \Gamma_n^P \rightarrow \Gamma_n^P$.

The following definitions give a basic classification of $*$ -orbits.

Definitions 2.6. Let P be a $*$ -orbit of $f_{m/n}$ with $\#(P \cap e_0) = N$, and label the points of $P \cap e_0$ as p_0, p_1, \dots, p_{N-1} in order from the initial to the final point of e_0 . Define a partition

$$\{0, \dots, N-1\} = A \cup B \cup C$$

by

$$i \in \begin{cases} A & \text{if } f_{m/n}^P \text{ is locally orientation-reversing at } p_i \\ B & \text{if } f_{m/n}^P \text{ is locally orientation-preserving at } p_i \\ C & \text{if } f_{m/n}^P \text{ is not locally injective at } p_i. \end{cases}$$

Observe that $i \in C$ if and only if $f_{m/n}(p_i)$ is the point of P closest to the initial point of e_m , and hence $\#C = 1$.

Given a rational $m/n \in (0, 1/2]$, an integer $k \in \{0, \dots, m-1\}$, and $\gamma \in \{A, B\}$ (with $\gamma = B$ if $k = 0$), write $\mathcal{P}(m/n, k, \gamma)$ for the set of all $*$ -orbits of $f_{m/n}$ with $f(p_0) \in e_k$ and $0 \in \gamma$, $\mathcal{P}(m/n, k) = \mathcal{P}(m/n, k, A) \cup \mathcal{P}(m/n, k, B)$, and $\mathcal{P}(m/n) = \bigcup_k \mathcal{P}(m/n, k)$, the set of all $*$ -orbits of $f_{m/n}$.

The rest of the description of a $*$ -orbit is contained in the next definitions.

Definitions 2.7. For each integer $n \geq 2$, define \mathcal{D}_n to be the set of all triples

$$d = ((N_0, \dots, N_{n-1}), \pi, (A, B, C)),$$

where N_r is a positive integer for $0 \leq r < n$; π is a cyclic permutation of the set

$$\mathcal{L} = \mathcal{L}_d = \{(r, s) : 0 \leq r < n, 0 \leq s < N_r\};$$

and (A, B, C) is a partition of $\{0, \dots, N_0 - 1\}$ with $\#C = 1$.

Let $P \in \mathcal{P}(m/n, k, \gamma)$. Then the *data* of P is the element

$$d(P) = ((N_0, \dots, N_{n-1}), \pi, (A, B, C))$$

of \mathcal{D}_n obtained as follows: $N_r = \#(P \cap e_r)$ for each r with $0 \leq r < n$. Identify P with $\mathcal{L} = \mathcal{L}_{d(P)} = \{(r, s) : 0 \leq r < n, 0 \leq s < N_r\}$ by labelling the points of $P \cap e_r$ as $(r, 0), (r, 1), \dots, (r, N_r - 1)$ from the initial to the final point of e_r , and let $\pi = f|_P: \mathcal{L} \rightarrow \mathcal{L}$. Finally, let (A, B, C) be the partition of $\{0, \dots, N - 1\}$ given by Definitions 2.6.

Having the same data is the basic equivalence relation which says that two $*$ -orbits have the same ‘shape’. In particular, it is clear that two $*$ -orbits with the same data have the same braid type as periodic orbits of $F_{m/n}: T_n \rightarrow T_n$. In this paper no distinction is made between orbits with the same data: thus, for example, the statement that $f_{m/n}$ has exactly one $*$ -orbit with a given property should be interpreted as meaning that $*$ -orbits with the given property exist, and all have the same data.

Example 2.8. Let $P \in \mathcal{P}(2/5, 1, A)$ be the periodic orbit depicted in Fig. 6. Then $N_0 = N_1 = N_3 = 2$ and $N_2 = N_4 = 1$; and $A = \{0\}$, $B = \emptyset$, and $C = \{1\}$. The cyclic permutation π is given by

$$\overset{A}{(0, 0)} \rightarrow (1, 1) \rightarrow (3, 1) \xrightarrow{C} (0, 1) \rightarrow (2, 0) \rightarrow (4, 0) \rightarrow (1, 0) \rightarrow (3, 0).$$

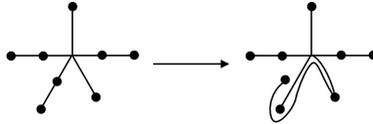


FIGURE 6. An example of a $*$ -orbit

The data of a periodic orbit will usually be written in the form above: the partition $A \cup B \cup C$ is denoted by letters above the elements $(0, s)$ of \mathcal{L} , and the integers N_r can be deduced from the elements of \mathcal{L} . For simplicity, it is not explicitly noted that $(3, 0) \rightarrow (0, 0)$. Note that by $*$ d), the cycle notation of π always ends $(m, 0) \rightarrow (2m, 0) \rightarrow \dots \rightarrow (n - m, 0)$. The reader seeking further clarification should consult Examples 2.12, 2.21, and 2.30, where other $*$ -orbits are described in this way.

Write $\tau_1, \tau_2: \mathcal{L} \rightarrow \mathbb{N}$ for the projections of \mathcal{L} onto the first and second components respectively; and write $\pi_1 = \tau_1 \circ \pi$ and $\pi_2 = \tau_2 \circ \pi$. Sometimes, with an abuse of notation, π_1^{-1} and π_2^{-1} will be used to denote $\tau_1 \circ \pi^{-1}$ and $\tau_2 \circ \pi^{-1}$ respectively.

Given $d \in \mathcal{D}_n$, one can ask whether or not there exists $P \in \mathcal{P}(m/n, k, \gamma)$ with $d(P) = d$. This imposes obvious conditions on d , which are expressed by the following result. While hardly concise, the conditions are in an ideal form for later use. If $d \in \mathcal{D}_n$ satisfies these conditions, it will be said that d is *legal data*.

Lemma 2.9. *An element $((N_r), \pi, (A, B, C))$ of \mathcal{D}_n is equal to $d(P)$ for some $P \in \mathcal{P}(m/n, k, \gamma)$ if and only if the following conditions hold:*

LD a) For all $r > 0$:

- i) $\pi_1(r, s) = r + {}_n m$ for all s , and $\pi_2(r, s)$ is increasing in s .*
- ii) $\pi_2(r, 0) = 0$.*

LD b) For $r = 0$:

- i) $0 \in \gamma$.*
- ii) $\pi_1(0, 0) = k$.*
- iii) If $c \in \{0, \dots, N_0 - 1\}$ is the unique element of C , then $\pi(0, c) = (m, 0)$.*
- iv) $\pi_1(0, s) \in \{k, \dots, m\}$ for all s , and is increasing in s .*
- v) If $s_1 < s_2$, $\pi_1(0, s_1) = \pi_1(0, s_2)$ and $s_1 \in B \cup C$ then $s_2 \in B$.*
- vi) If $s_1 < s_2$, $\pi_1(0, s_1) = \pi_1(0, s_2)$, and $s_1, s_2 \in A$ (respectively $s_1, s_2 \in B$) then $\pi_2(0, s_1) > \pi_2(0, s_2)$ (respectively $\pi_2(0, s_1) < \pi_2(0, s_2)$).*

Sketch Proof. The necessity of the conditions is obvious. Their sufficiency can be shown as follows: divide e_0 into $2m + 1$ subintervals, each mapped by $f_{m/n}$ over exactly one edge of Γ_n ; construct a Markov graph for $f_{m/n}$ using these subintervals and the edges e_r with $r > 0$ as the Markov partition. Given k, γ , and an element $d = ((N_r), \pi, (A, B, C))$ of \mathcal{D}_n satisfying LD a) and LD b), conditions a)i), b)iii) and b)iv) ensure that there are exactly two loops in the Markov graph which are compatible with the first component π_1 of π and with the partition (A, B, C) (two since the element $(0, s)$ of \mathcal{L} with $s \in C$ corresponds to the common endpoint of two intervals in the partition). Pick either of these loops, and let P be the periodic orbit of $f_{m/n}$ corresponding to it. Then conditions b)i) and b)ii) ensure that it belongs to $\mathcal{P}(m/n, k, \gamma)$, conditions b)iv)–vi) ensure that the data of P agree with the second component π_2 of π , and condition a)ii) ensures that P is a $*$ -orbit. \square

Example 2.10. In this example, Lemma 2.9 will be used to show that there is a bijection

$$\psi: \mathcal{P}(1/3, 0, B) \rightarrow \mathcal{P}(1/2, 0, B)$$

given by deleting all occurrences of $(2, s)$ in the cycle representation of π . Thus if $d(P) = ((N_0, N_1, N_2), \pi, (A, B, C))$, then $d(\psi(P)) = ((N_0, N_1), \pi', (A, B, C))$, where π' is obtained from π by deleting each occurrence of $(2, s)$ in its cycle representation.

ψ is first shown to be well defined: if $d(P) = ((N_0, N_1, N_2), \pi, (A, B, C))$ is legal data, then so is $((N_0, N_1), \pi', (A, B, C))$. Notice that the legality of $d(P)$ implies that $N_1 = N_2$, that $\pi(1, s) = (2, s)$ for all $s < N_1$, and that $\pi_1(2, s) = 0$ for all s . Hence

$$\pi'(r, s) = \begin{cases} \pi(r, s) & \text{if } r = 0 \\ \pi^2(r, s) & \text{if } r = 1. \end{cases}$$

Now condition LD a) holds since $\pi'_1(1, s) = \tau_1 \circ \pi^2(1, s) = 0$, $\pi'_2(1, s) = \tau_2 \circ \pi^2(1, s)$ is increasing in s by condition LD a)i) for $d(P)$, and $\pi'(1, 0) = \pi^2(1, 0) = (0, 0)$. Condition LD b) holds because $\pi' = \pi$ when $r = 0$.

Define $\phi: \mathcal{P}(1/2, 0, B) \rightarrow \mathcal{P}(1/3, 0, B)$ by following every $(1, s)$ with $(2, s)$ in the cycle representation of π . That is, if $d(P) = ((N_0, N_1), \pi, (A, B, C))$, then $d(\phi(P)) = ((N_0, N_1, N_1), \pi', (A, B, C))$, where π' is obtained from π by following each occurrence of $(1, s)$ with $(2, s)$ in its cycle notation. It can be shown as above that ϕ is well defined, and it is clearly an inverse to ψ .

This sort of argument is routine and tedious, and will be abbreviated when it occurs in earnest in the proof of Theorem 2.20.

2.3. The Train Track Condition. The main question addressed in this paper is: for which $P \in \mathcal{P}(m/n)$ is Γ_n^P itself a train track for the isotopy class of $F_{m/n}$ relative to P ? The essential property of a train track is *efficiency*: intuitively, this says that at any point x of Γ_n^P at which $(f_{m/n}^P)^i$ is not locally injective for some i , the image $(f_{m/n}^P)^i(U)$ of a small neighbourhood U of x ‘wraps around’ a point of P , and hence cannot be ‘pulled tight’ without passing through P . Efficiency can be detected as follows. By condition *d), the only point at which such pulling tight could occur is $(m, 0)$. For each point p of P which is not an endpoint of Γ_n^P (i.e. $p \neq (r, 0)$ for $0 \leq r < n$), an arc passing on one side of p will wrap around $(m, 0)$ under the appropriate iterate of $f_{m/n}^P$, while an arc passing on the other side will not, and its image can be pulled tight. Provided all of the arcs in the image $f_{m/n}^P(\Gamma_n^P)$ of Γ_n^P pass on the ‘correct’ side of each point of P , the map $f_{m/n}^P$ is efficient, and hence is a train track map for the isotopy class of $F_{m/n}: D^2 \setminus P \rightarrow D^2 \setminus P$.

The first part of the next definition gives a partition of $P \setminus \{(r, 0) : 0 \leq r < n\}$ into two sets, α and β : for those points p of P in α , an arc passing p to the left (according to the orientation of the edge e_r containing p) wraps around $(m, 0)$; while for those in β , an arc passing to the right wraps around $(m, 0)$. It is then possible to make a combinatorial definition of efficiency, as given by conditions TT a)–TT d) below.

Definitions 2.11. Let $P \in \mathcal{P}(m/n, k, \gamma)$ have data $d(P) = ((N_r), \pi, (A, B, C))$. Define a partition

$$\mathcal{L} \setminus \{(r, 0) : 0 \leq r < n\} = \alpha \cup \beta$$

inductively as follows. $\pi^{-1}(m, 0) \in \beta$. For each i with $2 \leq i \leq \#P - n$, the two elements $\pi^{-i}(m, 0)$ and $\pi^{-i+1}(m, 0)$ are in different sets if $\tau_1(\pi^{-i}(m, 0)) = 0$ and $\tau_2(\pi^{-i}(m, 0)) \in A$, and are in the same set otherwise.

P is a *train track orbit* (or P is *TT*) if and only if the following conditions hold:

TT a) For all $0 \leq s < N_0 - 1$,

$$\pi(0, s) \in \begin{cases} \alpha & \text{if } s \in B \\ \beta & \text{if } s \in A. \end{cases}$$

TT b) If $k < r < m$ then $N_{r+n-m} = 1$.

TT c) If $\pi_2(k + n - m, s) > \pi_2(0, 0)$ for some $s > 0$, then $\pi(k + n - m, s) \in \beta$.

TT d) If $\gamma = A$ then $\pi_2(k + n - m, s) > \pi_2(0, 0)$ for all $s > 0$.

The set of all train track orbits in $\mathcal{P}(m/n, k, \gamma)$, in $\mathcal{P}(m/n, k)$, and in $\mathcal{P}(m/n)$ will be denoted $TT(m/n, k, \gamma)$, $TT(m/n, k)$, and $TT(m/n)$.

Example 2.12. Let $P \in \mathcal{P}(1/3, 0, B)$ have data

$$(0, 0) \xrightarrow{B} (0, 1) \xrightarrow{A} (1, 1) \xrightarrow{\beta} (2, 1) \xrightarrow{\beta} (0, 2) \xrightarrow{C} (1, 0) \xrightarrow{\beta} (2, 0)$$

(see Fig. 7). The partition $\alpha \cup \beta$ is shown on the cycle representation of π . This partition is easy to write down: start with $(0, 2)$ (the unique element of C) in β , and move backwards through the permutation, switching from β to α or vice-versa at each occurrence of A .

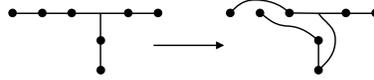


FIGURE 7. A periodic orbit satisfying the TT conditions

The TT conditions can easily be checked:

- a) $\pi(0, 0) \in \alpha$ and $\pi(0, 1) \in \beta$.
- b) Vacuous, since there are no r with $0 = k < r < m = 1$.
- c) $k+n-m = 0+3-1 = 2$, so the condition requires that $\pi(2, s) \in \beta$ whenever $\pi_2(2, s) > 1$: i.e. $\pi(2, 1) \in \beta$, which is true.
- d) Vacuous, since $\gamma = B$.

The TT conditions have intuitive motivations. TT a) ensures that arcs of $f_{m/n}^P(\Gamma_n^P)$ pass on the ‘correct’ side of images of points of $P \cap e_0$: for example, if $s \in B$ and $\pi_1(0, s) = r$, then an arc of $f_{m/n}^P(e_{r+n-m})$ passes to the left of $\pi(0, s)$ and hence $\pi(0, s)$ must lie in α to avoid a violation of efficiency. TT b) reflects the fact that for $k < r < m$, the image of e_0 passes images of points of $P \cap e_{r+n-m}$ on both sides: hence efficiency will be violated if there is any such point other than the one with image $(r, 0)$. TT c) reflects that the image of e_0 passes to the right of those images of points of $P \cap e_{k+n-m}$ which are further from the initial point of e_k than $\pi(0, 0)$; and TT d) is a stronger version of the same condition in the case $\gamma = A$, when the image of e_0 passes on both sides of points of e_k between the initial point and $\pi(0, 0)$.

The TT conditions above are stated in a way which reflects this intuitive motivation. However they can be replaced by equivalent alternatives which are easier to work with in practice. This will be done in Lemma 2.17 below.

The reader familiar with this approach to train tracks should accept without further argument that the TT conditions are the appropriate combinatorial expression of efficiency. However, a more precise statement is given by the following definition (which connects the intuitive notion of a train track with the definition due to Bestvina and Handel [1], and in particular, by introducing *peripheral loops* around the points of P , makes it possible to be precise about the concept of ‘wrapping around’ a point of P) and by Theorem 2.15. The rather contorted definition can most easily be understood by referring to Fig. 8.

Definitions 2.13. Let $P \in \mathcal{P}(m/n)$, and let $L \cup R$ be a partition of $\mathcal{L} \setminus \{(r, 0) : 0 \leq r < n\}$. The (L, R) -Bestvina-Handel star graph $BH(P, L, R) \subset T_n$ of P is defined as follows:

- a) There are $\#P + 1$ vertices; one vertex v contained in the central decomposition element of T_n , and $\#P$ vertices $\{v_{r,s} : (r, s) \in \mathcal{L}\}$ with $v_{r,s}$ contained in the same decomposition element of T_n as the point (r, s) of P . If $s > 0$ then $v_{r,s}$ is to the left or right of (r, s)

(with respect to the orientation of e_r) according as $(r, s) \in L$ or $(r, s) \in R$. (If $s = 0$ then $v_{r,s}$ can be chosen to the left or the right of (r, s) .)

- b) There are $\#P$ *peripheral edges* $\{p_{r,s} : (r, s) \in \mathcal{L}\}$: the peripheral edge $p_{r,s}$ has both endpoints at $v_{r,s}$, and forms a loop bounding a disk with the point (r, s) of P in its interior. The peripheral edges are small enough that no two of them intersect any given decomposition element of T_n .
- c) There are $\#P$ *main edges* $\{e_{r,s} : (r, s) \in \mathcal{L}\}$: the main edge $e_{r,s}$ goes from $v_{r,s}$ to $v_{r,s+1}$ if $s < N_r - 1$, and goes from $v_{r,s}$ to v if $s = N_r - 1$. The main edges are chosen so that each interval decomposition element of T_n contains at most one point of the union of the interiors of the main edges.

Since $BH(P, L, R)$ is a spine of $T_n \setminus P$, the thick tree map $F_{m/n}: T_n \rightarrow T_n$ induces a well defined homotopy class of graph maps $g: BH(P, L, R) \rightarrow BH(P, L, R)$. If g is required to restrict to a homeomorphism of the subgraph of peripheral edges, to send vertices to vertices, and to be locally injective away from its vertices, then the image edge-paths of g are also well-defined. $BH(P, L, R)$ is said to be a *Bestvina-Handel star train track* for P if such a graph map is a train track map: that is, if

- a) It is *absorbed*: the image edge path of each main edge begins and ends with main edges.
- b) It is *efficient*: there is no backtracking in the edge path $g^i(e_{r,s})$ for any $(r, s) \in \mathcal{L}$ and $i > 0$.

The $*$ -orbit $P \in \mathcal{P}(m/n)$ is said to *have a Bestvina-Handel star train track* if $BH(P, L, R)$ is a Bestvina-Handel star train track for some choice of the partition $\mathcal{L} \setminus \{(r, 0) : 0 \leq r < n\} = L \cup R$.

Remarks 2.14. a) If $BH(P, L, R)$ is efficient but not absorbed, then there is some other partition $L' \cup R'$ such that $BH(P, L', R')$ is both efficient and absorbed (obtained, for example, by applying the Bestvina-Handel operation of *absorbing into the peripheral subgraph*). Requiring train track maps to be absorbed as well as efficient means that if P has a Bestvina-Handel star train track, then there is a unique choice of partition $L \cup R$ such that $BH(P, L, R)$ is a Bestvina-Handel star train track (see the first paragraph of the proof of Theorem 2.15).

- b) It follows from the results of [1] that if P has a Bestvina-Handel star train track, and if the transition matrix for the main edges of that star train track is irreducible, then P has pseudo-Anosov braid type, and the train track yields a Markov partition for the pseudo-Anosov representative in the isotopy class of $F_{m/n}$ in $D^2 \setminus P$. Moreover, because the n edge germs at v are permuted by g they all lie in different gates: hence such a pseudo-Anosov must have an interior n -pronged singularity (whose prongs are rotated by m under the action of the pseudo-Anosov), 1-pronged singularities at each of the points of P , and a $(\#P - n)$ -pronged singularity at the boundary.

Theorem 2.15. *Let $P \in \mathcal{P}(m/n, k, \gamma)$. Then P has a Bestvina-Handel star train track if and only if P is TT.*

Sketch Proof. A straightforward argument shows that $BH(P, L, R)$ is absorbed if and only if $L = \beta$ and $R = \alpha$: the procedure for constructing inductively the unique partition $L \cup R$ for which $BH(P, L, R)$ is absorbed is identical to that by which the partition $\alpha \cup \beta$ is defined.

Thus it is only required to show that $BH(P, \beta, \alpha)$ is efficient if and only if P is TT. For each $u \neq v$ with $0 \leq u, v < n$, write $t_{u,v}$ for the edge-path $e_{u, N_u-1} \bar{e}_{v, N_v-1}$, and observe that

$BH(P, \beta, \alpha)$ is efficient if and only if the image edge-path $g(e_{r,s})$ of each main edge can be written in the form

$$g(e_{r,s}) = m_1 p_1 m_2 p_2 \dots, m_{k-1} p_{k-1} m_k,$$

where each p_i is a peripheral edge (or its inverse), and each m_i is either a main edge (or its inverse), or one of the edge-paths $t_{u,v}$. If this condition holds, then a straightforward induction shows that it holds also for $g^n(e_{r,s})$ for all $n > 0$, and hence there cannot be any cancellation in any of these edge-paths. If the condition fails, then the edge-path $g(e_{r,s})$ contains some word $e_{u,v} e_{u,v+1}$ (or its inverse), which under iteration yields the word $\bar{e}_{m,0} e_{m,0}$.

If P is not TT, then one of the conditions TT a) – TT d) fails. In each case, an argument arising from the intuitive motivation of the conditions can be used to show that there is some $g(e_{r,s})$ which contains a word of the form $e_{u,v} e_{u,v+1}$ or its inverse. Hence $BH(P, \beta, \alpha)$ is not efficient.

For the converse, suppose that P is TT. If $BH(P, \beta, \alpha)$ is not efficient, then there is some edge $e_{u,v}$ such that $g(e_{u,v})$ contains a word of the form $e_{r,s-1} e_{r,s}$ or $\bar{e}_{r,s} \bar{e}_{r,s-1}$. This implies that $u = 0$ or $u \geq k + n - m$ (i.e. that $0 \leq r \leq m$), since otherwise $g(e_{u,v}) = e_{u+n,m,v}$ is an edge-path of length 1.

There are two cases to consider. First suppose that $\pi_1^{-1}(r, s) \neq 0$. Then TT b) gives $r = k$ (and hence $\pi_1^{-1}(r, s) = k + n - m$). Thus $u = 0$, since only e_0 and e_{k+n-m} have images intersecting e_k . If $\gamma = B$ then, since $g(e_{0,v})$ passes to the right of (r, s) , it can only contain the word $e_{k,s-1} e_{k,s}$ if $\pi_2(0, v) < s$ and $(r, s) \in \alpha$. This contradicts TT c). If $\gamma = A$ then a similar contradiction to TT c) and TT d) arises.

The case where $\pi_1^{-1}(r, s) = 0$ can be treated similarly using TT a). □

Example 2.16. The Bestvina-Handel and Thurston train tracks for the TT orbit of Example 2.12 are shown in Fig. 8. The Bestvina-Handel train track is obtained by replacing each point of the orbit with a peripheral loop, and putting the vertex of this loop on the left or right according as the point of the orbit belongs to β or α . The Thurston train track is obtained from the Bestvina-Handel train track using the techniques of [1] (in the case of $*$ -orbits, this is simply a matter of replacing the central valence n vertex with an n -gon, and each of the loops with a 1-gon, as shown in the figure).

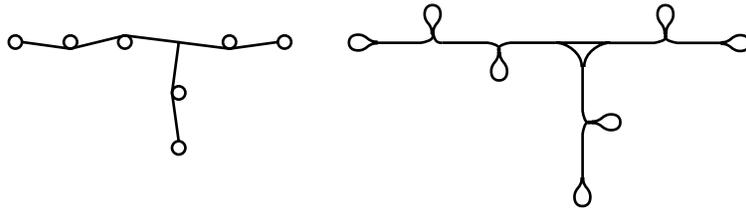


FIGURE 8. Bestvina-Handel and Thurston Train Tracks for the example of Fig. 7

As mentioned above, there is an alternative form of the TT conditions which is often easier to work with:

Lemma 2.17. *Consider the conditions*

TT' a) $\#A = 1$, and if $A = \{s\}$ and $j > 0$ is least such that $\tau_1(\pi^j(0, s)) = 0$, then $C = \{\tau_2(\pi^j(0, s))\}$.

TT' d) If $\gamma = A$ then $\pi(0, 0) = (k, 1)$.

Then $TTa)$ and $TT'a)$ are equivalent; and if they hold then $TTd)$ and $TT'd)$ are equivalent. Moreover, if $TTa)$ and $TTc)$ hold then there is at most one s with $\pi_2(k+n-m, s) \geq \pi_2(0, 0)$.

Proof. The equivalence of $TTa)$ and $TT'a)$ follows easily from the definition of the partition $\alpha \cup \beta$: note that $TT'a)$ says simply that reading the partition $A \cup B \cup C$ along the cycle notation of π , one sees $B^{N_0-2}AC$. In particular, since the cycle notation starts with $(0, 0)$, this implies that if $\gamma = A$ then $N_0 = 2$: since $\pi_1(0, 1) = m$ and $\pi_1(0, 0) = k < m$, the equivalence of $TTd)$ and $TT'd)$ follows. For the final statement, note that the elements of β are precisely those which lie between the element of A and the element and the element of C in the cycle notation of π , and only one such can lie in e_k : also it is not possible that $\pi_2(k+n-m, s) = \pi_2(0, 0)$, since this would imply $\pi(k+n-m, s) = \pi(0, 0)$. \square

If $P \in TT(m/n)$, then it follows from the results of [1] that the set of braid types forced by $\text{bt}(P)$ is precisely the set of braid types of periodic orbits of $f_{m/n}$ whose span is contained in the span of P . In particular, Lemma 2.3 gives:

Theorem 2.18. *The set $\{\text{bt}(P, F_{m/n}) : P \in TT(m/n)\}$ is totally ordered by the forcing relation.*

2.4. The TT condition in the horseshoe. The elements of $TT(1/2) = TT(1/2, 0, B)$ were calculated in [11]: in this case the partition $\alpha \cup \beta$ is precisely that defined on page 880 of [11], and conditions $TTa)$ and $TTc)$ are equivalent to the condition that π has *no bogus transitions*, also on page 880 of [11] (while $TTb)$ and $TTd)$ are vacuous). Theorem 2.1 and lemma 2.4 of [11] then give the following result (in which the words c_q are as defined in Section 1.1).

Theorem 2.19. *The function $\mathbb{Q} \cap (0, 1/2) \rightarrow \mathcal{P}(1/2, 0, B)$ which takes a rational q to the periodic orbit P_q with code $c_{q_1}^0$ is a bijection onto $TT(1/2, 0, B)$. If $r/s \in (0, 1/2)$, the data $d(P_{r/s}) = ((N_0, N_1), \pi, (A, B, C))$ of $P_{r/s}$ satisfies $N_0 = s - r + 1$ and $N_1 = r + 1$.*

2.5. Renormalizing horseshoe TT orbits (the case $k = m - 1$). Theorem 2.19 gives a relatively straightforward approach to determining the elements of $TT(m/n, m - 1, B)$ for each $m/n \in (0, 1/2)$: the next result defines a ‘renormalization operator’ ϕ , which is a bijection from the set $\mathcal{P}(1/2, 0, B)$ of horseshoe periodic orbits to the set $\mathcal{P}(m/n, m - 1, B)$ with the property that $\phi(P)$ is TT if and only if P is TT. Example 2.10 is the simplest case of this operator, when $m/n = 1/3$.

For notational simplicity, m/n will be taken to be a fixed rational in $(0, 1/2)$ throughout this section, and the dependence of some objects (such as ϕ) upon it will be dropped.

Theorem 2.20. *There is a bijection $\phi: \mathcal{P}(1/2, 0, B) \rightarrow \mathcal{P}(m/n, m - 1, B)$ given by replacing each occurrence of $(0, s)^\gamma$ (where $\gamma \in \{A, B, C\}$) in the cycle representation of π by*

$$(m - 1, s) \rightarrow (2m - 1, s) \rightarrow \cdots \rightarrow (n - m, s) \rightarrow (0, s)^\gamma$$

and each occurrence of $(1, s)$ by

$$(m, s) \rightarrow (2m, s) \rightarrow \cdots \rightarrow (n - m - 1, s) \rightarrow (n - 1, s).$$

Moreover, ϕ restricts to a bijection $TT(1/2, 0, B) \rightarrow TT(m/n, m - 1, B)$.

Example 2.21. Consider the periodic orbit $P \in \mathcal{P}(1/2, 0, B)$ given by

$$(0, 0) \xrightarrow{B} (0, 1) \xrightarrow{A} (1, 1) \xrightarrow{C} (0, 2) \xrightarrow{\beta} (1, 0)$$

(see Fig. 9).



FIGURE 9. A periodic orbit in the horseshoe

Then $\phi(P) \in \mathcal{P}(2/5, 1, B)$ is given by

$$\begin{array}{ccccccccccc} (1, 0) & \rightarrow & (3, 0) & \xrightarrow{B} & (0, 0) & \rightarrow & (1, 1) & \rightarrow & (3, 1) & \xrightarrow{A} & (0, 1) & \rightarrow & (2, 1) & \rightarrow & (4, 1) \\ \underbrace{\hspace{10em}}_{\text{from } (0,0)} & & \underbrace{\hspace{10em}}_{\text{from } (0,1)} & & \underbrace{\hspace{10em}}_{\text{from } (0,1)} & & \underbrace{\hspace{10em}}_{\text{from } (1,1)} \\ & & & & \xrightarrow{\beta} & & \xrightarrow{\beta} & & \xrightarrow{C} & & \xrightarrow{\beta} & & \xrightarrow{\beta} & & \xrightarrow{\beta} \\ & & & & \underbrace{\hspace{10em}}_{\text{from } (0,2)} & & \underbrace{\hspace{10em}}_{\text{from } (1,0)} \end{array}$$

(see Fig. 10). It can easily be checked that both of these orbits are TT. The figures also show the idea of the construction: the pattern of the periodic orbit of $f_{1/2}$ is replicated in the edges e_{m-1} and e_m of Γ_n .

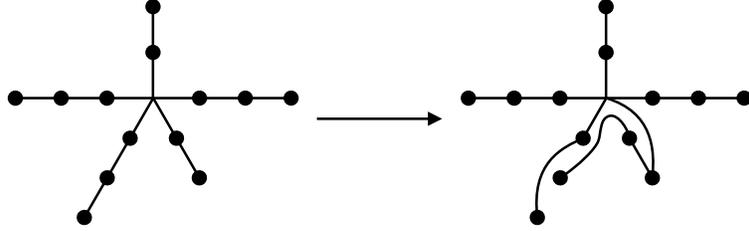


FIGURE 10. The periodic orbit in $\mathcal{P}(2/5, 1, B)$ obtained from that of Fig. 9

Proof. The first step is to show that the construction yields a well-defined function $\phi: \mathcal{P}(1/2, 0, B) \rightarrow \mathcal{P}(m/n, m-1, B)$: that is, that it transforms legal data to legal data.

Suppose, then, that $P \in \mathcal{P}(1/2, 0, B)$ has data $d(P) = ((N_0, N_1), \pi, (A, B, C))$. Since $\mathcal{O}_{m/n}[m-1, 0]$ and $\mathcal{O}_{m/n}[m, n-1]$ partition \mathbb{Z}_n , if the construction does give an element P' of $\mathcal{P}(m/n, m-1, B)$ it must have data

$$d(P') = ((N'_r), \pi', (A, B, C)),$$

where

$$N'_r = \begin{cases} N_0 & \text{if } r \in \mathcal{O}_{m/n}[m-1, 0] \\ N_1 & \text{if } r \in \mathcal{O}_{m/n}[m, n-1], \end{cases}$$

and

$$\pi'(r, s) = \begin{cases} (\pi_1(0, s) + m - 1, \pi_2(0, s)) & \text{if } r = 0 \\ (\pi_1(1, s) + m - 1, \pi_2(1, s)) & \text{if } r = n - 1 \\ (m +_n r, s) & \text{otherwise.} \end{cases}$$

(The sets A , B , and C are unchanged by construction.)

It is routine to show that conditions LD a) and LD b) of Lemma 2.9 hold for this data, and hence that $\phi: \mathcal{P}(1/2, 0, B) \rightarrow \mathcal{P}(m/n, m-1, B)$ is well defined as required. To show that it is a bijection, observe that if $P \in \mathcal{P}(m/n, m-1, B)$ has data $d(P) = ((N_r), \pi, (A, B, C))$, then conditions LD a)i) and LD b)iv) ensure that every occurrence of $(m-1, s)$ in the cycle representation of π is followed by

$$(m-1, s) \rightarrow (2m-1, s) \rightarrow \cdots \rightarrow (n-m, s) \rightarrow (0, s)^\gamma$$

for some $\gamma \in \{A, B, C\}$, and every occurrence of (m, s) is followed by

$$(m, s) \rightarrow (2m, s) \rightarrow \cdots \rightarrow (n-m-1, s) \rightarrow (n-1, s).$$

Replacing the first type of block with $(0, s)^\gamma$ and the second with $(1, s)$ thus defines an inverse $\psi: \mathcal{P}(m/n, m-1, B) \rightarrow \mathcal{P}(1/2, 0, B)$ of ϕ (checking that the data so obtained is legal is again routine).

Finally, it must be shown that the bijection ϕ preserves the TT conditions. Let $P \in \mathcal{P}(1/2, 0, B)$ and $P' = \phi(P) \in \mathcal{P}(m/n, m-1, B)$. Since the sets A , B , and C are the same for P and P' , the condition TT'a) is satisfied either for neither or for both of P and P' . Conditions TT b) and TT d) are vacuous for both P and P' . Condition TT c) for P reads

$$\pi(1, s) \in \beta \text{ whenever } \pi_2(1, s) > \pi_2(0, 0),$$

while for P' it reads

$$\pi'(n-1, s) \in \beta \text{ whenever } \pi'_2(n-1, s) > \pi'_2(0, 0).$$

Since $\pi'(n-1, s) \in \beta$ if and only if $\pi(1, s) \in \beta$; $\pi'_2(n-1, s) = \pi_2(1, s)$; and $\pi'_2(0, 0) = \pi_2(0, 0)$ (these last two by the expression for π' above), TT c) is also satisfied either for neither or for both of P and P' . \square

The next step is to use Lemma 1.17 to identify the orbits $\phi(P_{r/s})$ which make up $TT(m/n, m-1, B)$. The function $\xi_{u/v, p/q}$ in the statement of the next theorem is the one defined in Lemma 1.23.

Theorem 2.22. *Let $u/v = LFP(m/n)$ and $p/q = RFP(m/n)$. Let $r/s \in (0, 1/2)$. Then $\phi(P_{r/s})$ has height $\xi_{u/v, p/q}(r/s)$ and decoration $w_{m/n}$.*

Proof. By Theorem 2.19, $\phi(P_{r/s})$ has period

$$(s-r+1)\#\mathcal{O}_{m/n}[m-1, 0] + (r+1)\#\mathcal{O}_{m/n}[m, n-1],$$

which, by Lemma 1.22, is equal to $(s-r+1)v + (r+1)q$. It will be shown that $\phi(P_{r/s})$ has rotation interval

$$\left[\frac{(s-r)u + rp}{(s-r)v + rq}, \frac{u+p}{v+q} \right] = \left[\xi_{u/v, p/q} \left(\frac{r}{s} \right), \frac{m}{n} \right],$$

which will establish the result by Lemmas 1.17 and 2.4.

If P is a $*$ -orbit with data $d(P) = ((N_r), \pi, (A, B, C))$, then the rotation number of P about the fixed point is given by the number of times it cycles around Γ_n divided by its period: that is,

$$\rho(P) = \frac{\sum_{r=n-m}^{n-1} N_r}{\sum_{r=0}^{n-1} N_r}.$$

Let $P \in \mathcal{P}(1/2, 0, B)$. Then $\rho(P) = N_1/(N_0 + N_1)$, and if $\phi(P)$ has data $((N'_r), \pi, (A, B, C))$ then by Lemma 1.22, and the fact that $N'_r = N_0$ for all $r \in \mathcal{O}_{m/n}[m-1, 0]$ and $N'_r = N_1$ for all $r \in \mathcal{O}_{m/n}[m, n-1]$,

$$\frac{\sum_{r \in \mathcal{O}_{m/n}[m, n-1], r \geq n-m} N'_r}{\sum_{r \in \mathcal{O}_{m/n}[m, n-1]} N'_r} = \frac{p}{q} \quad \text{and} \quad \frac{\sum_{r \in \mathcal{O}_{m/n}[m-1, 0], r \geq n-m} N'_r}{\sum_{r \in \mathcal{O}_{m/n}[m-1, 0]} N'_r} = \frac{u}{v},$$

and hence

$$\rho(\phi(P)) = \frac{N_0 u + N_1 p}{N_0 v + N_1 q}.$$

If P has period b and rotation number a/b then $N_0 = b - a$ and $N_1 = a$, and so

$$\rho(\phi(P)) = \frac{(b-a)u + ap}{(b-a)v + aq} = \xi_{u/v, p/q}(a/b).$$

Now the rotation interval $[r/s, 1/2]$ of $P_{r/s}$ is the set of rotation numbers of periodic orbits of $f_{1/2}$ contained in the span of $P_{r/s}$, and, by a theorem of Boyland [4], for each such rotation number a/b there exists such an orbit with period b . Likewise, the rotation interval of $\phi(P_{r/s})$ is the set of rotation numbers of periodic orbits of $f_{m/n}$ contained in the span of $\phi(P_{r/s})$; but these orbits are precisely the images under ϕ of those defining the rotation interval of $P_{r/s}$. Hence the rotation interval of $\phi(P_{r/s})$ is

$$\left\{ \xi_{u/v, p/q} \left(\frac{a}{b} \right) : \frac{a}{b} \in \left[\frac{r}{s}, \frac{1}{2} \right] \right\} = \xi_{u/v, p/q} \left(\left[\frac{r}{s}, \frac{1}{2} \right] \right) = \left[\xi_{u/v, p/q} \left(\frac{r}{s} \right), \frac{m}{n} \right]$$

as required. \square

Remark 2.23. If $n \geq 2$ then $\mathcal{P}(1/n) = \mathcal{P}(1/n, 0, B)$, and hence the results of this section give a complete description of $TT(1/n)$: there is a bijection $q \mapsto P_q^{1/n}$ from $\mathbb{Q} \cap (0, 1/n)$ to $TT(1/n)$ such that $P_q^{1/n}$ has height q and decoration $w_{1/n} = 0^{n-3}$. In other words, $P_q^{1/n}$ has the same braid type as the horseshoe orbits of code $c_{q_1^0 0^{n-3} 1}$.

In general, since $\xi_{u/v, p/q} : (0, 1/2) \rightarrow (u/v, m/n)$ is an increasing bijection, the results of this section yield a bijection $q \mapsto P_q^{m/n}$ from $(\text{LFP}(m/n), m/n) \cap \mathbb{Q}$ to $TT(m/n, m-1, B)$ with the property that $P_q^{m/n}$ has height q and decoration $w_{m/n}$.

2.6. Admissible k . Condition TT b) gives restrictions on the values of k for which $TT(m/n, k)$ is non-empty: in this section it will be shown that the number of such *admissible* values of k is equal to the length of the left Farey sequence of m/n .

Definition 2.24. Let $m/n \in (0, 1/2)$. An integer $k \in [0, m-1]$ is *m/n -admissible* if

$$[k+1, m-1] \cap \mathcal{O}_{m/n}[k, 0] = \emptyset.$$

The set of all m/n -admissible integers k is denoted $\mathcal{A}_{m/n}$.

Lemma 2.25. Let $m/n \in (0, 1/2)$ and $k \in [0, m-1]$. If $k \notin \mathcal{A}_{m/n}$, then $TT(m/n, k)$ is empty.

Proof. Suppose $P \in TT(m/n, k)$ has data $d(P) = ((N_r), \pi, (A, B, C))$. Then by TT b), $N_{r+n-m} = 1$ for all r with $k < r < m$. Applying LD a) i) inductively, it follows that $N_s = 1$ for all $s \in \bigcup_{k < r < m} \mathcal{O}_{m/n}[m, r+n-m]$. However $N_k > 1$, since $\pi_2(0, 0) = \pi_2(k+n-m, 0) = k$. Thus $TT(m/n, k)$ must be empty unless $k \notin \bigcup_{k < r < m} \mathcal{O}_{m/n}[m, r+n-m]$. However, this condition is equivalent to $k \in \mathcal{O}_{m/n}[r, 0]$ for all r with $k < r < m$; which is in turn equivalent to $r \notin \mathcal{O}_{m/n}[k, 0]$ for all r with $k < r < m$: that is, to m/n -admissibility. \square

The remainder of this section is devoted to studying the structure of the set $\mathcal{A}_{m/n}$. It will be seen in Section 2.7 that the converse to Lemma 2.25 is also true: if $k \in \mathcal{A}_{m/n}$, then $TT(m/n, k)$ is non-empty.

Notice that $0, m-1 \in \mathcal{A}_{m/n}$ for all m/n : in particular $\mathcal{A}_{1/n} = \{0\}$. Lemma 2.26 below connects $\mathcal{A}_{m/n}$ with $\mathcal{A}_{\text{LFP}(m/n)}$, and hence yields an inductive description of $\mathcal{A}_{m/n}$ for all m/n .

Let $m/n \in (0, 1/2)$ with $\text{LFP}(m/n) = u/v$, and write $R_{m/n} = \mathcal{O}_{m/n}[m, n-1]$. By Lemma 1.22, $R_{m/n}$ has cardinality $n-v$. Moreover, $R_{m/n} \cap \mathcal{A}_{m/n} = \emptyset$: for if $k \in R_{m/n}$, then $\mathcal{O}_{m/n}[k, 0]$ contains the complement of R , and in particular contains $m-1$. Define a function $\psi_{m/n}: \mathbb{Z}_n \rightarrow \mathbb{Z}_v$ by

$$\psi_{m/n}(k) = k - \#(R_{m/n} \cap [0, k]).$$

Notice that

$$\psi_{m/n}(k) = \begin{cases} \psi_{m/n}(k-1) + 1 & \text{if } k \notin R_{m/n} \\ \psi_{m/n}(k-1) & \text{if } k \in R_{m/n} \end{cases}$$

It can easily be seen that $\psi_{m/n}(0) = 0$, $\psi_{m/n}(m-1) = \psi_{m/n}(m) = u$ (using Lemma 1.22), and $\psi_{m/n}(n-1) = v-1$. In particular, $\psi_{m/n}$ is an increasing surjection.

Lemma 2.26. *Let $m/n \in (0, 1/2)$ with $m > 1$, $u/v = \text{LFP}(m/n)$, and $k \in [0, m-2] \setminus R_{m/n}$. Then $k \in \mathcal{A}_{m/n}$ if and only if $\psi_{m/n}(k) \in \mathcal{A}_{u/v}$.*

Proof. Observe that $\psi_{m/n}|_{\mathbb{Z}_n \setminus R_{m/n}}: \mathbb{Z}_n \setminus R_{m/n} \rightarrow \mathbb{Z}_v$ is a bijection which conjugates the function $+_{u/v}: \mathbb{Z}_v \rightarrow \mathbb{Z}_v$ given by $+_{u/v}(k) = k +_v u$ to the function $\tilde{+}_{m/n}: \mathbb{Z}_n \setminus R_{m/n} \rightarrow \mathbb{Z}_n \setminus R_{m/n}$ given by

$$\tilde{+}_{m/n}(k) = \begin{cases} k +_n m & \text{if } k \neq 0 \\ m-1 & \text{if } k = 0. \end{cases}$$

Since $k \in [0, m-2]$ and $m-1 \notin R_{m/n}$, it follows that $\psi_{m/n}(k) \leq \psi_{m/n}(m-1) - 1 = u-1$. Now

$$\begin{aligned} k \in \mathcal{A}_{m/n} &\iff r \notin \mathcal{O}_{m/n}[k, 0] \text{ for } k < r < m \\ &\iff r \notin \mathcal{O}_{\tilde{+}_{m/n}}[k, 0] \text{ for } k < r < m \text{ and } r \notin R_{m/n} \\ &\iff \psi_{m/n}(r) \notin \mathcal{O}_{u/v}[\psi_{m/n}(k), \psi_{m/n}(0)] \text{ for } k < r < m \text{ and } r \notin R_{m/n} \\ &\iff s \notin \mathcal{O}_{u/v}[\psi_{m/n}(k), 0] \text{ for } \psi_{m/n}(k) < s < \psi_{m/n}(m) = u \\ &\iff \psi_{m/n}(k) \in \mathcal{A}_{u/v} \end{aligned}$$

(where the penultimate equivalence uses that $R_{m/n}$ contains neither k nor $m-1$). \square

It follows that $\mathcal{A}_{m/n} = \{m-1\} \cup \psi_{m/n}|_{\mathbb{Z}_n \setminus R_{m/n}}^{-1}(\mathcal{A}_{u/v})$, and in particular that $\mathcal{A}_{m/n}$ has one more element than $\mathcal{A}_{u/v}$. Since $\mathcal{A}_{1/n} = \{0\}$ for all n , the following result holds:

Corollary 2.27. *Let $m/n \in (0, 1/2)$. Then $\mathcal{A}_{m/n}$ has the same cardinality as $\text{LFS}(m/n)$.*

The following result will be important in the next section.

Lemma 2.28. *Let $\text{LFS}(m/n) = (0 = u_1/v_1, u_2/v_2, \dots, u_\alpha/v_\alpha)$, and let the elements of $\mathcal{A}_{m/n}$ be $0 = k_1 < k_2 < \dots < k_\alpha = m - 1$. Then $\mathcal{O}_{m/n}[k_i, 0]$ has cardinality v_i for each i . In particular, the elements of $\mathcal{A}_{m/n}$ appear in decreasing order along the orbit $\mathcal{O}_{m/n}[m, 0]$.*

Proof. That $\mathcal{O}_{m/n}[m - 1, 0]$ has cardinality v_α follows from Lemma 1.22. The proof of Lemma 2.26 shows that if $k \in \mathcal{A}_{m/n}$ is distinct from $m - 1$, then $\mathcal{O}_{u/v}[\psi_{m/n}(k), 0] = \psi_{m/n}(\mathcal{O}_{m/n}[k, 0])$, and the result follows by induction. \square

2.7. Train track orbits with $k < m - 1$. Let $m/n \in (0, 1/2)$ with $m > 1$, and $\text{LFS}(m/n) = (0 = u_1/v_1, u_2/v_2, \dots, u_\alpha/v_\alpha)$. By Corollary 2.27, $\mathcal{A}_{m/n}$ has α elements, which will be denoted $0 = k_1 < k_2 < \dots < k_\alpha = m - 1$. Let $i < \alpha$, and define

$$\begin{aligned} R &= \mathcal{O}_{m/n}[m, k_{i+1} + n - m] \\ S &= \mathcal{O}_{m/n}[k_{i+1}, k_i + n - m] \\ T &= \mathcal{O}_{m/n}[k_i, 0]. \end{aligned}$$

The dependence of the sets on m/n and i is suppressed, since these two variables will remain fixed throughout the section. R , S , and T are mutually disjoint by Lemma 2.28, and hence define a partition of \mathbb{Z}_n . The cardinalities of T and S are given by Lemma 2.28, and the cardinality of R can therefore be deduced:

$$\begin{aligned} \#R &= n - v_{i+1} \\ \#S &= v_{i+1} - v_i \\ \#T &= v_i. \end{aligned}$$

The following theorem gives a complete description of $TT(m/n, k_i, B)$:

Theorem 2.29. *There is a bijection from $(0, 1) \cap \mathbb{Q}$ to $TT(m/n, k_i, B)$, defined as follows: if $p/q \in (0, 1)$, then the element $P_{i,p/q,m/n}$ of $TT(m/n, k_i, B)$ corresponding to p/q has data $d(P_{i,p/q,m/n}) = ((N_r), \pi, (A, B, C))$ given by: $A = \{p\}$, $B = \{0, \dots, q - 1\} \setminus \{p\}$, $C = \{q\}$,*

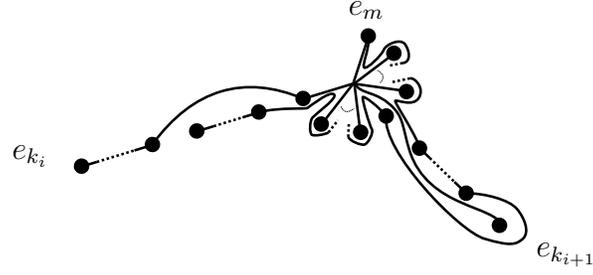
$$N_r = \begin{cases} 1 & \text{if } r \in R \\ q + 1 - p & \text{if } r \in S \\ q + 1 & \text{if } r \in T, \end{cases}$$

and

$$\begin{aligned} \pi(0, s) &= \begin{cases} (k_i, q - p + s) & \text{if } 0 \leq s < p \\ (k_{i+1}, q - p) & \text{if } s = p \\ (k_{i+1}, s - p) & \text{if } p + 1 \leq s < q \\ (m, 0) & \text{if } s = q \end{cases} \\ \pi(k_i + n - m, q - p) &= (k_i, q) \\ \pi(r, s) &= (r +_n m, s) \quad \text{for all other } (r, s). \end{aligned}$$

A schematic representation of $P_{i,p/q,m/n}$ is shown in Fig. 11, which depicts the image of e_0 and of e_{k_i+n-m} through to e_{n-1} . e_0 contains $q + 1$ points of P , of which p are mapped to e_{k_i} (the other $q + 1 - p$ points of P on e_{k_i} are images of points of P on e_{k_i+n-m}); $q - p$ are mapped to $e_{k_{i+1}}$ in the configuration shown; and one is mapped to e_m .

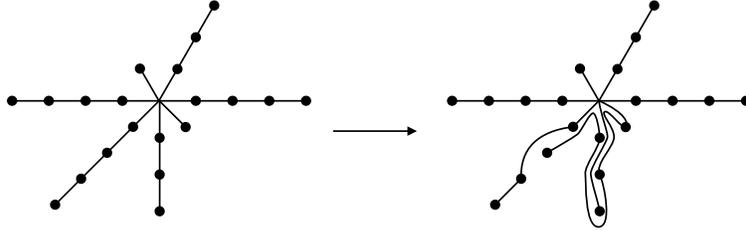
Example 2.30. Let $m/n = 3/7$. Now $\text{LFS}(3/7) = (0, 1/3, 2/5)$ has three elements, and hence $\mathcal{A}_{3/7} = \{0, 1, 2\}$, with $k_1 = 0$, $k_2 = 1$, and $k_3 = 2$. Pick $i = 2$. Then $R = \mathcal{O}_{3/7}[3, 2 + 7 - 3] = \{3, 6\}$, $S = \mathcal{O}_{3/7}[2, 1 + 7 - 3] = \{2, 5\}$, and $T = \mathcal{O}_{3/7}[1, 0] = \{1, 4, 0\}$.


 FIGURE 11. A schematic representation of an element of $TT(m/n, k_i, B)$

Consider the orbit $P_{1,1/3,3/7}$: by the statement of the theorem (with $p/q = 1/3$) it has data

$$\begin{aligned}
 & \overset{B}{(0, 0)} \rightarrow (1, 2) \rightarrow (4, 2) \rightarrow \overset{B}{(0, 2)} \rightarrow (2, 1) \rightarrow (5, 1) \rightarrow (1, 1) \\
 & \rightarrow (4, 1) \rightarrow \overset{A}{(0, 1)} \rightarrow (2, 2) \rightarrow (5, 2) \rightarrow (1, 3) \rightarrow (4, 3) \rightarrow \overset{C}{(0, 3)} \\
 & \rightarrow (3, 0) \rightarrow (6, 0) \rightarrow (2, 0) \rightarrow (5, 0) \rightarrow (1, 0) \rightarrow (4, 0)
 \end{aligned}$$

(see Fig. 12).


 FIGURE 12. The periodic orbit $P_{1,1/3,3/7}$

Proof. Suppose $P \in TT(m/n, k_i, B)$ has data $d(P) = ((N_r), \pi, (A, B, C))$. Let $\mu \leq N_0 - 1$ be greatest such that either $\mu \in A$ or $\pi_1(0, \mu) < k_{i+1}$. Then $\mu \in [1, N_0 - 2]$ since $0 \in B$ and $\pi_1(N_0 - 1, 0) = m$. It will be shown that the data of P must be as given in the statement of the theorem, with $p/q = \mu/(N_0 - 1)$. The proof is broken down into several short steps.

- i) $T \cap [k_i + 1, m - 1] = \emptyset$ (by the definition of m/n -admissibility of k_i).
- ii) If $r \in R$, then $N_r = 1$ and $\pi(r, 0) = (r +_n m, 0)$.
 Since $k_i < k_{i+1} < m$, TT b) gives $N_{k_{i+1} +_n m} = 1$, and applying LD a) i) inductively gives the result.
- iii) If $r \in T$ then $N_r = N_0$, and if $r \neq 0$ then $\pi(r, s) = (r +_n m, s)$ for all $s < N_0$.
 By i), if $j \in T$ with $j \neq 0$, the points of $P \cap e_{j+_n m}$ are precisely the images of the points of $P \cap e_j$.
- iv) $\pi(k_i +_n m, N_{k_i +_n m} - 1) = (k_i, N_0 - 1)$.
 $N_m = 1$ by ii), so $\pi(0, N_0 - 1) = (m, 0)$. Hence $N_0 - 1 \in \beta$ (by definition of the partition $\alpha \cup \beta$), and it follows from iii) that $(r, N_0 - 1) \in \beta$ for all $r \in T$: in particular, $(k_i, N_0 - 1) \in \beta$. By TT a), if $\pi_1^{-1}(k_i, N_0 - 1) = 0$ then $\pi_2^{-1}(k_i, N_0 - 1) \in A$, contradicting LD b) i), ii), v). Hence $\pi_1^{-1}(k_i, N_0 - 1) = k_i +_n m$, and the result follows.
- v) If $s < N_{k_i +_n m} - 1$, then $\pi(k_i +_n m, s) = (k_i, s)$ (by iv) and Lemma 2.17).

vi) $S \cap [k_i, m] = \{k_{i+1}\}$.

Suppose not: let k be the last element in the orbit segment from k_{i+1} to $k_i + n - m$ under addition of m modulo n which lies in $[k_i, m]$. Then $\mathcal{O}_{m/n}[k, 0]$ contains no elements of $[k + 1, m - 1]$, and hence $k \in \mathcal{A}_{m/n}$; moreover, $k_i < k < k_{i+1}$ by Lemma 2.28. This is a contradiction, since k_i and k_{i+1} are successive elements of $\mathcal{A}_{m/n}$.

vii) $[k_i + 1, m] \setminus \{k_{i+1}\} \subseteq R$ (by i) and vi)).

viii) $A = \{\mu\}$ and $\pi_1(0, s)$ is equal to k_i if $0 \leq s < \mu$, to k_{i+1} if $\mu \leq s < N_0 - 1$, and to m if $s = N_0 - 1$.

That $\pi_1(0, s)$ is equal to k_i, k_{i+1} or m for all s follows from ii) and vii). Then $A = \{\mu\}$ by the definition of μ , ii), and LD b) i), iii), v): the stated values of $\pi_1(0, s)$ are then immediate.

ix) If $r \in S$ then $N_r = N_0 - \mu$, and if $r \neq k_i + n - m$ then $\pi(r, s) = (r +_n m, s)$ for all $s < N_0 - \mu$.

By viii), $N_{k_{i+1}} = 1 + \#[\mu, N_0 - 2] = N_0 - \mu$. The result follows by vi) and LD a) i).

x) $\pi(0, \mu) = (k_{i+1}, N_0 - \mu - 1)$

By ii), iii), iv) and ix), there is a segment of the orbit of π as follows:

$$\begin{aligned} (0, s) &\rightarrow \underbrace{(k_{i+1}, N_0 - \mu - 1) \rightarrow \cdots \rightarrow (k_i + n - m, N_0 - \mu - 1)}_s \\ &\rightarrow \underbrace{(k_i, N_0 - 1) \rightarrow \cdots \rightarrow (0, N_0 - 1)}_T \rightarrow (m, 0) \end{aligned}$$

for some s . However $s \in A$ by TT'a), and hence $s = \mu$ by viii).

These results are enough to show that $d(P)$ is as given in the statement of the theorem with $p = \mu$ and $q = N_0 - 1$. Conversely, the data given in the statement defines an element of $TT(m/n, k_i, B)$ by construction, provided only that π is a cyclic permutation. The proof is therefore completed by showing that given two integers p and q with $0 < p < q$, the permutation π given in the statement is cyclic if and only if p and q are coprime. It is clear that for any (r, s) , there is some K such that $\tau_1(\pi^K(r, s)) = 0$; and that $\pi^{v_i}(0, 0) = (0, q - p) \neq (0, 0)$. Hence the cyclicity of π is equivalent to the cyclicity of the first return permutation of π on $\{(0, s) : 1 \leq s \leq q\}$. Let ρ be this first return permutation, i.e. $\rho(s) = \tau_2(\pi^{K_s}(0, s))$, where $K_s > 0$ is least such that $\tau_1(\pi^{K_s}(0, s)) = 0$ and $\tau_2(\pi^{K_s}(0, s)) \neq 0$. Then a straightforward calculation shows that ρ is a rotation of $[1, q]$ by $-p$, and hence is cyclic if and only if p and q are coprime. \square

The next step is to identify the horseshoe braid types of these orbits: it will be shown that $P_{i,p/q,m/n}$ has period $n + v_i p + v_{i+1}(q - p)$ and rotation interval $[\frac{u_i p + u_{i+1}(q - p)}{v_i p + v_{i+1}(q - p)}, \frac{m}{n}] = [\xi_{u_i/v_i, u_{i+1}/v_{i+1}}(p/q), m/n]$, and hence has height $\xi_{u_i/v_i, u_{i+1}/v_{i+1}}(p/q)$ and decoration $w_{m/n}$ by Lemma 1.17. In contrast to the situation in Section 2.5, there is no renormalization operator to provide a short cut to these rotation intervals: instead, they will be calculated by Markov partition techniques. Although the calculation is rather complicated the techniques are quite standard, and as such only a sketch proof, outlining the main steps, is given, so as to enable the enthusiastic reader to reconstruct the proof without too much difficulty.

Theorem 2.31. $P_{i,p/q,m/n} \in TT(m/n, k_i, B)$ has height $\xi_{u_i/v_i, u_{i+1}/v_{i+1}}(p/q)$ and decoration $w_{m/n}$.

Sketch Proof. That $P = P_{i,p/q,m/n}$ has period $n + v_i p + v_{i+1}(q - p)$ is immediate from Theorem 2.29 and the cardinalities of the sets R , S , and T . The rotation interval of P can be determined using Markov partition techniques, using the partition of Γ_n^P into intervals whose endpoints are the points of P and the valence n vertex v : the interval with endpoints (r, s) and either $(r, s + 1)$ or v is labelled $\langle r, s \rangle$. The Markov graph with these intervals as vertices can in principle be determined from the expression for π given in Theorem 2.29. Each loop in the Markov graph corresponds to a periodic orbit Q of $f_{m/n}$ whose braid type is forced by that of P : the rotation interval of P is the set of rotation numbers of such orbits about v . To calculate the rotation number corresponding to a given loop, one counts the number of times it goes around the star (i.e. the number of occurrences of $\langle r, s \rangle$ in the loop where $r \geq n - m$), and divides by its length.

It is clear that no such orbit can have rotation number greater than m/n , and that there is an orbit with this rotation number (namely the one given by the loop through the intervals with endpoint v). Thus it only remains to calculate the smallest possible rotation number; by standard arguments, this will be realized by a *minimal* loop (i.e. one which passes through each vertex at most once). The full Markov graph is too complicated to study in its entirety, so a sequence of simplifications is made.

- i) Every loop must contain $\langle 0, s \rangle$ for some s (i.e. $\{\langle 0, s \rangle: 0 \leq s \leq q\}$ is a *rome* [2] for the Markov graph): the structure of all loops is therefore given by considering only these intervals, and making a list of minimal paths between such intervals. Each such *basic* path has associated a rotation number (the number of occurrences of $\langle r, s \rangle$ with $r \geq n - m$ divided by its length), and the rotation number of a loop made by concatenating basic paths is the Farey sum of the corresponding rotation numbers.

Because of the structure of π (the only transition from $\langle r, s \rangle$ is to $\langle r +_n m, s \rangle$ unless $r = 0$, $k_i + n - m$, or $k_{i+1} + n - m$), most of these paths correspond to passing either through T , or through $S \cup T$, or through $R \cup S \cup T$: a short calculation (similar in spirit to the proof of Lemma 1.22) shows that the rotation numbers corresponding to these three types of path are $u_i/v_i < u_{i+1}/v_{i+1} < m/n$ respectively.

- ii) There are basic paths from $\langle 0, q \rangle$ to $\langle 0, s \rangle$ for all $s \in [0, q]$, each with rotation number m/n . Any loop passing through $\langle 0, q \rangle$ can be replaced by one with smaller rotation number, and $\langle 0, q \rangle$ can therefore be ignored.
- iii) There is a *rotation* loop which passes through each $\langle 0, s \rangle$ with $0 \leq s < q$ exactly once, made of basic paths from $\langle 0, s \rangle$ to $\langle 0, s -_q p \rangle$, each of which has rotation number either u_i/v_i or u_{i+1}/v_{i+1} . Hence given any loop which uses a basic path with rotation number greater than u_{i+1}/v_{i+1} , another loop of smaller rotation number can be constructed by replacing this basic path with a segment of the rotation loop. Thus the loop with minimal rotation number cannot use basic paths with rotation numbers greater than u_{i+1}/v_{i+1} .
- iv) The only basic paths which now remain to be considered are the following:

$$\langle 0, s \rangle \rightarrow \langle 0, q - p + s \rangle \text{ for } 0 \leq s \leq p - 1, \text{ with rotation } u_i/v_i$$

$$\langle 0, p \rangle \rightarrow \langle 0, t \rangle \text{ for } 0 \leq t \leq q - 1, \text{ with rotation } u_{i+1}/v_{i+1}$$

$$\langle 0, s \rangle \rightarrow \langle 0, s - p \rangle \text{ for } p + 1 \leq s \leq q - 2, \text{ with rotation } u_{i+1}/v_{i+1}$$

$$\langle 0, q - 1 \rangle \rightarrow \langle 0, t \rangle \text{ for } q - p - 1 \leq t \leq q - 1, \text{ with rotation } u_{i+1}/v_{i+1}.$$

One can then argue that basic paths from $\langle 0, p \rangle$ to $\langle 0, t \rangle$ for $t \neq 0$, or from $\langle 0, q - 1 \rangle$ to $\langle 0, t \rangle$ for $t \neq q - p - 1$ could be replaced by compound paths with lower rotation number. Hence the rotation loop realizes the minimum rotation number.

In the rotation loop, p of the basic paths have rotation number u_i/v_i , and the other $q - p$ have rotation number u_{i+1}/v_{i+1} . Hence $P_{i,p/q,m/n}$ has rotation interval

$$\left[\frac{u_i p + u_{i+1}(q - p)}{v_i p + v_{i+1}(q - p)}, \frac{m}{n} \right] = \left[\xi_{u_i/v_i, u_{i+1}/v_{i+1}}(p/q), \frac{m}{n} \right].$$

Since it has period $n + v_i p + v_{i+1}(q - p)$, the result follows by Lemma 1.17. \square

Hence $TT(m/n, k_i, B)$ consists of exactly one orbit with height r/s and decoration $w_{m/n}$ for each $r/s \in (u_i/v_i, u_{i+1}/v_{i+1})$. Combining this with the results of Section 2.7, $\bigcup_k TT(m/n, k, B)$ consists of one orbit with height r/s and decoration $w_{m/n}$ for each $r/s \in (0, m/n) \setminus \text{LFS}(m/n)$. The ‘missing’ heights $r/s \in \text{LFS}(m/n) \setminus \{0\}$ are supplied by orbits in $\bigcup_k TT(m/n, k, A)$, as described by Theorem 2.32 below. The case $\gamma = A$ is much easier than the case $\gamma = B$, since by Lemma 2.17 (and its proof), if $P \in TT(m/n, k, A)$ has data $d(P) = ((N_r), \pi, (A, B, C))$, then $k > 0$, $N_0 = 2$, and $\pi(0, 0) = (k, 1)$: thus $TT(m/n, k, A)$ is empty if $k \notin \mathcal{A}_{m/n}$, and has at most one element if $k \in \mathcal{A}_{m/n}$. The details are left to the reader.

Theorem 2.32. *Let $0 < i < \alpha$. Then $TT(m/n, k_i, A)$ consists of a single periodic orbit P , whose data $d(P) = ((N_r), \pi, (A, B, C))$ satisfies: $N_r = 2$ for $r \in T$, $N_r = 1$ for $r \notin T$, $A = \{0\}$, $B = \emptyset$, $C = \{1\}$, and*

$$\begin{aligned} \pi(0, 0) &= (k_i, 1) \\ \pi(0, 1) &= (m, 0) \\ \pi(r, s) &= (r +_n m, s) \text{ if } r > 0. \end{aligned}$$

P has height u_i/v_i and decoration $w_{m/n}$.

Example 2.33. The unique element of $TT(2/5, 1, A)$ is the $*$ -orbit of Example 2.8.

Combining the results of Theorems 2.19, 2.20, 2.22, 2.29, 2.31, 2.32, 2.18 and 1.8b) gives the main result of this paper:

Theorem 2.34. *Let $m/n \in (0, 1/2]$. Then $TT(m/n)$ consists exactly of one orbit $P_q^{m/n}$ of height q and decoration $w_{m/n}$ for each rational $q \in (0, m/n)$. The set of braid types of elements of $TT(m/n)$ is totally ordered by the forcing relation, with $\text{bt}(P_q^{m/n}) \leq \text{bt}(P_{q'}^{m/n})$ if and only if $q \geq q'$.*

3. STARS AND THE FULL HORSESHOE

The aim of this section is to clarify the relationship between the thick tree maps $F_{m/n}: T_n \rightarrow T_n$ and the full horseshoe $F_{1/2}: T_2 \rightarrow T_2$. The strongest connection is provided by *pruning theory*. Each thick tree map $F_{m/n}$ can be obtained from the full horseshoe by pruning: that is, by performing an isotopy which destroys all of the dynamics of the horseshoe in an open subset U of D^2 (after the isotopy every point of U is wandering), while leaving the dynamics unchanged elsewhere (the isotopy is supported on U). Although this process is conceptually valuable, it is not required in the main body of the paper: since it is rather complicated and assumes an understanding of the methods and results of [7, 8], the treatment given in Section 3.1 below is on an intuitive level.

The only aspect of the relationship between stars and the full horseshoe which was used in Section 2 is the fact that every $*$ -orbit of $f_{m/n}$ has the same braid type as some horseshoe periodic orbit. A rather straightforward proof of this is given in Section 3.2: it has the additional advantage of providing, for each $*$ -orbit P , the code of a horseshoe periodic orbit of the same braid type as P .

3.1. The pruning approach. As stated above, the approach taken throughout this subsection is intuitive: the interested reader is referred to [8] for details of the constructions outlined.

The aim is to start with the full horseshoe map $F_{1/2}$, to perform a sequence of isotopies which decrease the dynamics monotonically, remaining within the category of thick tree maps, and to arrive at a given thick tree map $F_{m/n}: T_n \rightarrow T_n$. Two distinct types of operation are used. The first redefines the thick tree structure (that is, the subset of D^2 which is regarded as the thick tree and the decomposition elements which give it its structure), but leaves the dynamics unchanged. The second is an isotopy which destroys the dynamics in some region which corresponds to an interval in the underlying tree endomorphism whose image backtracks.

Both of these operations have counterparts on the level of the underlying tree maps, and will be described on this level for the sake of both conceptual and diagrammatic simplicity. The important point is that the operations performed on the tree endomorphisms can be realised by isotopies of the corresponding homeomorphisms of the disk. Both the trees and their endomorphisms have additional structure due to the fact that they are induced by thick tree maps: in particular, the edges incident on a vertex have a cyclic order; and when several edge images backtrack over a common edge, there is a well-defined notion of the ‘innermost’ backtracking. This observation is, of course, reflected in the way that tree maps have been drawn throughout this paper.

Suppose that the image $f(e_r) = \dots \bar{e}_s e_s \dots$ of an edge e_r contains an innermost backtracking over the edge e_s . The following two operations can be performed:

Glueing e_r : Let I be the subinterval of e_r whose image is $\bar{e}_s e_s$. Identify those points of I which have the same image (so all points of I except the preimage of the initial point of e_s are identified in pairs). In the constructions described below, I always has the central vertex v of the star as an endpoint: glueing therefore preserves the star structure, but increases the valence of v by 1. It does not change the dynamics of f .

Pulling tight e_r : Delete the word $\bar{e}_s e_s$ from $f(e_r)$. This decreases the dynamics of f , while leaving the tree unchanged.

A second family of endomorphisms $g_{m/n}$ of Γ_n are needed for the construction: they differ from $f_{m/n}$ only in the image of the edge e_{n-1} .

Given $m/n \in (0, 1/2)$, let $g_{m/n}: \Gamma_n \rightarrow \Gamma_n$ be the tree map defined by

$$\begin{aligned} g(e_0) &= e_0 \bar{e}_1 e_1 \bar{e}_2 e_2 \dots \bar{e}_m e_m \\ g(e_r) &= e_{r+nm} \quad (1 \leq r \leq n-2) \\ g(e_{n-1}) &= e_{m-1} \bar{e}_m e_m. \end{aligned}$$

As an example, $g_{4/11}$ is depicted in Fig. 13. These maps are endowed with additional ‘2-dimensional’ structure as indicated in the figure.

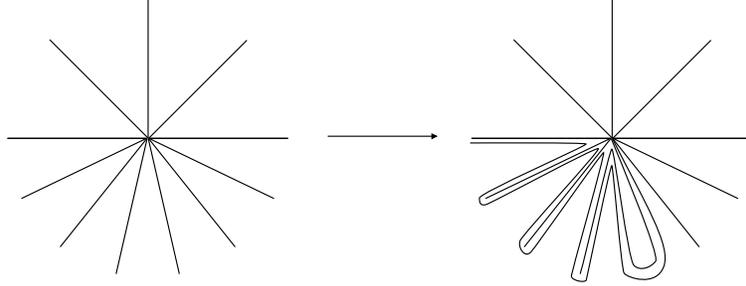


FIGURE 13. The tree map $g_{4/11}: \Gamma_{11} \rightarrow \Gamma_{11}$

The aim, then, is to construct the entire family $f_{m/n}$ of star maps starting from $f_{1/2}$ by applying the operations of glueing and pulling tight. This can be accomplished by a combination of two compound procedures (together with the observation that one can pass from $g_{m/n}$ to $f_{m/n}$ by pulling tight e_{n-1}).

Procedure L: starts with $f_{m/n}$ and yields $g_{p/q}$, where p/q is the immediate left Farey child of m/n .

Procedure R: starts with $g_{m/n}$ and yields $g_{u/v}$, where u/v is the immediate right Farey child of m/n .

To construct $f_{m/n}$ one navigates through the Farey graph from $1/2$ to m/n , passing at each step from parent to immediate child: at each step to the left (respectively right) one applies Procedure L (respectively Procedure R). Thus, for example, to obtain $f_{3/10}$ from $f_{1/2}$ one applies Procedure L to obtain $g_{1/3}$; pulls tight to obtain $f_{1/3}$; applies Procedure L to obtain $g_{1/4}$; applies Procedure R to obtain $g_{2/7}$; applies Procedure R to obtain $g_{3/10}$; and pulls tight to obtain $f_{3/10}$.

Procedure L: Start with $f_{m/n}$ and glue e_0 (that is, identify pairs of points in e_0 which have the same image in e_m). This creates a new edge, which is labelled e_n . Now the image of e_{n-m} backtracks over e_n ; glue e_{n-m} . This creates another new edge e_{n+1} , and the image of e_{n-2m} backtracks over it. Continue this procedure until edge e_{m-1} has been glued; each time an edge e_j is glued with $1 \leq j \leq m-1$, pull tight the innermost backtracking of the image of e_0 over e_j before proceeding. This yields $g_{p/q}$, where p/q is the immediate left Farey child of m/n .

Applying Lemma 1.22, observe that if $\text{LFP}(m/n) = a/b$, a total of $\#\mathcal{O}_{m/n}[m-1, 0] = b$ glueings are performed, and hence at the end of the procedure the star has $n+b = q$ edges; and $\#(\mathcal{O}_{m/n}[m-1, 0] \cap [0, m-1]) = a$ of the edges glued are between e_0 and e_m so that, after relabelling the edges in cyclic order, the image of e_0 crosses edges e_0 to $e_{m+a} = e_p$.

Example 3.1. Let $m/n = 3/7$: thus procedure L gives a construction of $g_{5/12}$ from $f_{3/7}$. The edges $e_0, e_4, e_1, e_5,$ and e_2 are glued successively: the new edges thereby created are labelled e_7, e_8, e_9, e_{10} and e_{11} respectively. Since $m-1 = 2$, the only pulling tight occur after glueing e_1 and e_2 . The procedure is shown in Fig. 14.

Procedure R: is similar. Start with $g_{m/n}$. Since $g_{m/n}(e_{n-1})$ backtracks over e_m one can successively glue $e_{n-1}, e_{n-m-1}, e_{n-2m-1}, \dots, e_m$, pulling tight e_0 each time an edge e_j is glued with $1 \leq j \leq m-1$. This yields $g_{u/v}$, where u/v is the immediate right Farey child of m/n .

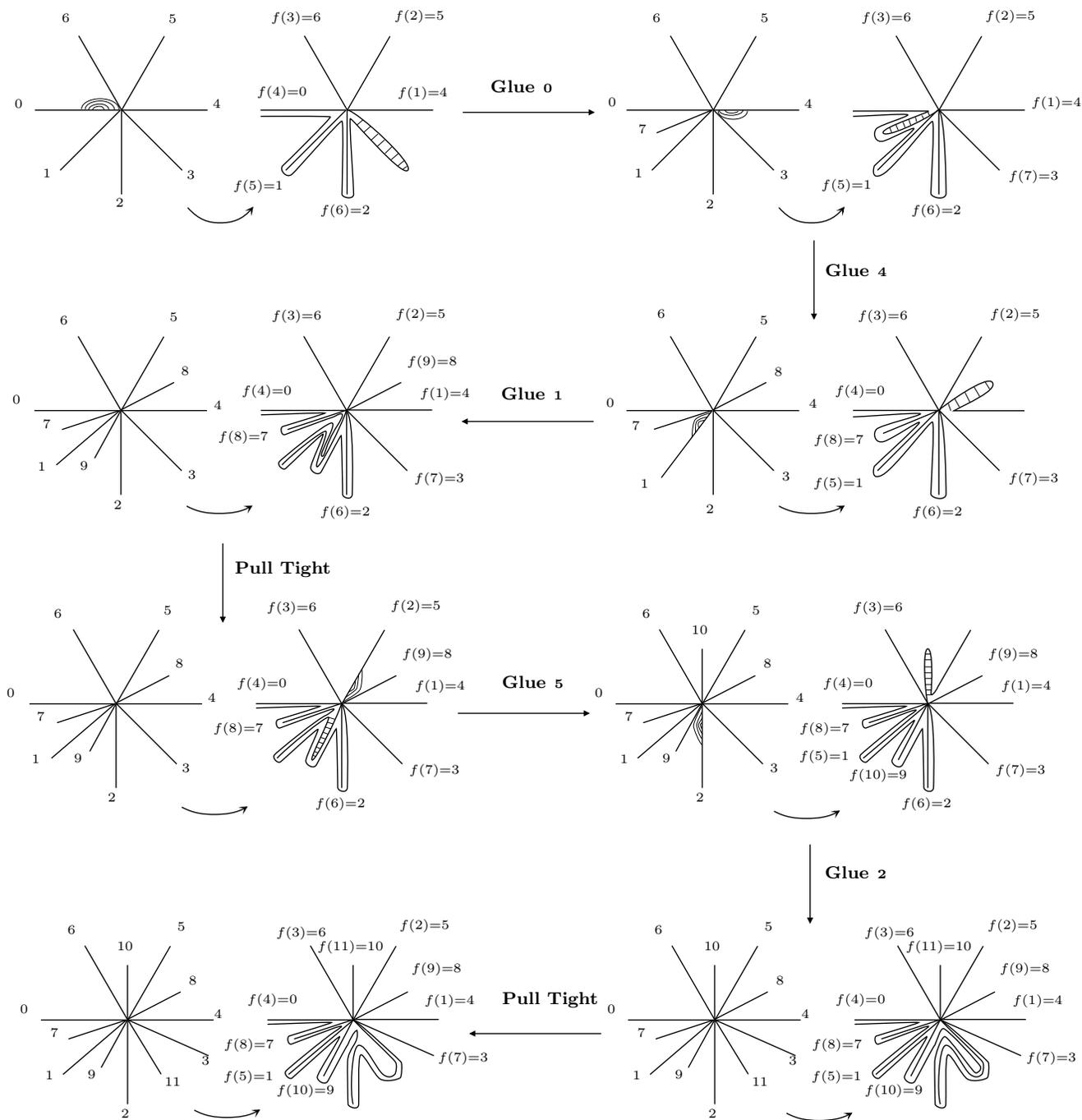


FIGURE 14. Constructing $g_{5/12}$ from $f_{3/7}$

3.2. Horseshoe symbolics for $*$ -orbits. In this section a more direct approach to the problem of showing that $*$ -orbits have horseshoe braid type is outlined. It makes use of the notion of the *line diagram* of an isotopy class of homeomorphisms of the punctured disk. For each $n \geq 2$, let D_n be a standard model of the n -punctured disk in which the punctures (or marked points) are equally spaced along the horizontal diameter of the disk. The line diagram of a homeomorphism $f : D_n \rightarrow D_n$ is the sequence $([f(\alpha_1)], \dots, [f(\alpha_{n-1})])$

of homotopy classes of the images of the horizontal arcs α_i joining the i^{th} to the $(i+1)^{\text{th}}$ puncture. It is straightforward to show that two homeomorphisms $f, g: D_n \rightarrow D_n$ are isotopic if and only if they have the same line diagram.

If $F: D^2 \rightarrow D^2$ is an orientation-preserving homeomorphism having a period n orbit P , then conjugating F so that the points of P coincide with the punctures of D_n yields a homeomorphism of D_n and thence a line diagram. Different conjugacies naturally give rise to different line diagrams, but periodic orbits P and Q of homeomorphisms F and G have the same braid type if and only if the conjugacies can be chosen so as to give the same line diagrams.

If F is the horseshoe, then there is a natural choice of (isotopy class of) conjugacy, namely one which sends vertical leaves to vertical leaves in an order-preserving manner. Such a conjugacy gives rise to ‘unimodal’ line diagrams which reflect the underlying unimodal structure of the horseshoe: for example, the periodic orbit of code 10010 gives rise to the line diagram of Fig. 9.

Thus to show that a $*$ -orbit P of $F_{m/n}$ has horseshoe braid type, it is necessary to construct a conjugacy which places the points of P along the horizontal diameter of D^2 in such a way that the resulting line diagram is unimodal. This will be achieved by drawing an arc through the points of the orbit which will be mapped by the conjugacy onto an interval of the horizontal diameter of D^2 . Not only does this show that P has horseshoe braid type, it also yields the code of a horseshoe periodic orbit of the same braid type as P : those points whose images are on the increasing segment of the unimodal line diagram are coded 0, while those on the decreasing segment are coded 1 – the critical point can be coded either 0 or 1.

The crucial observation is that if $\alpha_{m/n}$ denotes the path

$$\alpha_{m/n} = e_0 \bar{e}_1 e_1 \bar{e}_2 e_2 \dots \bar{e}_{n-m-1} e_{n-m-1} \bar{e}_{n-1} e_{n-1} \bar{e}_{n-2} e_{n-2} \dots \bar{e}_{n-m+1} e_{n-m+1} \bar{e}_{n-m},$$

then its image

$$f_{m/n}(\alpha_{m/n}) = e_0 \bar{e}_1 e_1 \dots \bar{e}_{n-1} e_{n-1} \bar{e}_{m-1} e_{m-1} \bar{e}_{m-2} e_{m-2} \dots \bar{e}_1 e_1 \bar{e}_0$$

can be obtained from the image

$$\begin{aligned} h_{m/n}(\alpha_{m/n}) &= \alpha_{m/n} \bar{\alpha}_{m/n} = e_0 \bar{e}_1 e_1 \bar{e}_2 e_2 \dots \bar{e}_{n-m-1} e_{n-m-1} \\ &(\bar{e}_{n-1} e_{n-1} \bar{e}_{n-2} e_{n-2} \dots \bar{e}_{n-m+1} e_{n-m+1}) \bar{e}_{n-m} e_{n-m} \bar{e}_{n-m+1} \dots \bar{e}_{n-1} e_{n-1} \\ &(\bar{e}_{n-m-1} e_{n-m-1} \dots \bar{e}_m e_m) \bar{e}_{m-1} e_{m-1} \dots \bar{e}_1 e_1 \bar{e}_0 \end{aligned}$$

of $\alpha_{m/n}$ under the horseshoe $h_{m/n}$ by removing the bracketed words (see Fig. 15, which illustrates this for $m/n = 2/5$). Notice that if Γ_n is drawn (as in this paper) so that e_0 and e_{n-m} are horizontal, then $\alpha_{m/n}$ passes along e_0 , then around each of the edges below the horizontal in the positive direction, then around each of the edges above the horizontal in the negative direction, and finally along e_{n-m} .

It follows that if P is a $*$ -orbit of $f_{m/n}$, then P has horseshoe braid type provided that an arc α projecting to $\alpha_{m/n}$ can be passed through the points of P in the thick tree T_n in such a way that $F_{m/n}(\alpha)$ projects to $f_{m/n}(\alpha_{m/n})$. Moreover, horseshoe symbolics for P can be obtained by coding with 0 those points of P which either lie in $e_0 \cup \dots \cup e_{n-2m-1}$, or lie in e_{n-2m} and have α pass through them with the orientation of \bar{e}_{n-2m} (these are precisely the points whose images lie in the increasing segment of the unimodal line diagram); and coding with 1 the other points of P .

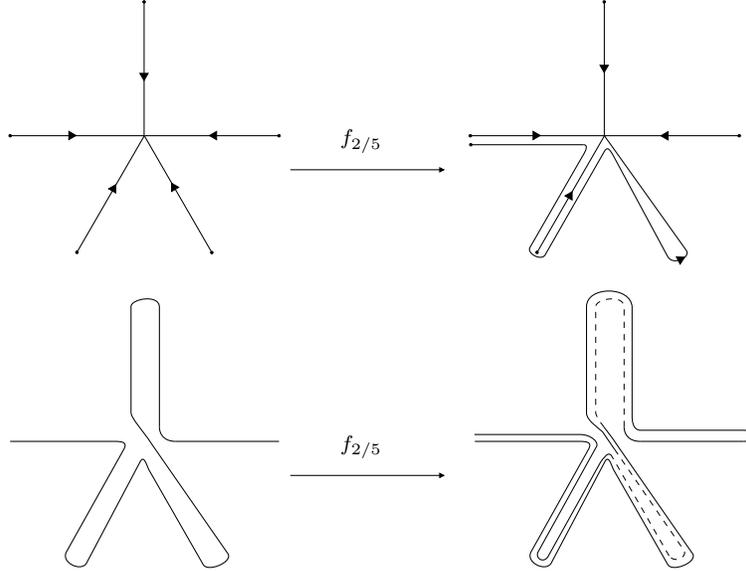


FIGURE 15. A path whose image under $f_{2/5}$ is a subhorseshoe

The only issue in constructing such an arc is to decide whether it should pass through each point of P with the orientation of e_r , or with the orientation of \bar{e}_r . The following partition of $\mathcal{L} \setminus \{(r, 0) : 0 \leq r < n\}$ gives the unique coherent way of doing this: points which α passes through with the orientation of e_r lie in I .

Definition 3.2. Let $P \in \mathcal{P}(m/n)$ have data $d(P) = ((N_r), \pi, (A, B, C))$. Define a partition

$$\mathcal{L} \setminus \{(r, 0) : 0 \leq r < n\} = I \cup O$$

inductively as follows. $\pi(0, 0)$ lies in I (respectively O) if $0 \in B$ (respectively $0 \in A$). For each i with $2 \leq i \leq \#P - n$, let $\pi^{i-1}(0, 0) = (r, s)$. Then $\pi^i(0, 0)$ and $\pi^{i-1}(0, 0)$ are in the same set if $r = 0$ and $s \in B$, or if $0 < r < n - 2m$. They are in different sets if $r = 0$ and $s \in A$, or if $n - 2m < r < n$. If $r = n - 2m$, then $\pi^i(0, 0) \in O$.

The following theorem then follows from the discussion above.

Theorem 3.3. Let P be a period N orbit in $\mathcal{P}(m/n)$ for some $m/n \neq 1/2$. Define a map $c: \mathcal{L} \rightarrow \{0, 1\}$ by

$$c(r, s) = \begin{cases} 0 & \text{if } r < n - 2m \text{ or } r = n - 2m \text{ and } (r, s) \in O \text{ or } (r, s) = (n - 2m, 0) \\ 1 & \text{if } r > n - 2m \text{ or } r = n - 2m \text{ and } (r, s) \in I. \end{cases}$$

Then P has the same braid type as the horseshoe periodic orbit of code

$$c(n - m, 0)c(\pi(n - m, 0))c(\pi^2(n - m, 0)) \dots c(\pi^{N-1}(n - m, 0)).$$

Example 3.4. Let $P \in \mathcal{P}(2/5, 1, A)$ be the periodic orbit with data

$$\begin{array}{cccccccc} (0, 0) & \rightarrow & (1, 1) & \rightarrow & (3, 1) & \rightarrow & (0, 1) & \rightarrow & (2, 0) & \rightarrow & (4, 0) & \rightarrow & (1, 0) & \rightarrow & (3, 0). \\ I & & O & & O & & I & & & & & & & & \end{array}$$

The following lemma will also be needed:

Lemma 4.2. *Let $m/n < 1/2$, and write $c_{m/n} = 10^{\kappa_1}1^20^{\kappa_2}1^2 \dots 1^20^{\kappa_m}1$ (where the $\kappa_i = \kappa_i(m/n)$ are given by formula (1) on page 5). Let $1 \leq r \leq m$. Then the word*

$$c = 10^{\kappa_r+1}1^20^{\kappa_r+1}1^2 \dots 1^20^{\kappa_m}1$$

disagrees with $c_{m/n}$ within the shorter of their lengths, and is greater than it in the unimodal order.

Proof. If the two words didn't disagree, then it would follow that $\kappa_r + 1 = \kappa_1$ and that $\kappa_m = \kappa_{m-r+1}$, contradicting the fact that $c_{m/n}$ is palindromic.

Observe that formula (1) gives, for each s with $1 \leq s \leq m+1-r$,

$$\sum_{i=1}^s \kappa_i = \left\lfloor \frac{sn}{m} \right\rfloor - (2s-1),$$

and

$$\begin{aligned} \kappa_r + 1 + \sum_{i=r+1}^{r+s-1} \kappa_i &= \left\lfloor \frac{(r+s-1)n}{m} \right\rfloor - \left\lfloor \frac{(r-1)n}{m} \right\rfloor - 2s + 1 \\ &\geq \left\lfloor \frac{(r+s-1)n}{m} - \frac{(r-1)n}{m} \right\rfloor - 2s + 1 \\ &= \left\lfloor \frac{sn}{m} \right\rfloor - (2s-1). \end{aligned}$$

Hence at the point where they first disagree c has a longer block of 0s than $c_{m/n}$, and so is greater in the unimodal order. \square

Lemma 1.17 *Let P be a period $N > 1$ orbit of the horseshoe with non-trivial rotation interval $\rho_i(P) = [u/v, m/n]$. Then $N \geq v+n$. Moreover, $N = v+n$ if and only if c_P is one of the four words $c_{u/v_1^0}w_{m/n_1^0}$ (or one of the two words c_{u/v_1^0} in the case $m/n = 1/2$).*

Proof. By induction on N , with the case $N = 2$ vacuous since the only period 2 horseshoe orbit has trivial rotation interval. The case $m/n = 1/2$ follows immediately from parts b) and d) of Lemma 4.1, so it will be assumed that $m/n < 1/2$. By parts b) and d) of Lemma 4.1, it then follows that $N \geq v+3$: let $k = N - v \geq 3$. By Theorem 1.8 a), the code of P is of the form

$$c_P = c_{u/v_1^0}w_1^0,$$

for some word w of length $k-3$. The four cases $c_P = c_{u/v}0w0$, $c_{u/v}0w1$, $c_{u/v}1w0$, and $c_{u/v}1w1$ need to be considered separately.

Case a) Suppose that

$$\begin{aligned} c_P &= c_{u/v}0w0 \\ &= 10^{\kappa_1}1^20^{\kappa_2}1^2 \dots 1^20^{\kappa_u} 10^{\lambda_1}1^{\nu_1}0^{\lambda_2}1^{\nu_2} \dots 1^{\nu_{t-1}}0^{\lambda_t}, \end{aligned}$$

where $\kappa_i = \kappa_i(u/v)$ for $1 \leq i \leq u$, and λ_i and ν_i are chosen to be positive for all i . Note that the initial word of c_P up to 0^{κ_u} has length v , and the complementary word has length k . For $1 \leq i \leq u$, let $[\xi_i^1, \eta_i^1]$ be the contribution to the rotation interval corresponding to the block

0^{κ_i} of 0s; and for $1 \leq i \leq t$, let $[\xi_i^2, \eta_i^2]$ be the contribution corresponding to 0^{λ_i} . Thus

$$\rho i(P) = \bigcup_{i=1}^u [\xi_i^1, \eta_i^1] \cup \bigcup_{i=1}^t [\xi_i^2, \eta_i^2].$$

Now for $2 \leq i \leq u$, $\eta_i^1 = q(10^{\kappa_{i-1}+1}1^2 0^{\kappa_{i-2}}1^2 \dots 1^2 0^{\kappa_1}10\dots)$, which, by Lemmas 4.2 and 4.1 c) and the fact that $c_{u/v}$ is palindromic, is less than u/v . Hence, since $\rho i(P) = [u/v, m/n]$, the interval $[\xi_i^1, \eta_i^1]$ must be empty. On the other hand, $\xi_1^1 = u/v$, and

$$\eta_1^1 = q(10^{\lambda_t}1^{\mu_{t-1}} \dots 1^{\mu_1}0^{\lambda_1}10\dots)$$

Let $\eta_1^1 = r/s$: then by Lemma 4.1 c), $k+1 \geq s+1$, or $k \geq s$. It follows that $\bigcup_{i=1}^u [\xi_i^1, \eta_i^1]$ is equal either to $\{u/v\}$, or to $[u/v, r/s]$ for some r/s with $s \leq k = N - v$.

Now consider the intervals $[\xi_i^2, \eta_i^2]$. First, $\eta_1^2 = u/v$, while $\xi_1^2 \geq u/v$ (since $\overline{c_P}$ has height u/v and $q: \{0, 1\}^{\mathbb{N}} \rightarrow (0, 1/2]$ is order-reversing). Hence $[\xi_1^2, \eta_1^2]$ is either empty or equal to $\{u/v\}$. By Lemma 4.1 c), all of the other $[\xi_j^2, \eta_j^2]$ are either empty or of the form $[a/b, c/d]$, where both b and d are less than k .

Hence $\rho i(P) = [u/v, m/n]$, where $n \leq k = N - v$, and the first statement of the lemma follows. For the second statement, observe from the argument above that if $n = k$ then

$$m/n = \eta_1^1 = q(10^{\lambda_t}1^{\mu_{t-1}} \dots 1^{\mu_1}0^{\lambda_1}10\dots),$$

and using Lemma 4.1 c) it follows that $w = w_{m/n}$ as required.

Case b) Suppose that $c_P = c_{u/v}0w1$ for some word w of length $k - 3$. If $c_{u/v}0w0$ is a maximal (non-repetitive) word, then it is the code of an orbit of the same braid type as P , and the result follows from the case a). If it is not maximal, then P is the period-doubling of a horseshoe orbit of half its period with the same rotation interval, and the result follows from the inductive hypothesis.

Case c) If $c_P = c_{u/v}1w0 = 10^{\kappa_1}1^2 0^{\kappa_2}1^2 \dots 1^2 0^{\kappa_u} 1^{\nu_1}0^{\lambda_1}1^{\nu_2}0^{\lambda_2}1^{\nu_3} \dots 1^{\nu_t}0^{\lambda_t}$, decompose $\rho i(P)$ into intervals $[\xi_i^1, \eta_i^1]$ for $1 \leq i \leq u$ and $[\xi_i^2, \eta_i^2]$ for $1 \leq i \leq t$ as in case a). Just as in that case, it can be shown that $[\xi_i^1, \eta_i^1]$ is empty for $2 \leq i \leq u$. Thus either η_1^1 or some η_i^2 is equal to m/n . By Lemma 4.1 a) and c), this means that the reverse $\widehat{c_P}$ of the code of P has the property that $\overline{\widehat{c_P}}$ contains one of the words $01_1^0 w_{m/n} 1_1^0$: equivalently (since $w_{m/n}$ is palindromic), $\overline{c_P}$ contains one of the words $1_1^0 w_{m/n} 1_1^0$. If there is such a word which is disjoint from the prefix $10^{\kappa_1}1^2 0^{\kappa_2}1^2 \dots 1^2 0^{\kappa_u}1$ of c_P , then the proof can be completed as in the case a). It remains to show, therefore, that if such a word overlaps the prefix, then the corresponding contribution to $\rho i(P)$ can not have right hand endpoint m/n . (It is not necessary to consider words contained entirely within the prefix, since it has already been shown that $\eta_i^1 < u/v$ for $i \geq 2$.)

Write $\kappa'_i = \kappa_i(m/n)$ for $1 \leq i \leq m$, and suppose first that $\overline{\widehat{c_P}}$ contains a word

$$010w_{m/n}01 = 010^{\kappa'_1}1^2 0^{\kappa'_2}1^2 \dots 1^2 0^{\kappa'_m}1$$

whose associated interval has right hand endpoint m/n , and suppose that the final block $0^{\kappa'_m}$ of 0s in this (reverse) word coincides with the block 0^{κ_r} of 0s in the prefix. Then $r \geq 2$, since otherwise P would have height m/n . By Lemma 4.1 a), the right hand endpoint of the associated interval being m/n gives

$$q(10^{\kappa'_1}1^2 \dots 1^2 0^{\kappa'_m}1^2 0^{\kappa_{r-1}}1^2 \dots 1^2 0^{\kappa_1}10\dots) = q(10w_{m/n}0110^{\kappa_{r-1}}1^2 \dots 1^2 0^{\kappa_1}10\dots) = m/n.$$

Hence by Lemma 4.1 c), $0^{\kappa_{r-1}}1^2 \dots 1^2 0^{\kappa_1} 10 \dots \succ \overline{w_{m/n} 011}$, or equivalently

$$10^{\kappa_{r-1}+1} 1^2 0^{\kappa_{r-2}} 1^2 \dots 1^2 0^{\kappa_1} 10 \dots \prec \overline{10 w_{m/n} 011}.$$

Applying Lemma 4.1 c) again gives $q(10^{\kappa_{r-1}+1} 1^2 0^{\kappa_{r-2}} 1^2 \dots 1^2 0^{\kappa_1} 10 \dots) \geq m/n > u/v$, so that

$$10^{\kappa_{r-1}+1} 1^2 0^{\kappa_{r-2}} 1^2 \dots 1^2 0^{\kappa_1} 10 \dots \prec 10^{\kappa_1} 1^2 0^{\kappa_2} 1^2 \dots 0^{\kappa_m} 10 \dots,$$

contradicting Lemma 4.2.

Exactly the same argument works if the contributing word in $\overline{c_P}$ is $011w_{m/n}01$. If the word is one of $01_1^0 w_{m/n} 11$, then comparing the overlapping segments of this word and the prefix of P gives $\kappa_r = \kappa'_m - 1$ and hence (using again that $w_{m/n}$ is palindromic)

$$\begin{aligned} 10^{\kappa'_1} 1^2 0^{\kappa'_2} 1^2 \dots 1^2 0^{\kappa'_{m-r}} 1^2 0^{\kappa'_{m-r+1}} \dots &= 10^{\kappa_r+1} 1^2 0^{\kappa_{r+1}} 1^2 \dots 1^2 0^{\kappa_{m-1}} 1^2 0^{\kappa_m} \dots \\ &\succ 10^{\kappa_1} 1^2 0^{\kappa_2} 1^2 \dots 1^2 0^{\kappa_m} \dots, \end{aligned}$$

(where the final inequality is by Lemma 4.2). Taking heights gives

$$m/n = q(10^{\kappa'_1} 1^2 0^{\kappa'_2} 1^2 \dots 1^2 0^{\kappa'_{m-r}} 1^2 0^{\kappa'_{m-r+1}} \dots) < q(10^{\kappa_1} 1^2 0^{\kappa_2} 1^2 \dots 1^2 0^{\kappa_m} \dots) = u/v,$$

a contradiction.

Case d) The proof for $c_P = c_{u/v} 1w1$ follows from case c) in the same way that case b) follows from case a). \square

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