## ALGORITHM 72

## COMPOSITION GENERATOR

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procedure comp ( $c, k$ ); value $k$; integer array $c$; integer k ;
comment Given a $k$-part composition e of the positive integer $n$, comp generates a consequent composition if there is one. If comp operates on each consequent composilion after it is found, all compositions will be gencrated, provided that $1,1, \ldots, 1$, $n-k+1$ is the initial $c$. If $c$ is of the form $n-k+1,1,1, \ldots, 1$, there is no consequent, and $c$ will be replaced by a $k$ vector of O's. Reference: John Riordan, An Introduction to Combinatorial Analysis, John Wiley and Sons, Ine., New York, 1958, Chapter 6;
hegin integer $j$; integer array $d[1: k]$;
if $k=1$ then go to last;
for $\mathrm{j}:=1$ step 1 until k do $\mathrm{d}[\mathrm{j}]:=\mathrm{c}[\mathrm{j}]-1$;
test: if $\mathrm{d}[j]>0$ then go to set; $\mathrm{j}:=\mathrm{j}-1$; go to if $j=1$ then last else test;
set: $d[j]:=0$; $\mathrm{d}[\mathrm{j}-1]:=\mathrm{d}[\mathrm{j}-1]+1$; $\mathrm{d}[\mathrm{k}]:=\mathrm{c}[\mathrm{j}]-2$; for $j:=1$ step 1 until $k$ do $e[j]:=d[j]+1$; go to exit;
last: for $\mathrm{j}:=1$ step 1 until k do $\mathrm{c}[\mathrm{j}]:=0$;
exit: end comp

## CERTIFICATION OF ALGORITHM 42

INVERT (T. C. Wood, Comm. ACM, Apr., 1961)
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INVERT was hand-coded for the LGP-30 using machine language and the 24.0 floating-point interpretive system, which carries 24 bits of significance for the fractional part of a number and five bits for the exponent. The following changes were found necessary:
(a) if $\mathrm{j}=\mathrm{n}+1$ then $\mathrm{a}[\mathrm{i}, \mathrm{j}]:=1.0$ else $\mathrm{a}[\mathrm{i}, \mathrm{j}]:=0.0$;
should be
if $\mathrm{j}=\mathrm{n}+\mathrm{i}$ then $\mathrm{a}[\mathrm{i}, \mathrm{j}]:=1.0$ else $a[\mathrm{i}, \mathrm{j}]:=0.0$;
(b) for $\mathrm{k}:=\mathrm{j}$ step 1 until $2 \times \mathrm{n}$ do $\mathrm{a}[\mathrm{i}, \mathrm{k}]:=\mathrm{a}[\mathrm{i}, \mathrm{k}] / \mathrm{a}[\mathrm{i}, \mathrm{j}] ;$
should be
for $k:=2 \times n$ step -1 until $i d o$ $a[i, k]:=a[i, k] / a[i, i] ;$
(c) if $1 \neq \mathrm{i}$ then for $\mathrm{k}:=1$ step 1 until $2 \times \mathrm{n}$ do $a[l, k]:=a[l, k]-a[i, k] \times a[l, j] ;$ should be if $1 \neq \mathrm{i}$ then for $\mathrm{k}:=2 \times \mathrm{n}$ step -1 untilido $a[l, k]:=a[l, k]-a[i, k] \times a[l, i] ;$

Given these changes, $j$ becomes superfluous in the second i loop, and the other references to $j$ may be changed to references to $i$.

INVERT obtained the following results:
The computer inverted a 17 -by-17 matrix whose elements were integers less than ten in absolute value. When the matrix and its inverse were multiplied together, the largest nondiagonal element in the product was - 000003 . Most nondiagonal elements were less than 00001 in absolute value.

INVERT was tested using finite segments of the Hibert matrix.
The following results were obtained in the $4 \times 4$ case:

| 16.005 | -120.052 | 240.125 | -140.082 |
| ---: | ---: | ---: | ---: |
| -120.052 | 1200.584 | -2701.407 | 1680.917 |
| 240.126 | -2701.411 | 6483.401 | -4202.217 |
| -140.082 | 1680.920 | -4202.219 | 2801.446 |

The exact inverse is:

$$
\begin{array}{rrrr}
16 & -120 & 240 & -140 \\
-120 & 1200 & -2700 & 1680 \\
240 & -2700 & 6480 & -4200 \\
-140 & 1680 & -4200 & 2800
\end{array}
$$

INVERT was also coded for the LGP 30 in machine language and the 24.1 extended range interpretive system. This system, which uses 30 significant bits for the fraction, obtained the lollowing as the inverse of the $4 \times 4$ Hilbert matrix:

| 16.000 | -120.001 | 240.001 | -140.001 |
| ---: | ---: | ---: | ---: |
| -120.001 | 1200.006 | -2700.015 | 1680.010 |
| 240.001 | -2700.016 | 6480.037 | -4200.024 |
| -140.001 | 1680.010 | -4200.024 | 2800.016 |

The program coded in the 24.0 interpretive system successfully inverted a matrix consisting of ones on the minor diagonal and zeros everywhere else.

## REMARK ON ALGORITHM 52

A SET OF TEST MATRICES (John IR. Herndon, Comm. ACM, Apr. 1961)
G. H. Dubay

University of St. Thomas, Houston, Tex.
In the assignment stalement

$$
\mathrm{c}:=\mathrm{t} \times(\mathrm{t}+1) \times(\mathrm{t}+\mathrm{t}-5) / 6 ;
$$

the $t$ is undefined. A suitable definition would be provided by preceding (a) with $t:=n$;

## CERTIFICATION OF ALGORITHM 68

AUGMENTATION (H. (. . Rice, Comm. ACM, Aug. 1961)
L. M. Breed

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AUGMENTATION was transliterated into BALGOL for the Burroughs 220 , and proved successful in a number of test cases. However, the following algorithm has exactly the same effect and is considerably simpler:
real procedure $\operatorname{Aug}(x, y) ;$ value $x, y ;$ integer $x, y$;
begin if $x<0$ then $I$ : go to $L$ else Aug $:=x+y$ end Aug

Contributions to this department must be in the form stated in the Algorithms Department policy statement (Communications, February, 1960) except that ALGO1, 60 notation should be used (see Communications, May, 1960). Contributions should be sent in duplicate to J. H. Wegstein, Computation Laboratory, National Bureau of Standards, Washington 25, D. C. Algorithms should be in the Publication form of ALGOL 60 and written in a style patterned after the most recent algorithms appearing in this department.

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